

## Basic of Graphs

**Definition 1.** A *graph*  $G = (V, E)$  consists of a vertex set  $V$  and an edge set  $E$ , where the elements of  $V$  are called *vertices* and the elements of

$$E \subseteq \binom{V}{2} = \{(x, y) : x, y \in V\}$$

are called edges.

If  $E$  contains unordered pairs, then  $G$  is an undirected graph, otherwise,  $G$  is directed graph. In this course, all graphs are undirected.

In this course, all graphs are simple, that is, it contains no loops and multiple edges.

We say vertices  $i$  and  $j$  are adjacent if  $(i, j) \in E$ , write as  $i \sim_G j$ . We say the edge  $(i, j)$  is incident to the endpoints  $i$  and  $j$ . Let  $e(G) = \#$  edges in graph  $G = (V, E)$ , i.e.,  $e(G) = |E|$ . The degree of a vertex  $v$  in graph  $G$  denote by  $d_G(v)$ , is the number of edges of  $G$  incident to  $v$ .

The neighborhood of a vertex  $v$  is the set of vertices  $u$  s.t.  $u$  and  $v$  are adjacent, i.e.,  $N_G(v) = \{u \in V(G) | u \sim_G v\}$ .

$$\Rightarrow d_G(v) = |N_G(v)|.$$

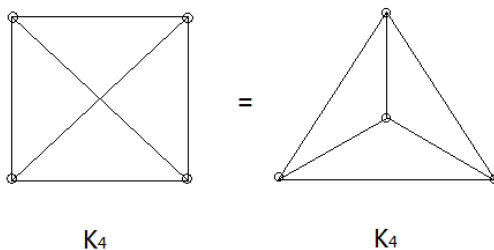
A graph  $G' = (V', E')$  is a subgraph of  $G = (V, E)$  if  $V' \subseteq V$  and  $E' \subseteq E \cap \binom{V'}{2}$ , and we write as  $G' \subseteq G$ .

A graph with  $n$  vertices is a complete graph or a clique denote by  $K_n$  if all pairs of vertices are adjacent. So  $e(K_n) = \binom{n}{2}$ .

A graph with  $n$  vertices is called an independent set, denoted by  $I_n$ , if it contains no edge at all.

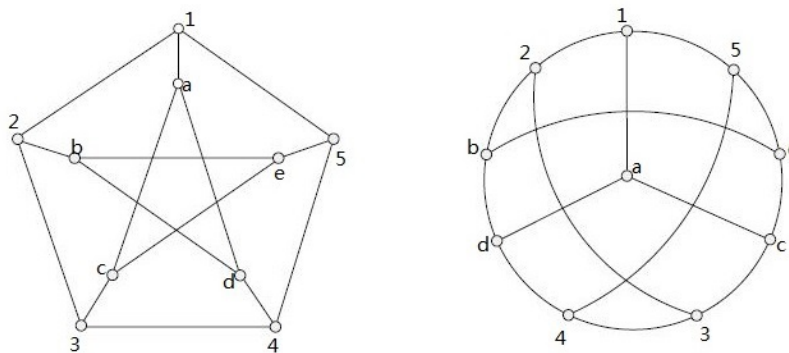
Given a graph  $G = (V, E)$ , its complement is a graph  $\bar{G} = (V, E^c)$  with the same vertex set  $V$  such that  $E^c = \binom{V}{2} \setminus E$ . So clearly  $e(G) + e(\bar{G}) = \binom{n}{2}$  where  $n = |V(G)|$ .

**Definition 2.** Two graphs  $G = (V, E)$  and  $G' = (V', E')$  are *isomorphic* if there exists a bijection  $f : V \rightarrow V'$  such that  $i \sim_G j$  if and only if  $f(i) \sim_{G'} f(j)$ .



$K_4$

$K_4$

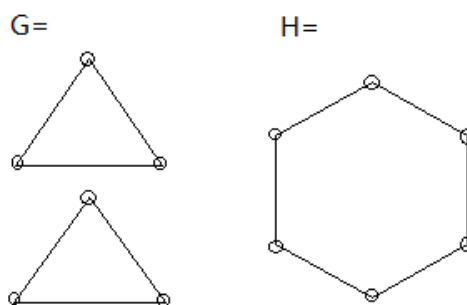


Petersen graph

The degree sequence of a graph  $G = (V, E)$  is a sequence of degrees of all vertices listed in a non-decreasing order.

**Problem.** If two graphs  $G$  and  $H$  have the same degree sequence, then they are isomorphic?

The answer is NO.

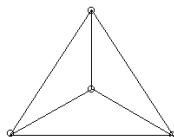


**Definition 3.** A *path*  $P_k$  of length  $k - 1$  is a graph  $v_1 - v_2 - v_3 - \dots - v_k$ , where  $v_i \sim v_{i+1}$ . Note that the length means the number of edges it contains.

**Definition 4.** A *cycle*  $C_k$  of length  $k$  is a graph consisting of  $V(C_k) = \{v_1, v_2, \dots, v_k\}$ ,  $E(C_k) = \{(v_1, v_2), (v_2, v_3), \dots, (v_{k-1}, v_k), (v_k, v_1)\}$ .

**Definition 5.** A graph  $G$  is a *planar* graph, if we can draw  $G$  on the plane such that its edges intersect only at their endpoints.

**Example.**  $K_4$  is planar.



**Exercise.**  $K_5$  is not planar.

**Lemma 6** (Hand-shaking Lemma). *In any  $G = (V, E)$ ,*

$$\sum_{v \in V} d(v) = 2|E|.$$

*Proof.* Let  $F = \{(e, v) : e \in E(G), v \in V(G) \text{ s.t. } e \text{ is incident to } v\}$ , i.e.,  $v$  is one of the two endpoints of  $e$ . Then

$$|F| = \sum_{e \in E(G)} 2 = 2|E|, \quad |F| = \sum_{v \in V(G)} d(v) \Rightarrow \text{Done.}$$

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**Corollary 7.** *In any graph  $G$ , the number of vertices with odd degree is even.*

*Proof.* Let  $\mathcal{O} = \{v \in V(G) \text{ s.t. } d(v) \text{ is odd}\}$ ,  $\mathcal{E} = \{v \in V(G) \text{ s.t. } d(v) \text{ is even}\}$ . Then

$$\begin{aligned} 2|E| &= \sum_{v \in V} d(v) = \sum_{v \in \mathcal{O}} d(v) + \sum_{v \in \mathcal{E}} d(v) \\ &\Rightarrow \sum_{v \in \mathcal{O}} d(v) \text{ is even} \Rightarrow |\mathcal{O}| \text{ is even.} \end{aligned}$$

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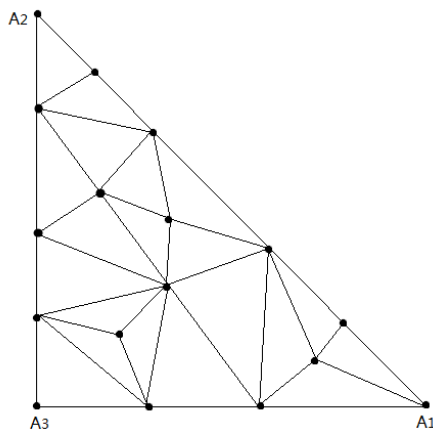
**Corollary 8.** *In any graph  $G$ , if there exists a vertex with odd degree, there are at least two vertices with odd degree.*

Let us consider the following application.

Let us draw a large triangle  $A_1, A_2, A_3$  in the plane, with 3 vertices  $A_1, A_2, A_3$ . Then we divided this large triangle arbitrarily into small triangles, s.t. no triangles can have a vertex inside an edge of any other triangle.

Then we assign colors 1, 2, 3 to all vertices of these triangles, under the following rules:

- (1). The vertex  $A_i$  is assigned the color  $i$  for  $\forall i \in 1, 2, 3$ .
- (2). All vertices lying on the edges  $A_iA_j$  of the large triangle must be assigned the color  $i$  or  $j$ .
- (3). All interior vertices can be assigned any color 1, 2, or 3.



**Lemma 9** (Sperner's Lemma (a planar version)). *For any assignment of colors described as above, there always exists a small triangle whose three vertices are assigned all three colors 1, 2, 3.*

We call such triangle as a rainbow triangle.

$V(G) = \{ \text{the face of any small triangles, and the outer face} \}$ . We define the edges of  $G$  as follows:

Two vertices of  $G$ , i.e. 2 faces of the drawing, are adjacent in  $G$  if they are neighboring faces and the two endpoints of their common edge have colors 1 and 2.