

Basic of Graphs

Lemma 1 (Sperner's Lemma (a planar version)). *For any assignment of colors described as above, there always exists a small triangle whose three vertices are assigned all three colors 1, 2, 3.*

Proof. Define an auxiliary graph G :

- Its vertices are the faces of small triangle and the outer face. Let z be the vertex representing the outer face.
- Two vertices of G are adjacent, if the two corresponding faces are neighboring faces and the two endpoints of their common edge are colored by 1 and 2.

We consider the degree of a vertex $v \in V(G) - \{z\}$

- (1). If the face of v has NO two vertices with color 1 and 2, then $d_G(v) = 0$.
- (2). The face of v has 2 vertices with color 1 and 2. Let the color of the third vertex be k . If $k \in \{1, 2\}$, then $d_G(v) = 2$. Otherwise $k = 3$, then $d_G(v) = 1$

We claim that: $d_G(v)$ is odd iff $d_G(v) = 1$, iff the face of v has colors 1,2,3.

Then we consider $d_G(z)$ and we claim it must be odd. Why? The edge of G incident to z obviously have to go across A_1A_2 . Consider the sequence of

the colors of the points on A_1A_2 , from A_1 to A_2 . So $d_G(z) = \#$ of alternations between 1 and 2 in this sequence, which must be odd.

By Corollary, since the graph G has the vertex z with odd degree, there must be another vertex $v \in V(G) - \{z\}$ of odd degree. Then the face of v has colors 1,2,3. ■

Theorem 2 (Brouwer's Fixed Point Theory in 2-dimension). *Every continuous function $f : \Delta \rightarrow R$ has a fixed point x , that is $f(x) = x$.*

Proof. Consider a sequence of refinement of Δ . Define three auxiliary functions $\beta_i : \Delta \rightarrow R$ for $i \in \{1, 2, 3\}$ as following: for $\forall a = (x, y) \in \Delta$,

$$\begin{cases} \beta_1(a) = x \\ \beta_2(a) = y \\ \beta_3(a) = 1 - x - y \end{cases}$$

For a given continuous $f : \Delta \rightarrow \Delta$, define $M_1 = \{a \in \Delta : \beta_1(a) \geq \beta_1(f(a))\}$ for $i \in \{1, 2, 3\}$

Fact 1: $\forall a \in \Delta, \exists i \in \{1, 2, 3\}$ s.t. $a \in M_i$

Fact 2: if $a \in M_1 \cap M_2 \cap M_3$, then a is a fixed point.

We can define a coloring $C : \Delta \rightarrow \{1, 2, 3\}$ such that

- (1). Any $a \in \Delta$ colored by i must be $a \in M_i$.
- (2). The coloring C satisfies the condition of Sperner's Lemma for each Δ .

We show the following can be done:

- For the point A_i , say $i = 1$, not $A_i = (1, 0) \in M_1$, so we can let $C(A_i) = i$.
- Consider a vertex $a = (x, y) \in A_1 A_2$, i.e. $x + y = 1$. Then $a \in M_1 \cap M_2$, otherwise $x + y < 1$ contradiction.

We have proved that for any $f : \Delta \rightarrow \Delta$ such $C : \Delta \rightarrow \{1, 2, 3\}$ exists. Apply Sperner's Lemma for the C on any Δ_i .

$\Rightarrow \exists$ a small triangle, say $A_1^{(i)} A_2^{(i)} A_3^{(i)}$ in Δ_i , which has 3 colors 1,2,3.

Consider the sequence $A_1^{(1)}, A_1^{(2)}, \dots, A_1^{(i)}, \dots$. Since everyone's bounded, there is a subsequence $A_1^{(i_1)}, A_1^{(i_2)}, \dots, A_1^{(i_j)}, \dots$ such that $\lim_{j \rightarrow \infty} A_1^{(i_j)} = p \in \Delta$

Since the diameter of $A_1^{(i)} A_2^{(i)} A_3^{(i)}$ is going to be 0 as $j \rightarrow \infty$, we see that $\lim_{j \rightarrow \infty} A_2^{(i_j)} = p$ and $\lim_{j \rightarrow \infty} A_3^{(i_j)} = p$

Note that $\beta_1(A_1^{(i)}) \geq \beta_1(f(A_1^{(i)}))$ so $\beta_1(p) \geq \beta_1(f(p))$. Similarly, $\beta_2(p) \geq \beta_2(f(p))$ and $\beta_3(p) \geq \beta_3(f(p)) \Rightarrow p \in M_1 \cap M_2 \cap M_3$. By Fact 2, p is a fixed point of f , i.e. $f(p) = p$. ■

Double Counting

Suppose that we can give two finite sets A and B , and a subset $S \subseteq A \times B$. And if $(a, b) \in S$, we say a and b are incident. Let $N_a = \#$ of elements $b \in B$, $N_b = \#$ of elements $a \in A$. Then $\sum_{a \in A} N_a = |S| = \sum_{b \in B} N_b$. Define a table $X = (x_{ij})$ where

$$x_{ij} = \begin{cases} 1 & i|j \\ 0 & \text{Otherwise} \end{cases}$$

Let $T(j) = \#$ of divisions of $j = \#$ of i 's in j^{th} column.

Let $\overline{T(n)} = \frac{1}{n} \sum_{j=1}^n T(j)$.

Fact: $|\overline{T(n)} - H(n)| < 1$, where $H(n) = \sum_{i=1}^n \frac{1}{i}$ is the n^{th} Harmonic number.

Proof. $\sum_{j=1}^n T(j) = \# \text{ 1's in } n \times n \text{ table} = \sum_{i=1}^n \lfloor \frac{n}{i} \rfloor$

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Sperner Theorem

Def: Let $\mathcal{F} \subseteq 2^{[n]}$ be a family of subsets of $[n]$. We say \mathcal{F} is independent (or \mathcal{F} is an independent system). If $\forall A, B \in \mathcal{F}, A \subsetneq B$ and $B \subsetneq A$, in other words, \mathcal{F} is independent \Leftrightarrow there is no "containment" relationship between any 2 subsets of \mathcal{F} .

Fact: For a fixed $k \in [n], \binom{[n]}{k}$ is an independent system.

Theorem 3 (Sperner's Theorem). *For any independent system \mathcal{F} of $[n]$, we have*

$$|\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}.$$

First proof: (Double-Counting).

Def: (1). A chain of subsets of $[n]$ is a sequence of distinct subsets $A_1 \subseteq A_2 \subseteq A_3 \subseteq \dots \subseteq A_k$

(2). A maximal chain is a chain with the property that No other subsets of $[n]$ can be inserted into it.

Fact 1: Any maximal chain looks like:

$$\phi \subseteq \{x_1\} \subseteq \{x_1, x_2\} \subseteq \dots \subseteq \{x_1, \dots, x_k\} \subseteq \dots \subseteq \{x_1, \dots, x_n\}.$$

Fact 2: There are exactly $n!$ maximal chains.

Why? Each maximal chain, say $\mathcal{C} : \emptyset \subseteq \{x_1\} \subseteq \{x_1, x_2\} \subseteq \dots \subseteq \{x_1, x_2, \dots, x_n\}$ defines a unique permutation:

$$\pi : [n] \rightarrow [n], \pi(i) = x_i$$

We first notice that this bound $|\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$. Next, we double count by considering the number of pairs (\mathcal{C}, A) such that:

- (1). \mathcal{C} is a maximal chain of $[n]$.
- (2). $A \in \mathcal{C} \cap \mathcal{F}$.

Recall the double counting.

$$\sum_{\mathcal{C}} N_{\mathcal{C}} = \#\text{pairs}(\mathcal{C}, A) = \sum_A N_A.$$

- $N_{\mathcal{C}} = \#\text{subsets } A \in \mathcal{C} \cap \mathcal{F} = |\mathcal{C} \cap \mathcal{F}| \leq 1$.
- $N_A = \#\text{maximal chains } \mathcal{C} \text{ s.t. } A \in \mathcal{C} = |A|!(n - |A|)!$

So we have

$$\begin{aligned} n! &= \sum_{\mathcal{C}} 1 \geq \sum_{\mathcal{C}} N_{\mathcal{C}} = \sum_{A \in \mathcal{F}} N_A \\ &= \sum_{A \in \mathcal{F}} |A|!(n - |A|)! = \sum_{A \in \mathcal{F}} \frac{n!}{\binom{n}{|A|}} \\ &\geq \sum_{A \in \mathcal{F}} \frac{n!}{\binom{n}{\lfloor \frac{n}{2} \rfloor}} = \frac{n!}{\binom{n}{\lfloor \frac{n}{2} \rfloor}} |\mathcal{F}| \\ &\Rightarrow |\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor} \end{aligned}$$

Second proof:

Def:A chain is symmetric if it consists of subsets of sizes $k, k+1, \dots, \lfloor \frac{n}{2} \rfloor, \dots, n-k-1, n-k$ for some k .

For example, $n=3$ $\{\{2\}, \{2, 3\}, [3]\}$ NOT Symmetric. $\{\phi, [3]\}$ NOT Symmetric.

Def:A partition of $2^{[n]}$ into symmetric chains is a way of expressing $2^{[n]}$ as a disjoint union of symmetric chains.

Theorem 4. *The family $2^{[n]}$ has a partition into symmetric chains.*

Proof of Sperner's Thm(Assuming Thm 2)

Note that any symmetric chain contains exactly one subset of size $\lfloor \frac{n}{2} \rfloor$. Since there are $\binom{n}{\lfloor \frac{n}{2} \rfloor}$ many subsets of size $\lfloor \frac{n}{2} \rfloor$, we see that any partition of $2^{[n]}$ into symmetric chains has to consist of exactly $\binom{n}{\lfloor \frac{n}{2} \rfloor}$ symmetric chains.

For each $A \in 2^{[n]}$, we define a sequence " $a_1 a_2 \dots a_n$ " consisting of left and right parentheses by defining

$$a_i = \begin{cases} "(" , & \text{if } i \in A \\ ")" , & \text{otherwise} \end{cases}$$

e.g. $n=7, A = \{2, 5, 6\}$, the sequence is $)()((()))((\rightarrow))()$.

We then define the "partial pairing of parentheses" as following:

- (1). First, we pair up all pairs " $()$ " of adjoint parentheses.
- (2). Then, we delete these already paired parentheses.
- (3). Repeat the above process until nothing can be done.

Note that when this process stops, the remaining unpaired parentheses must look like this:

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We say two subsets $A, B \in 2^{[n]}$ have the same partial pairing, if the paired parentheses are the same (even in the same positions).

$$n=11: A_1 = \{5, 6, 8\} \text{)})(()())$$

$$A_2 = \{5, 6, 8, 11\} \text{)})(()()()$$

$$A_3 = \{4, 5, 6, 8, 11\} \text{)})((()()()$$

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$$A_6 = \{1, 2, 3, 4, 5, 6, 8, 11\} ((((((()()()$$

$\Rightarrow \{A_1, A_2, \dots, A_6\}$ is a symmetric chain.

we can define an equivalence " \sim " on $2^{[n]}$ by letting $A \sim B$ iff A, B have the same partial pairing.

Exercise: Each equivalence class indeed forms a symmetric chain. This proves Thm2.

Littlewood-Offord Problem Fix a vector $\mathbf{a} = (a_1, a_2, \dots, a_n)$ with each $|a_i| \geq 1$. Let $S = \{\boldsymbol{\epsilon} = (\epsilon_1, \epsilon_2, \dots, \epsilon_n) : \epsilon_i \in \{1, -1\}, \boldsymbol{\epsilon} \cdot \mathbf{a} \in (-1, 1)\}$, Then $|S| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$.

Proof. For $\forall \boldsymbol{\epsilon} \in S$, define $A_{\boldsymbol{\epsilon}} = \{i \in [n] : a_i \epsilon_i > 0\}$. Let $\mathcal{F} = \{A_{\boldsymbol{\epsilon}} : \boldsymbol{\epsilon} \in S\}$.

$$\Rightarrow |S| = |\mathcal{F}|$$

We want \mathcal{F} is independent system.

Suppose NOT, say $A_{\epsilon_1}, A_{\epsilon_2} \in \mathcal{F}, A_{\epsilon_1} \subseteq A_{\epsilon_2}$

$$\begin{cases} \epsilon_1 \cdot \mathbf{a} \in (-1, 1) \\ \epsilon_2 \cdot \mathbf{a} \in (-1, 1) \end{cases}$$

$$\epsilon_1 \cdot \mathbf{a} = \sum_{i \in A_{\epsilon_1}} |a_i| - \sum_{i \notin A_{\epsilon_1}} |a_i| = 2 \sum_{i \in A_{\epsilon_1}} |a_i| - \sum_{i=1}^n |a_i|$$

$$\epsilon_2 \cdot \mathbf{a} - \epsilon_1 \cdot \mathbf{a} = 2 \left(\sum_{i \in A_{\epsilon_2}} |a_i| - \sum_{j \in A_{\epsilon_1}} |a_j| \right) \geq 2|a_j| \geq 2 \text{ for some } j \in A_{\epsilon_2} \setminus A_{\epsilon_1}$$

But this is a contradiction as $|\epsilon_2 \cdot \mathbf{a} - \epsilon_1 \cdot \mathbf{a}| < 2$.

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