

# Combinatorics 2017 Fall

## week11 note

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**Recall Fisher's Inequality:** Let  $A_1, \dots, A_m$  be subsets of  $X$  such that  $|A_i \cap A_j| = k$  for all  $i \neq j \in [m]$ , then  $m \leq n$ .

**Def:** Let  $\mathcal{F} \subset 2^X$ , and  $L \subset \{0, 1, \dots\}$  be a finite set of integers. say  $\mathcal{F}$  is  $L$ -intersecting if  $|A \cap B| \in L$  for all  $A \neq B \in \mathcal{F}$ . Note that Fisher's inequality tells that  $|\mathcal{F}| \leq n$ , if  $|L| = 1$ .

**Thm1:(Frankl-Wilson)** If  $\mathcal{F}$  is an  $L$ -intersecting family of subsets of an  $n$ -elements set, then  $|\mathcal{F}| \leq \sum_{i=0}^{|L|} \binom{n}{i}$ . [Ex: Show the bound is best possible.]

**Def:** A function space is the set of all functions from  $\Omega \rightarrow F$ . A set of functions  $f_1, f_2, \dots, f_m$  is linearly independence if  $\exists \lambda_i$  s.t.  $\lambda_1 f_1 + \lambda_2 f_2 + \dots + \lambda_m f_m = 0$ , then  $\lambda_i = 0, i \in [m]$ .

**Lemmal:(Independence criterion)** If  $i \in [m]$ , Let  $f_i : \Omega \rightarrow F$ , (where  $F$  is a field) be functions and  $v_i \in \Omega$ . such that

- (i)  $f_i(v_i) \neq 0, \forall i \in [m]$ .
- (ii)  $f_i(v_j) = 0, \forall 1 \leq j < i \leq m$ .

Then  $f_1, \dots, f_m$  are linearly independent in the function space  $F^\Omega$ .

**proof:** If there is a nontrivial linear relation

$$\lambda_1 f_1 + \lambda_2 f_2 + \dots + \lambda_m f_m = 0$$

Suppose  $i$  is the smallest index such that  $\lambda_i \neq 0$ . then

$$0 = \lambda_{i+1}f_{i+1}(v_i) + \lambda_{i+2}f_{i+2}(v_i) + \cdots + \lambda_n f_n(v_i) = -\lambda_i f_i(v_i) \neq 0$$

□

**proof of Thm1:** Consider  $\{f(x_1, \dots, x_n)$  polynomials with degree  $\leq d\}$ , then each of  $f$  is combination of  $x_1^{t_1} \cdots x_n^{t_n}$ , with  $t_1 + \cdots + t_n \leq d$ . Suppose  $\mathcal{F} = \{A_1, \dots, A_m\}$ , with  $|A_1| \leq \cdots \leq |A_m|$ , and  $L = \{l_1, \dots, l_s\}$ . That is,  $\forall i \neq j, \exists k \in [s], s.t. |A_i \cap A_j| = l_k$ . Let  $v_i$  be the indicator vector of  $A_i, i \in [m]$ , then  $\langle v_i, v_j \rangle = |A_i \cap A_j|$ . For  $i = 1, \dots, m$ . define  $f_i$  with  $n$  variables by

$$f_i : \mathbb{R}^n \rightarrow \mathbb{R}.$$

$$f_i(x) = \prod_{k: l_k < |A_i|} (\langle v_i, x \rangle - l_k).$$

if  $j < i$ , then  $f_i(v_j) = 0$  and  $f_i(v_i) \neq 0, i \in [m]$ . Then  $f_1, \dots, f_m$  are linearly independent over  $\mathbb{R}$ .

$f_i$  are polynomials of degree at most  $s$ . so  $x_1^{r_1} \cdots x_n^{r_n}$  with  $r_1 + \cdots + r_n \leq s$  are basic monomials, which has in total  $\sum_{i=0}^s \binom{i+n-1}{i} = \binom{s+n}{s}$ , but we can do it better! define new polynomials  $\bar{f}_i$  from  $f_i$  by replacing all term  $x_i^k$  by  $x_i$  since  $v_i$  are 0-1 vectors. we have  $\bar{f}_i(v_i) = f_i(v_i), \forall i, j$  so  $\bar{f}_1, \dots, \bar{f}_m$  are linearly independent. who lie in a space with basis  $x_1^{r_1} \cdots x_n^{r_n}$  with  $r_1 + \cdots + r_n \leq s$  and  $r_i \in \{0, 1\}$ . so  $m \leq \sum_{i=0}^s \binom{n}{i}$ . □

**Thm2:** Let  $p$  be a prime and  $L \subset \mathbb{Z}_p = \{0, 1, \dots, p-1\}$ . Assume  $\mathcal{F} = \{A_1, \dots, A_m\} \subset 2^{[n]}$  such that

- (a)  $|A_i| \notin L \pmod{p}, \forall i \in [m]$ .
- (b)  $|A_i \cap A_j| \in L \pmod{p}, \forall i \neq j$ .

$$\text{Then } |\mathcal{F}| \leq \sum_{i=0}^{|L|} \binom{n}{i}$$

[Hint: define  $f_i(x) = \prod_{l \in L} (\langle v_i, x \rangle - l) \pmod{p}, i \in [m]$ .]

**Recall Ramsey number**  $R(s, t)$

1.  $R(t, t) > (t - 1)^2$ .
2.  $R(s, t) > \Omega(t^3)$  1972, Zsigmond Nagy
3.  $R(t, t) > \Omega(2^{\frac{t}{2}})$ .
4.  $R(t, t) > t^{\Omega(\ln t / \ln \ln t)}$ . 1977, Frankl, 1981 F&Wlison.

A graph is Ramsey graph with respect  $t$ , if it has no clique of size  $t$  and no independent set of size  $t$ .

**Thm3:** For any prime  $p$ , there is a graph  $G$  on  $n = \binom{p^3}{p^2 - 1}$  vertices s.t. the size of maximum clique or maximum independent set is  $\leq \sum_{i=0}^{p-1} \binom{p^3}{i}$ .

**proof:** Let  $G = (V, E)$  be as follows

- $V = \binom{[p^3]}{p^2 - 1}$
- for  $A, B \in V$ ,  $A \sim B$  iff  $|A \cap B| \not\equiv p - 1 \pmod{p}$ .

consider the clique  $A_1, \dots, A_m \in V$ ,  $|A_i| = p^2 - 1 \equiv p - 1 \pmod{p}$ ,  $|A \cap B| \not\equiv p - 1 \pmod{p}$  means  $|A \cap B| \in L \pmod{p}$ . where  $L = \{0, 1, \dots, p - 2\} \subset \mathbb{Z}_p$ .

By Thm2, we have  $m \leq \sum_{i=0}^{p-1} \binom{p^3}{i}$ .

Now consider an independent set  $B_1, \dots, B_s$ ,  $|B_i \cap B_j| \equiv p - 1 \pmod{p}$ ,

so  $|B_i \cap B_j| \in \{p - 1, 2p - 1, \dots, p(p - 1) - 1\} = L^* \subset \mathbb{Z}_{\geq 0}$ . and  $|L^*| = p - 1$ , By Thm1, we have  $s \leq \sum_{i=0}^{p-1} \binom{p^3}{i}$ .  $\square$

**Colloary4:**  $R(t + 1, t + 1) > t^{\Theta(\ln t / \ln \ln t)}$ .

**proof:** Let  $t = \sum_{i=0}^{p-1} \binom{p^3}{i}$ ,  $V = \binom{p^3}{p^2 - 1}$ .

Recall  $(\frac{n}{k})^k \leq \binom{n}{k} \leq (\frac{en}{k})^k$

$$t \approx \binom{p^3}{p} \approx (p^2)^p = p^{2p}, V \approx \binom{p^3}{p^2} \approx p^{p^2},$$

$$\ln t \approx 2p \ln p = O(p \ln p), \ln \ln t \approx \ln p.$$

$$\Rightarrow p = \Theta(\frac{\ln t}{\ln \ln t}), V \approx (p^{2p})^{p/2} \approx t^{\theta(\ln t / \ln \ln t)}.$$

□

Recall a vector space over a field  $(\mathbb{C}, \mathbb{R}, \mathbb{Q}, \mathbb{F}_q)$

**Thm5:(Odd/Even Town)** Let  $\mathcal{F} \subset 2^{[n]}$ , s.t.  $|A|$  is odd for all  $A \in \mathcal{F}$ , and  $|A \cap B|$  is even for  $A \neq B \in \mathcal{F}$ , Then  $|\mathcal{F}| \leq n$ .

**proof:**  $\forall A \in \mathcal{F}$ , let  $e_A$  be the indicator vector of  $A$ , consider  $e_A$  as a vector in  $\mathbb{F}_2^n$ , Then

$$\langle e_A, e_A \rangle = 1, \quad \forall A \in \mathcal{F}$$

$$\langle e_A, e_B \rangle = 0, \quad \forall A \neq B \in \mathcal{F}$$

Let  $|\mathcal{F}| = m$ , we will show vectors  $e_A, A \in \mathcal{F}$  are linearly independent in  $\mathbb{F}_2^n$ . If  $\alpha_A \in \{0, 1\}$ , s.t.  $\sum_{A \in \mathcal{F}} \alpha_A e_A = 0$

$$\Rightarrow \langle e_A, \sum_{A \in \mathcal{F}} \alpha_A e_A \rangle = 0 \Rightarrow \alpha_A = 0$$

This shows  $e_A, A \in \mathcal{F}$  are linearly independent, so  $|\mathcal{F}| \leq n$ . □

**Colloary6:**  $R(t+1, t+1) > \binom{t}{3}$ .

**proof:** Let  $G = (V, E)$  be as follows

$$V = \binom{[t]}{3}, A \sim B, \text{ iff } |A \cap B| = 1,$$

Consider a clique  $A_1, \dots, A_m$ .  $|A_i \cap A_j| = 1$  by Fisher's inequality,  $m \leq t$ .

Consider an independent set  $B_1, \dots, B_s$

$$|B_i| = 3 \quad \text{Odd}, |B_i \cap B_j| = 0 \text{ or } 2 \quad \text{Even}.$$

By Odd/Even Thm,  $s \leq t$

□

**2017.12.1**

**Graham-Pollak Thm:**

Biclique:i.e. a complete bipartite graph  $K_{A,B}$  with vertex set  $A \cup B$ , and edge set  $A \times B$ .

Def:Let  $G$  be a graph and  $G_1, \dots, G_t$  be subgraph of  $G$ , say  $G_1, \dots, G_t$  is an edge-decomposition of  $G$  if each edge of  $G$  occurs in exactly one subgraph of  $G_i$ .

**Thm1:(Graham-Pollak Thm)** The edges of  $K_n$  can not be decomposed into fewer than  $n - 1$  edge-disjoint bicliques.

[Best possible:define  $K_{A_i, B_i}$  with  $A_i = \{i\}, B_i = \{i + 1, \dots, n\}, i = 1, \dots, n - 1$ , then  $K_{A_i, B_i}, i \in [n - 1]$  form a decomposition].

**proof:**Associate each vertex  $i$  of  $K_n$  with an indeterminate  $x_i$  and each edges  $i \sim j$  with  $x_i x_j$ , Let  $S(x) = \sum_{1 \leq i < j \leq n} x_i x_j$ . suppose  $K_{A_i, B_i}, i \in [t]$  is an edge decomposition of  $K_n$ .  $\forall i \in [t]$ , Let  $L_i(x) = \sum_{j \in A_i} x_j, R_i(x) = \sum_{j \in B_i} x_j$ , then  $L_i(x) \cdot R_i(x) = (\sum_{j \in A_i} x_j)(\sum_{j \in B_i} x_j) = \sum_{\{i, j\} \in K_{A_i, B_i}} x_i x_j$ . and

$$S(x) = \sum_{i=1}^t L_i(x) \cdot R_i(x).$$

$$\text{Let } T(x) = \left(\sum_{i=1}^n x_i\right)^2 - 2 \sum_{i < j} x_i x_j = \left(\sum_{i=1}^n x_i\right)^2 - 2S(x) = \left(\sum_{i=1}^n x_i\right)^2 - 2 \sum_{i=1}^t L_i(x) \cdot R_i(x).$$

consider  $t + 1$  linear equations over  $\mathbb{R}$ .  $L_i(x) = 0, i \in [t], \sum_{i=1}^n x_i = 0$ , If  $t \leq n - 2$ , then it has a nonzero solution  $x \in \mathbb{R}^n$ , Put this vector into equation  $T(x) = \left(\sum_{i=1}^n x_i\right)^2 - 2 \sum_{i=1}^t L_i(x) \cdot R_i(x)$ , we have  $LRH = T(x) \neq 0$ . Contradiction!  $\square$

Fix a field  $F$ , Let  $F[x_1, \dots, x_n] = \{\text{multivariate polynomial } f : F^n \rightarrow F\}$ .

A polynomial  $f$  vanishes on  $E \subset F^n$  if  $f(x_1, \dots, x_n) = 0, \forall (x_1, \dots, x_n) \in E$ .

A point  $(x_1, \dots, x_n) \in F^n$  with  $f(x_1, \dots, x_n) = 0$  is a root of  $f$ .

A polynomial  $f$  is a zero polynomial if all coefficients are 0.

Recall: univariate case  $f(x) \neq 0, \deg(f) = d$ , then  $f$  has  $\leq d$  roots.

1. If  $0 \neq f$  vanishes on a subset  $S$ , then  $|S| \leq \deg(f)$ .
2. If  $\nexists f \neq 0$  that vanishes on  $S$ , then  $|S| > \deg(f)$ .

**Lemma1:** Given  $E \subset F^n$  of size  $|E| < \binom{n+d}{d}$ , there exists a  $0 \neq f \in F[x_1, \dots, x_n]$  with  $\deg(f) \leq d$  that vanishes on  $E$ .

**proof:** Let  $V_d = \{f \in F[x_1, \dots, x_n], \deg(f) \leq d\}$ , then dimension of  $V_d$  is  $\binom{n+d}{d}$ .

Let  $F^E$  be the set of vectors of length  $|E|$ , where each component of a vector  $u \in F^E$  has value in  $F$  and is indexed by an element in  $E$ , that is  $u = (u_x)_{x \in E}$ , and  $u_x \in F$ .

Consider the evaluation map  $V_d \rightarrow F^E, f \mapsto (f(a))_{a \in E}$ . since dimension of  $F^E$  is  $|E| < \binom{n+d}{d} = \dim(V_d)$ .

We have the map is non-injective. i.e.  $\exists f_1 \neq f_2$ , such that  $(f_1(a))_{a \in E} = (f_2(a))_{a \in E}$ , then  $f = f_1 - f_2$  vanishes on  $E$   $\square$

**Lemma2:**  $\forall 0 \neq f(x_1, \dots, x_n) \in \mathbb{F}_q[x_1, \dots, x_n]$  with  $\deg(f) = d$ , has at most  $dq^{n-1}$  roots. ( $d \leq q$ , if  $d > q$  trivial.)

**proof:** Assume  $n \geq 2$ , Write  $f = g + h$ , where  $g$  is homogenous of degree  $d$ , and  $\deg(h) \leq d-1$ . Since  $f \neq 0, g \neq 0, \exists \omega \in \mathbb{F}_q^n \setminus \{0\}, g(\omega) \neq 0$ .

$\forall u \in \mathbb{F}_q^n$ , Let  $L_u = \{u + t\omega : t \in \mathbb{F}_q\}, |L_u| = q$

If  $v \notin L_u$  then  $L_u \cap L_v = \emptyset$ . Hence  $\mathbb{F}_q^n$  is partitioned into  $q^n/q = q^{n-1}$  lines. It remains to show that the number of roots of  $f$  on each line is at most  $d$ .

$\forall u \in \mathbb{F}_q^n$ , define a univariate polynomial  $p_u(t) = f(u + t\omega)$ , then  $\deg(p_u) \leq d$ . since the coefficient  $t^d$  in  $p_u(t)$  is  $g(\omega) \neq 0$ ,  $p_u(t) \neq 0$ . so  $p_u(t)$  has at most  $d$  roots in  $L_u$ , implying that  $f$  has at most root in  $L_u$ . since we have  $q^{n-1}$  lines,  $f$  has at most  $dq^{n-1}$  roots in  $\mathbb{F}_q^n$   $\square$

**Lemma3:**  $\forall S \subset F, |S| \geq d, \forall 0 \neq f \in F[x_1, \dots, x_n]$  of degree  $d$  can have at most  $d|S|^{n-1}$  roots in  $S^n$ .

**proof:** Induction on  $n$ , the number of variables in  $f$ .

$n = 1$  is true, let  $n \geq 2$ . write  $f$  according to the powers of  $x_n$ .  $f = f_0 + f_1x_n + \dots + f_tx_n^t, t \leq d, f_t \neq 0$ . and  $f_i, i = 0, 1, \dots, t$  are polynomials in  $F[x_1, \dots, x_{n-1}]$ .

Now estimate the number of points  $(a, b) \in S^{n-1} \times S, s.t. f(a, b) = 1$ .

1.  $f_t(a) = 0, \deg(f_t) \leq d-t$ , by assumption,  $f_t$  at most  $(d-t)|S|^{n-2}$  roots in  $|S|^{n-1}$ , so there are at most  $(d-t)|S|^{n-1}$  points  $(a, b) \in S^{n-1} \times S$ , for which  $f(a, b) = 0$  and  $f_t(a) = 0$ .
2.  $f_t(a) \neq 0$  Fix a point  $a \in S^{n-1}$  satisfying  $f_t(a) \neq 0$ . then  $g_a(x_n) = f(a, x_n)$  is of degree  $t$ , So  $g_a(x_n)$  has at most  $t$  roots, since there are at most  $|S|^{n-1}$  such points  $a$ , the number of points  $(a, b) \in S^{n-1} \times S$ , for which  $f(a, b) = 0$  and  $f_t(a) \neq 0$ . is at most  $t|S|^{n-1}$ .

Together, there are  $\leq d|S|^{n-1}$  points  $(a, b) \in S^{n-1} \times S$ , for which  $f(a, b) = 0$   $\square$

### Combinatorial Nullstellensatz

**Lemma4:** Let  $f \in F[x_1, \dots, x_n]$  be a polynomial, and let  $t_i$  be the maximum degree of  $x_i$  in  $f$ . Let  $S_i \subset F$  with  $|S_i| \geq t_i + 1$ . If  $f(x) = 0, \forall x \in S_1 \times \dots \times S_n$ , then  $f$  is zero polynomial.

**proof:** By induction.  $n = 1$  is true, Assume that the lemma holds for  $n - 1$ , we prove it for  $n (n \geq 2)$ . Write  $f$  as a polynomial in

$x_n, f = \sum_{i=0}^{t_n} f_i(x_1, \dots, x_{n-1})x_n^i$ . where in each  $f_i$ , the maximum degree of  $x_j$  is  $t_j, j \in [n - 1]$ .

For each fixed  $(n - 1)$ -tuple  $(x_1, \dots, x_{n-1}) \in S_1 \times \dots \times S_{n-1}$ , Let  $g(x_n) = f(x_1, \dots, x_{n-1}, x_n)$  vanishes on  $S_n$ , Hence  $g(x_n) \equiv 0$  That is all coefficients  $f_i(x_1, \dots, x_{n-1})$  of  $g(x_n)$  is zero, for all  $(x_1, \dots, x_{n-1}) \in$

$S_1 \times \cdots \times S_{n-1}$ , Hence by induction hypothesis,  $f_i \equiv 0, \forall i$ , which implies that  $f \equiv 0$ .  $\square$