

Combinatorics

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1 Double Counting

The basic setting of the double counting technique is as follows. Suppose that we are given two finite sets A and B , and a subset $S \subseteq A \times B$. If $(a, b) \in S$, then we say that a and b are incident. Let N_a be the number of elements $b \in B$ such that $(a, b) \in S$, and N_b be the number of elements $a \in A$ such that $(a, b) \in S$. Then we have

$$\sum_{a \in A} N_a = |S| = \sum_{b \in B} N_b.$$

Theorem 1.1. *Let $T(j)$ be the number of divisors of a positive integer j . Let $\overline{T(n)} = \frac{1}{n} \sum_{j=1}^n T(j)$. Then we have $|\overline{T(n)} - H(n)| < 1$, where $H(n) = \sum_{i=1}^n \frac{1}{i}$ is the n^{th} Harmonic number.*

Proof. Define a table $X = (x_{ij})$ where

$$x_{ij} = \begin{cases} 1 & \text{if } i|j \\ 0 & \text{otherwise.} \end{cases}$$

Then

$$\sum_{j=1}^n T(j) = \sum_{1 \leq i \leq j \leq n} x_{ij} = \sum_{i=1}^n \lfloor \frac{n}{i} \rfloor,$$

which implies that

$$\overline{T(n)} = \frac{1}{n} \sum_{i=1}^n \lfloor \frac{n}{i} \rfloor.$$

Then we have

$$|\overline{T(n)} - H(n)| < 1. \quad \blacksquare$$

Exercise 1.2. *Prove that*

$$\left| \frac{1}{n} \sum_{i=1}^n \lfloor \frac{n}{i} \rfloor - \sum_{i=1}^n \frac{1}{i} \right| < 1.$$

1.1 Sperner's Theorem

Definition 1.3. *Let $\mathcal{F} \subseteq 2^{[n]}$ be a family of subsets of $[n]$. We say \mathcal{F} is independent (or \mathcal{F} is an independent system), if for any two $A, B \in \mathcal{F}$, we have $A \not\subseteq B$ and $B \not\subseteq A$. In other words, \mathcal{F} is independent if and only if there is no "containment" relationship between any two subsets of \mathcal{F} .*

Fact 1.4. *For a fixed $k \in [n]$, $\binom{[n]}{k}$ is an independent system.*

Theorem 1.5 (Sperner's Theorem). *For any independent system \mathcal{F} of $[n]$, we have*

$$|\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}.$$

First proof of Sperner's Theorem (Double-Counting). A *chain* of subsets of $[n]$ is a sequence of distinct subsets

$$A_1 \subseteq A_2 \subseteq A_3 \subseteq \dots \subseteq A_k.$$

A *maximal chain* is a chain with the property that no other subsets of $[n]$ can be inserted into it to find a longer chain. We have the following observations.

(1). Any maximal chain looks like:

$$\phi \subseteq \{x_1\} \subseteq \{x_1, x_2\} \subseteq \dots \subseteq \{x_1, \dots, x_k\} \subseteq \dots \subseteq \{x_1, \dots, x_n\}.$$

(2). There are exactly $n!$ maximal chains.

This is because any such a maximal chain, say $\mathcal{C} : \phi \subseteq \{x_1\} \subseteq \{x_1, x_2\} \subseteq \dots \subseteq \{x_1, x_2, \dots, x_n\}$, defines a unique permutation:

$$\pi : [n] \rightarrow [n], \pi(i) = x_i, \forall i \in [n].$$

Now we double count the number of pairs (\mathcal{C}, A) satisfying that:

- \mathcal{C} is a maximal chain of $[n]$.
- $A \in \mathcal{C} \cap \mathcal{F}$.

Recall the rule of double counting given at the beginning that

$$\sum_{\mathcal{C}} N_{\mathcal{C}} = \text{the number of pairs } (\mathcal{C}, A) = \sum_A N_A,$$

where $N_{\mathcal{C}}$ is the number of subsets $A \in \mathcal{C} \cap \mathcal{F}$ and N_A is the number of maximal chains \mathcal{C} contains A . It is key to observe that

- $N_{\mathcal{C}} \leq 1$,
- $N_A = |A|!(n - |A|)!$

So we have

$$\begin{aligned} n! &= \sum_{\mathcal{C}} 1 \geq \sum_{\mathcal{C}} N_{\mathcal{C}} = \sum_{A \in \mathcal{F}} N_A = \sum_{A \in \mathcal{F}} |A|!(n - |A|)! \\ &= \sum_{A \in \mathcal{F}} \frac{n!}{\binom{n}{|A|}} \geq \sum_{A \in \mathcal{F}} \frac{n!}{\binom{n}{\lfloor \frac{n}{2} \rfloor}} = \frac{n!}{\binom{n}{\lfloor \frac{n}{2} \rfloor}} |\mathcal{F}|, \end{aligned}$$

which implies that

$$|\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}.$$

This finishes the proof. ■

Now we give another proof of Sperner's Theorem.

Definition 1.6. A chain is symmetric if it consists of subsets of sizes $k, k+1, \dots, \lfloor \frac{n}{2} \rfloor, \dots, n-k-1, n-k$ for some $k \geq 0$.

For example, when $n = 3$, $\{\{2\}, \{2, 3\}, \{1, 2, 3\}\}$ is not symmetric. And when $n = 4$, $\{\emptyset, \{1, 2, 3\}\}$ is not symmetric.

Theorem 1.7. The family $2^{[n]}$ can be partitioned into a disjoint union of symmetric chains.

Proof of Theorem 1.7. For each $A \in 2^{[n]}$, we define a sequence " $a_1 a_2 \dots a_n$ " consisting of left and right parentheses by defining

$$a_i = \begin{cases} "(", & \text{if } i \in A \\ ")", & \text{otherwise} \end{cases}$$

We then define the "partial pairing of parentheses" as following:

- (1). First, we pair up all pairs "(" of adjoint parentheses.
- (2). Then, we delete these already paired parentheses.
- (3). Repeat the above process until nothing can be done.

Note that when this process stops, the remaining unpaired parentheses must look like this:

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We say two subsets $A, B \in 2^{[n]}$ have the same partial pairing, if the paired parentheses are the same (even in the same positions).

We can define an equivalence " \sim " on $2^{[n]}$ by letting $A \sim B$ if and only if A, B have the same partial pairing.

Exercise 1.8. Each equivalence class indeed forms a symmetric chain.

Using this fact, now we see that $2^{[n]}$ can be partitioned into disjoint equivalence classes, which are disjoint symmetric chains. This fishes the proof. ■

Theorem 1.7 can rapidly imply Sperner's Theorem.

Second proof of Sperner's Theorem. Note that by definition, any symmetric chain contains exactly one subset of size $\lfloor \frac{n}{2} \rfloor$. Since there are $\binom{n}{\lfloor \frac{n}{2} \rfloor}$ many subsets of size $\lfloor \frac{n}{2} \rfloor$, by Theorem 1.7, we see that any partition of $2^{[n]}$ into symmetric chains has to consist of exactly $\binom{n}{\lfloor \frac{n}{2} \rfloor}$ symmetric chains. Each symmetric chain can contain at most one subset from \mathcal{F} and thus we see $|\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$. ■

1.2 Littlewood-Offord Problem

Theorem 1.9. Fix a vector $\vec{a} = (a_1, a_2, \dots, a_n)$ with each $|a_i| \geq 1$. Let $S = \{\vec{\epsilon} = (\epsilon_1, \epsilon_2, \dots, \epsilon_n) : \epsilon_i \in \{1, -1\} \text{ and } \vec{\epsilon} \cdot \vec{a} \in (-1, 1)\}$, then $|S| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$.

Remark: Note that this is tight for many vectors \vec{a} .

Proof. For any $\vec{\epsilon} \in S$, define $A_{\vec{\epsilon}} = \{i \in [n] : a_i \epsilon_i > 0\}$. Let $\mathcal{F} = \{A_{\vec{\epsilon}} : \vec{\epsilon} \in S\}$. Then we have

$$|S| = |\mathcal{F}|.$$

Now we claim that \mathcal{F} is an independent system. Suppose for a contradiction that there exist $A_{\vec{\epsilon}_1}, A_{\vec{\epsilon}_2} \in \mathcal{F}$ with $A_{\vec{\epsilon}_1} \subseteq A_{\vec{\epsilon}_2}$. That also says,

$$\begin{cases} \vec{\epsilon}_1 \cdot \vec{a} \in (-1, 1) \\ \vec{\epsilon}_2 \cdot \vec{a} \in (-1, 1) \end{cases}$$

which imply that

$$|\epsilon_1 \cdot \vec{a} - \epsilon_2 \cdot \vec{a}| < 2.$$

By definition, we have

$$\vec{\epsilon}_1 \cdot \vec{a} = \sum_{i \in A_{\vec{\epsilon}_1}} |a_i| - \sum_{i \notin A_{\vec{\epsilon}_1}} |a_i| = 2 \sum_{i \in A_{\vec{\epsilon}_1}} |a_i| - \sum_{i=1}^n |a_i|.$$

Since $A_{\vec{\epsilon}_1} \subseteq A_{\vec{\epsilon}_2}$, we also have that

$$\vec{\epsilon}_2 \cdot \vec{a} - \vec{\epsilon}_1 \cdot \vec{a} = 2 \left(\sum_{i \in A_{\vec{\epsilon}_2}} |a_i| - \sum_{j \in A_{\vec{\epsilon}_1}} |a_j| \right) \geq 2|a_j| \geq 2, \text{ for some } j \in A_{\vec{\epsilon}_2} \setminus A_{\vec{\epsilon}_1},$$

a contradiction. By Sperner's Theorem, we have $|S| = |\mathcal{F}| \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$. This finishes the proof ■

1.3 Turán Type Problem

Definition 1.10. A graph G is bipartite if its vertex set can be partitioned into two parts (say A and B) such that each edge joints one vertex in A and another in B .

This is equivalent to say that $V(G)$ can be partitioned into two independent subsets. And we say (A, B) is a bipartition of G . For example, all even cycles C_{2k} are bipartite, while all odd cycles C_{2k+1} are not.

Definition 1.11. Let $K_{a,b}$ be the complete bipartite graph with two parts of sizes a and b . This is a bipartite graph with edge set $\{(i, j) : i \in A, j \in B\}$ where $|A| = a$ and $|B| = b$.

Definition 1.12. Given a graph H , we say a graph G is H -free if G does not contain a copy of H as its subgraph.

For example, $K_{a,b}$ is K_3 -free.

Definition 1.13. For fixed graph H , let the Turán number of H , denoted by $\text{ex}(n, H)$, be the maximum number of edges in an n -vertex H -free graph G .

Theorem 1.14. $\text{ex}(n, C_4) \leq \frac{n}{4}(1 + \sqrt{4n - 3})$.

Proof. Let G be a C_4 -free graph with n vertices. We need to show that $e(G) \leq \frac{n}{4}(1 + \sqrt{4n - 3})$. Consider $S = \{(\{u_1, u_2\}, w) : u_1 w u_2 \text{ is a path of length 2 in } G\}$. Since G is C_4 -free, for fixed $\{u_1, u_2\}$, there is at most one vertex w such that $(\{u_1, u_2\}, w) \in S$. So we have

$$|S| = \sum_{\{u_1, u_2\}} \text{the number of } (\{u_1, u_2\}, w) \in S \leq \sum_{\{u_1, u_2\}} 1 = \binom{n}{2}.$$

On the other hand, fixed a vertex w , the number of $\{u_1, u_2\}$ such that $(\{u_1, u_2\}, w) \in S$ exactly equals $\binom{d(w)}{2}$, which implies that

$$|S| = \sum_{w \in V(G)} \binom{d(w)}{2} = \frac{1}{2} \sum_{w \in V(G)} d^2(w) - e(G).$$

Putting the above together, we have

$$\binom{n}{2} \geq |S| = \frac{1}{2} \sum_{w \in V(G)} d^2(w) - e(G).$$

Using Cauchy-Schwarz inequality, we have

$$\frac{n^2 - n}{2} \geq \frac{n}{2} \sum_{w \in V(G)} \frac{d^2(w)}{n} - e(G) \geq \frac{n}{2} \sum_{w \in V(G)} \left(\frac{d(w)}{n}\right)^2 - e(G),$$

which implies that

$$\frac{2e^2(G)}{n} - e(G) \leq \frac{n^2 - n}{2}.$$

Solving it, we can derive easily that $e(G) \leq \frac{n}{4}(1 + \sqrt{4n - 3})$. ■

Exercise 1.15. Prove that $\text{ex}(n, C_4) < \frac{n}{4}(1 + \sqrt{4n - 3})$.

Corollary 1.16. We have $\text{ex}(n, C_4) \leq (\frac{1}{2} + o(n))n^{\frac{3}{2}}$, where $o(n) \rightarrow 0$ as $n \rightarrow \infty$.

Theorem 1.17 (Mantel's Thm). $\text{ex}(n, K_3) = \lfloor \frac{n^2}{4} \rfloor$.

Proof. We first consider the lower bound $\text{ex}(n, K_3) \geq \lfloor \frac{n^2}{4} \rfloor$ as the complete bipartite graph $K_{\lfloor \frac{n}{2} \rfloor, \lfloor \frac{n}{2} \rfloor}$ is K_3 -free and has $\lfloor \frac{n^2}{4} \rfloor$ edges.

Next, we show $\text{ex}(n, K_3) \leq \lfloor \frac{n^2}{4} \rfloor$. We prove by induction on n that any n -vertex K_3 -free graph G has at most $\frac{n^2}{4}$ edges. First it holds trivially when $n \in \{1, 2\}$. Now we assume that any K_3 -free graph H with less than n vertices has at most $|V(H)|^2/4$ edges. Let G be K_3 -free with n vertices. Take any edge of G , say $xy \in E(G)$. Since G is K_3 -free, we say $N_G(x) \cap N_G(y) = \emptyset$, implies that $|d(x)| + |d(y)| \leq n$.

Let H be a graph obtained from G by deleting vertex x and y . Note that H is also K_3 -free and has $n - 2$ vertices. By induction, $e(H) \leq \frac{(n-2)^2}{4}$. Thus we have that

$$e(G) = e(H) + |d(x)| + |d(y)| - 1 \leq \frac{(n-2)^2}{4} + n - 1 = \frac{n^2}{4}.$$

This finishes the proof. ■

Exercise 1.18. *The unique n -vertex K_3 -graph which attains the maximum number of edges $\text{ex}(n, K_3)$ is the complete bipartite graph $K_{\lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil}$.*