

Partially Ordered Sets

Let X be a finite set.

Definition 1. R is called a "relation" on the set X , if $R \subseteq X \times X$ where $X \times X = \{(x_1, x_2) : x_1, x_2 \in X\}$. Denote the Cartesian product if $(x, y) \in R$, then xRy .

Definition 2. A partially ordered set (poset for short) is an ordered pair (X, R) , where X is a finite set and R is a relation on X such that the following holds:

- (1) R is reflective: xRy for $\forall x \in X$
- (2) R is antisymmetric: if xRy and yRx , then $x = y$
- (3) R is transitive: if xRy and yRz , then xRz

Examples. Consider the poset $(2^{[n]}, \subseteq)$, where " \subseteq " denotes the inclusion relationship.

We often use " \preceq " to replace the use of " R ". So $(X, R) \Rightarrow (X, \preceq)$.

If $x \preceq y$ but $x \neq y$, then $x \prec y$, and we say x is predecessor of y .

Definition 3. Let (X, \preceq) be a poset, we say element x is an immediate predecessor of y , if

- (1) $x \prec y$

(2) NO $t \in X$ s.t. $x \prec t \prec y$

If x is an immediate predecessor of y , then we write $x \triangleleft y$.

Fact: For $x, y \in (X, \preceq)$, $x \prec y$ if and only if there exists $x_1, x_2, \dots, x_k \in X$ s.t. $x \triangleleft x_1 \triangleleft x_2 \triangleleft \dots \triangleleft x_k \triangleleft y$ (Note that $k = 0$ i.e. $x \triangleleft y$).

Proof. (\Leftarrow) trivial

(\Rightarrow) For $x \prec y$, let $M_{xy} = \{t \in X : x \prec t \prec y\}$. We prove by induction on $|M_{xy}|$. Because $|M_{xy}| = 0 \Rightarrow x \triangleleft y$. So suppose it holds for $x \prec y$ with $|M_{xy}| < n$. Consider $x \prec y$ with $|M_{xy}| = n \geq 1$. Pick any $t \in M_{xy}$, consider M_{xt} , M_{ty} . Clearly $M_{xt} \subsetneq M_{xy}$ and $M_{ty} \subsetneq M_{xy}$ (Because of transitivity). By induction on M_{xt} , M_{ty} , there exists $x_1, x_2, \dots, x_k \in X$ and $y_1, y_2, \dots, y_l \in X$ s.t. $x \triangleleft x_1 \triangleleft x_2 \triangleleft \dots \triangleleft x_k \triangleleft t$ and $t \triangleleft y_1 \triangleleft y_2 \triangleleft \dots \triangleleft y_l \triangleleft y$, $\Rightarrow x \triangleleft x_1 \triangleleft x_2 \triangleleft \dots \triangleleft x_k \triangleleft t \triangleleft y_1 \triangleleft \dots \triangleleft y_l \triangleleft y$. We are done. \blacksquare

One property that posets have is that we can express them in diagrams

Definition 4. The Hasse diagrams of a poset (X, \preceq) is a drawing in the plane such that

- (1) Each element of X is drawn as a node in the plane
- (2) Each pair x, y with $x \triangleleft y$ is connected by a line segment
- (3) If $x \triangleleft y$, then the node x must appear lower in the plane than the node y

The fact that $x \prec y$ iff $x \triangleleft x_1 \triangleleft x_2 \triangleleft \dots \triangleleft x_k \triangleleft y$ now can be restated as follows: $x \prec y$ if and only if we can find a path in the Hasse diagram from node x to node y , strictly from bottom to top.

Definition 5. Let (X, \preceq) and (X', \preceq') be two posets. A mapping $f : X \rightarrow X'$ is called an embedding of (X, \preceq) and (X', \preceq') if

- (1) f is injective
- (2) $f(x) \preceq f(y)$ iff $x \preceq y$.

Theorem 6. For every poset (X, \preceq) there exists an embedding into the poset $\mathcal{B}_X = (2^X, \subseteq)$

Proof. Consider the mapping $f : X \rightarrow 2^X$ by $f(x) = \{y \in X : y \preceq x\}$. Let us verify that such f is an embedding of (X, \preceq) into \mathcal{B}_X .

Firstly, f is injective

Suppose $f(x) = f(y)$ for $x, y \in X \Rightarrow x \in f(y) \Rightarrow x \preceq y$, similarly $y \preceq x$.

Thus $x = y$.

Secondly, $f(x) \subseteq f(y)$ iff $x \preceq y$.

If $x \preceq y$, then $\forall t \in f(x)$ has $t \preceq x \preceq y \Rightarrow t \in f(y) \Rightarrow f(x) \subseteq f(y)$.

If $f(x) \subseteq f(y)$, then $x \in f(x) \subseteq f(y) \Rightarrow x \preceq y$ ■

Definition 7. Let $P = (X, \preceq)$ be a poset.

- (1) For distinct $x, y \in X$, if $x \prec y$ or $y \prec x$, then we say x, y are comparable. Otherwise x, y are incomparable
- (2) The set $A \subseteq X$ is an antichain of P , if any two elements of A are incomparable. Let $\alpha(P)$ be the maximum size of an antichain in P
- (3) The set $A \subseteq X$ is a chain of P , if any two elements of A are comparable. Let $\omega(P)$ be the maximum size of a chain of P

Consider the Hasse diagram, $\omega(P)$ means the max length of a path (from bottom to top) in this diagram. So $\omega(P)$ is also called the height of P . And $\alpha(P)$ is called the width of P .

Definition 8. An element $x \in X$ is minimal in $P = (X, \preceq)$, if x has NO predecessor in P .

Fact: The set of minimal elements of $P = (X, \preceq)$ forms an antichain of P .

Theorem 9. For \forall poset $P = (X, \preceq)$, $\alpha(P) \cdot \omega(P) = |X|$

Proof. We will inductively define a sequence of poset P_i and set M_i for $1 \leq i \leq l$ s.t. M_i is the set of minimal elements of $P_i = (X_i, \preceq)$ and $X_i = X - \bigcup_{j=1}^{i-1} M_j$ as following. First, set $P_1 = P = (X, \preceq)$, $X_1 = X$ and $M_1 = \emptyset$. Assume posets $P_i = (X_i, \preceq)$ and M_{i-1} are defined for all $1 \leq i \leq k$. Let $M_i = \{ \text{all minimal elements of } P_i \}$ and let $X_{i+1} = X - M_1 \cup \dots \cup M_i$. Then let P_{i+1} be the subposet of P restricted on X_{i+1} . We keep doing this until $X_{l+1} = \emptyset$. By Fact 2, each M_i for $1 \leq i \leq l$ is an antichain of P_i and thus it is also an antichain of P . So $|M_i| \leq \alpha(P)$ ■