

# Combinatorics, 2017 Fall, USTC

## Week 13 Note 2

2017.12.15, Friday

### Probabilistic Methods

#### • Ideas:

- ① Image we need to find some combinatorial object satisfying certain property, call them “good” objects. We consider a random object. If the probability that the random object is “good” is positive, then there must exist “good” objects.
- ② To compute the probability of being “good”, we often compute the probability of being “bad” and aim to prove the probability is strictly less than 1.

**Definition 1.** A finite probability space is  $(\Omega, \Pr)$ , where  $\Omega$  is a finite set and  $\Pr : \Omega \rightarrow [0, 1]$  such that  $\sum_{x \in \Omega} \Pr[x] = 1$ .

- $A \subseteq \Omega$  is called an event,  $\Pr[A] = \sum_{x \in A} \Pr[x]$ .
- $\Pr[\emptyset] = 0$ ,  $\Pr[\Omega] = 1$ .
- $\Pr[A \cup B] = \Pr[A] + \Pr[B] - \Pr[A \cap B] \leq \Pr[A] + \Pr[B]$  (union bound).  
“=” holds if  $A \cap B = \emptyset$ .
- Two events  $A$  and  $B$  are **independent** if  $\Pr[A \cap B] = \Pr[A] \cdot \Pr[B]$ .
- A random variable is a function  $X : \Omega \rightarrow \mathbb{R}$ .
- $\forall A \subseteq \Omega$ , its indicator random variable  $X_A$  is defined as

$$X_A(x) = \begin{cases} 1, & x \in A \\ 0, & x \notin A \end{cases}.$$

- **Expectation:**  $E[X] = \sum_{\omega \in \Omega} \Pr[\omega] X(\omega)$ .
- **linearity of expectation:**  $E[aX + bY] = aE[X] + bE[Y]$ ,  $a, b \in \mathbb{R}$ .

**Theorem 1.** Let  $n, s \in \mathbb{Z}$  satisfying  $\binom{n}{s} 2^{1-\binom{s}{2}} < 1$ , then  $R(s, s) > n$ .

**Proof.** We need to find a 2-edge-coloring of  $K_n$  such that it has NO monochromatic clique  $K_s$  (i.e. a clique with all edges of the same color). Consider a

random 2-edge-coloring of  $K_n$ : each edge is colored blue or red, each with probability  $\frac{1}{2}$ , independent of other edges.

Let  $A$  be the event that the defined  $K_n$  has a monochromatic  $K_s$ . For each  $B \in \binom{[n]}{s}$ , let  $A_B$  be the event that  $K_n$  has a monochromatic  $K_s$  with vertex set  $B$ . Then  $A = \bigcup_{B \in \binom{[n]}{s}} A_B$ .

$$\Pr[A] = \Pr\left[\bigcup_{B \in \binom{[n]}{s}} A_B\right] \leq \sum_{B \in \binom{[n]}{s}} \Pr[A_B] = \binom{n}{s} 2^{1-\binom{s}{2}} < 1. \text{ Thus } \Pr[\bar{A}] =$$

$1 - \Pr[A] > 0$ , i.e. the probability that  $K_n$  has NO monochromatic  $K_s$  is positive. So there must exist a 2-edge-coloring of  $K_n$  such that it has NO monochromatic  $K_s$ .  $\square$

**Corollary 2.**  $R(k, k) \geq \frac{1}{e\sqrt{2}} k 2^{\frac{k}{2}}$ .

**Proof.** Let  $n = \frac{1}{e\sqrt{2}} k 2^{\frac{k}{2}} \left(\frac{e}{2}\right)^{\frac{1}{k}}$ . Recall  $\binom{n}{k} < \frac{n^k}{k!}$  and  $k! > e \left(\frac{k}{e}\right)^k$ . We have

$$\binom{n}{k} 2^{1-\binom{k}{2}} < \frac{n^k}{e \left(\frac{k}{e}\right)^k} 2^{1-\binom{k}{2}} = \left(\frac{en}{k}\right)^k \left(\frac{2}{e}\right)^k 2^{-\binom{k}{2}}.$$

Since  $n = \frac{1}{e\sqrt{2}} k 2^{\frac{k}{2}} \left(\frac{e}{2}\right)^{\frac{1}{k}}$ , we have  $\binom{n}{k} 2^{1-\binom{k}{2}} < 1$ . Then  $R(k, k) > n \geq \frac{1}{e\sqrt{2}} k 2^{\frac{k}{2}}$ .  $\square$

**Definition 2.** A **tournament** of  $n$  vertices is an orientation of  $K_n$ . Say vertex  $x$  beats vertex  $y$  if  $(x, y) \in E$ . Say a tournament  $T$  has the property  $P_k$ , if for any  $k$ -subset  $A \subset V$ , there is a vertex who beats all vertices in  $A$ .

**Question:** For  $\forall k \geq 2$ , does there exists a  $T$  with property  $P_k$ ?

**Theorem 3.**  $\forall k \geq 2$ , if  $\binom{n}{k} \left(1 - \frac{1}{2^k}\right)^{n-k} < 1$ , then there exists a tournament with  $n$  players with property  $P_k$ .

**Proof.** Consider a random tournament of  $n$  players. For any  $i < j$ , the arc  $i \rightarrow j$  occurs with probability  $\frac{1}{2}$ , independent of other choices. Let  $B$  be the event that  $T$  doesn't satisfy  $P_k$ . For  $A \in \binom{[n]}{k}$ , let  $B_A$  be the event that  $T$  satisfies that all vertices in  $[n] \setminus A$  can not beat every vertices in  $A$ . Then  $B = \bigcup_{A \in \binom{[n]}{k}} B_A$ .

For  $x \in [n] \setminus A$ , let  $B_{A,x}$  be the event that  $T$  satisfies that  $x$  can not beat every vertices in  $A$ . Then  $B_A = \bigcap_{x \in [n] \setminus A} B_{A,x}$ . Clearly,  $\Pr[B_{A,x}] = 1 - \left(\frac{1}{2}\right)^k$ .

Note that only arcs between  $x$  and  $A$  will affect the event  $B_{A,x}$ , and these arcs for distinct  $x$ 's are disjoint. Thus all events  $B_{A,x}$ 's are independent,  $x \in [n] \setminus A$ . So  $\Pr[B_A] = \prod_{x \in [n] \setminus A} \left(1 - \frac{1}{2^k}\right) = \left(1 - \frac{1}{2^k}\right)^{n-k}$ .

By union bound,

$$\Pr[B] \leq \sum_{A \in \binom{[n]}{k}} \Pr[B_A] = \binom{n}{k} \left(1 - \frac{1}{2^k}\right)^{n-k} < 1.$$

Thus  $\Pr[\bar{B}] > 0$ , i.e.  $\exists$  a tournament  $T$  with property  $P_k$ .  $\square$

**Corollary 4.** If  $n \geq k^2 2^{k+1}$ , then  $\exists$  a tournament  $T$  with  $n$  players with property  $P_k$ .

**Proof.** Exercise.  $\square$

**Definition 3.**  $\mathcal{F} \subseteq 2^{[n]}$ , say  $\mathcal{F}$  is 2-colorable if  $\exists$  a function  $f : [n] \rightarrow \{\text{blue}, \text{red}\}$ , s.t. every set  $A \in \mathcal{F}$  is not monochromatic.

**Note:**  $\mathcal{F} \subseteq \binom{[n]}{2}$ ,  $\mathcal{F}$  is 2-colorable iff  $G = ([n], \mathcal{F})$  is bipartite.

**Theorem 5.** Every  $k$ -uniform family  $\mathcal{F}$  with  $|\mathcal{F}| < 2^{k-1}$  is 2-colorable.

**Proof.** Let  $\mathcal{F}$  be any fixed  $k$ -uniform family of subsets of some finite set  $X$ . Consider a random function  $f : X \rightarrow \{\text{blue}, \text{red}\}$ , such that each  $x \in X$  is colored by blue or red with probability  $\frac{1}{2}$ , and the coloring of different elements are independent. Let  $B$  be the event that  $f$  is "bad", i.e.  $\exists A \in \mathcal{F}$  is monochromatic.

For  $A \in \mathcal{F}$ , let  $B_A$  be the event that  $A$  is monochromatic. So  $B = \bigcup_{A \in \mathcal{F}} B_A$ .  
 $\forall A \in \mathcal{F}$ ,  $\Pr[B_A] = 2^{1-k}$ . Hence  $\Pr[B] \leq \sum_{A \in \mathcal{F}} \Pr[B_A] = |\mathcal{F}| 2^{1-k} < 1$ . So  $\Pr[\bar{B}] > 0$ , i.e.  $\exists$  a 2-coloring s.t. no  $A \in \mathcal{F}$  is monochromatic.  $\square$

**Theorem 6.** If  $k$  is sufficiently large, then there exists a  $k$ -uniform family  $\mathcal{F}$  such that  $|\mathcal{F}| \leq k^2 2^k$  and  $\mathcal{F}$  is not 2-colorable.

**Proof.** To be continued.  $\square$

Let  $B(k)$  be the minimum possible number of sets in a  $k$ -uniform family with is not 2-colorable. Then  $2^{k-1} \leq B(k) \leq k^2 2^k$ .