

Combinatorial Networks
Week 6, April 22-23

Even-cycles-free

- For bipartite H (so $\chi(H) = 2$), it is unknown the order of $ex(n, H)$.
- For any H with $\chi(H) = k + 1 \geq 3$, $ex(n, H) = (1 - \frac{1}{k} + o_H(1)) \frac{n^2}{2}$.
- For special bipartite graphs, the even cycles C_{2t} , it is known that $ex(n, C_{2t}) \leq O(n^{1+\frac{1}{t}})$
- **Conjecture.** $\forall t \geq 2$, $ex(n, C_{2t}) = \Theta(n^{1+\frac{1}{t}})$
- **Definition.** Let $\mathcal{F} = \{\text{some graphs } H\}$, $ex(n, \mathcal{F}) =$ the maximum of $e(G)$, G has n vertices, $\forall H \in \mathcal{F}$, G is H -free.
- **Theorem.** $\forall t \geq 2$, $ex(n, \{C_3, C_4, \dots, C_{2t}\}) \leq n^{1+\frac{1}{t}} + n$
- **Lemma.** Every graph with at least dn edges has a subgraph whose minimum degree is at least d .
- **Proof of Lemma.** Delete vertices whose degrees are less than d . The subgraph left is not empty because the number of deleted edges is less than dn . ■
- **Proof of Theorem.** Suppose there is a $\{C_3, C_4, \dots, C_{2t}\}$ -free n -vertices graph G with more than $n^{1+\frac{1}{t}} + n$ edges. By lemma, G has a subgraph G' with minimum degree $\delta(G') \geq n^{\frac{1}{t}} + 1$. Consider G' and the BFS-tree(Breadth-First Search).
Fix $v \in V(G')$, define $L_0 = v$, $L_1 = N(v)$
$$L_{i+1} = \{w \in V(G') - \bigcup_{j=0}^i L_j : \exists u \in L_i, uw \in E(G')\} = N(L_i) - \bigcup_{j=0}^i L_j$$
And G' is $\{C_3, C_4, \dots, C_{2t}\}$ -free, so $|L_i| \geq n^{\frac{1}{t}} |L_{i-1}|$, $|L_i| \geq n^{\frac{i}{t}}$, $|L_t| \geq n$, a contradiction.
This shows that $ex(n, \{C_3, C_4, \dots, C_{2t}\}) \leq n^{1+\frac{1}{t}} + n$. ■

$K_{t,t}$ -free

- Recall: $ex(n, K_{s,t}) \leq Cn^{2-\frac{1}{t}}$ for $s \geq t \geq 2$
- **Conjecture.** $\forall t \geq 2$, $ex(n, K_{t,t}) = \Theta(n^{2-\frac{1}{t}})$
open for $t \geq 4$, need a lower bound.
- **Theorem(Erdős).** $ex(n, K_{2,2}) = \Theta(n^{\frac{3}{2}})$
- **Proof.** We construct an n -vertex graph with NO $K_{2,2}$ as follows.
Let $p \geq 2$ be a prime, $V(G) = \{(x, y) : x, y \in \mathbb{Z}_p, x \neq y\}$, $n = p(p-1)$. For (x, y) , $(a, b) \in V(G)$, define $(x, y) \overset{G}{\sim} (a, b)$ iff $xa + yb \equiv 1 \pmod{p}$. And $\forall v \in V(G)$, $d(v) = p$, we call G is p -regular.
Next we verify G is $K_{2,2}$ -free. Suppose NOT, say $\exists(a_1, b_1), (a_2, b_2) \in V(G)$, s.t.

$$\begin{cases} a_1x + b_1y \equiv 1 \pmod{p} \\ a_2x + b_2y \equiv 1 \pmod{p} \end{cases}$$

has 2 solutions, a contradiction. ■

- **Theorem(Brown).** $ex(n, K_{3,3}) = \Theta(n^{\frac{5}{3}})$
- **Proof.** Let $p \geq 2$ be a prime, define G by $V(G) = \{(x, y, z) : x, y, z \in \mathbb{Z}_p, x \neq y \neq z\}$, $n = p(p-1)(p-2)$. For $(x, y, z), (a, b, c) \in V(G)$, $(x, y, z) \stackrel{G}{\sim} (a, b, c)$ iff $(x-a)^2 + (y-b)^2 + (z-c)^2 \equiv 1 \pmod{p}$. G is p -regular with $d = \Theta(p^2)$, $e(G) = \frac{1}{2}nd = \Theta(n^{\frac{5}{3}})$. Verify G has NO $K_{3,3}$. Suppose NOT, say $\exists(a_1, b_1, c_1), (a_2, b_2, c_2) \in V(G)$, s.t.

$$\begin{cases} (x - a_1)^2 + (y - b_1)^2 + (z - c_1)^2 \equiv 1 \pmod{p} \\ (x - a_2)^2 + (y - b_2)^2 + (z - c_2)^2 \equiv 1 \pmod{p} \\ (x - a_3)^2 + (y - b_3)^2 + (z - c_3)^2 \equiv 1 \pmod{p} \end{cases}$$

has 3 solutions, a contradiction.(Consider in \mathbb{R}^3 , since 3 spheres in general position intersect in at most 2 points.) ■

1-Distance Problems

- Recall: $ex(n, K_{s,t}) = O(n^{1-\frac{1}{t}})$.
In particular, $ex(n, K_{2,3}) = O(n^{\frac{3}{2}})$
- **Problem 1.** In plane, given n vertices in general position, how many pairs of points have distance 1?
- Define a graph G on these n vertices, $a \stackrel{G}{\sim} b$ iff the distance between a and b is 1.
Fact. G is $K_{2,3}$ -free. Because two circle at a and b with radius 1 intersect in at most 2 points.
- **Theorem 1.** There are at most $O(n^{\frac{3}{2}})$ pairs at distance 1.
- **Proof.** Define G as above, so the number of pairs at distance 1 equals to $e(G)$, where G is $K_{2,3}$ -free. ■
- **Erdős.** Construct n points in plane with $n^{1+\frac{1}{\log \log n}}$ pairs at distance 1.
- **Open Problem.** Find an example of n points with n^{1+c} pairs at distance 1 for some $c > 0$.
- **Problem 2.** How many points in \mathbb{R}^n s.t. the distance between any 2 points is 1.
- **Theorem 2.** In \mathbb{R}^n , there are at most $n + 1$ points such that the distance between any 2 points is 1.
- **Proof.** Consider \mathbb{R}^n , say $\exists m$ points with distance 1 to each other. We may assume that one of the m points is $(0, 0, \dots, 0)$, then the other $m - 1$ points can be viewed as vectors, say $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_{m-1}$, and

$$\forall i, \langle \mathbf{v}_i, \mathbf{v}_i \rangle = 1; \forall i \neq j, \langle \mathbf{v}_i, \mathbf{v}_j \rangle = \frac{1}{2}$$

Let

$$A = \begin{pmatrix} \mathbf{v}_1 \\ \mathbf{v}_2 \\ \vdots \\ \mathbf{v}_{m-1} \end{pmatrix}_{(m-1) \times n} \quad AA^T = \begin{pmatrix} 1 & 1/2 & \cdots & 1/2 \\ 1/2 & 1 & \cdots & 1/2 \\ \vdots & \vdots & \ddots & \vdots \\ 1/2 & 1/2 & \cdots & 1 \end{pmatrix}_{(m-1) \times (m-1)}$$

Since $\det(AA^T) \neq 0$, $m - 1 = \text{rank}(AA^T) \leq \text{rank}(A) \leq n$, so $m \leq n + 1$. ■

2-Distance Problem

- How many points can we have in \mathbb{R}^n , such that their distance are one of 2 choices.
- **Example 1.** We can find $\binom{n}{2}$ such points.
Consider 0-1 vectors whose 1-norm are exactly 2.
- **Example 2.** We can find $\binom{n+1}{2}$ such points in a n -dimension space.
In \mathbb{R}^{n+1} , we can find $\binom{n+1}{2}$ such points according to Example 1. And those points are in the plane $\{\mathbf{x} = (x_1, x_2, \dots, x_{n+1}) : \sum_{i=1}^{n+1} x_i = 2\}$, which is a n -dimension space.
- **Theorem.** There are at most $\binom{n+2}{2} + n + 1$ points in \mathbb{R}^n such that there distances are of 2 choices.
- **Proof.** Let $a^{(1)}, a^{(2)}, \dots, a^{(m)} \in \mathbb{R}^n$ form a 2-distance set with distances d_1, d_2 , write $a^{(i)} = (a_1^{(i)}, a_2^{(i)}, \dots, a_n^{(i)})$. Define $f_i : \mathbb{R}^n \rightarrow \mathbb{R}$,

$$f_i(x) = (\|x - a^{(i)}\|^2 - d_1^2)(\|x - a^{(i)}\|^2 - d_2^2)$$

we have $f_i(a^{(i)}) = d_1^2 d_2^2$ for $\forall i$, and $f_i(a^{(j)}) = 0$ for $\forall i \neq j$.

Let $\lambda_1, \lambda_2, \dots, \lambda_m \in \mathbb{R}$, and $\sum_{i=1}^m \lambda_i f_i = 0$, that is $\sum_{i=1}^m \lambda_i f_i(x) \equiv 0$ for $\forall x \in \mathbb{R}^n$.

$$\forall j, 0 = \sum_{i=1}^m \lambda_i f_i(a^{(j)}) = \lambda_j d_1^2 d_2^2 \Rightarrow \lambda_j = 0$$

So f_1, f_2, \dots, f_m are linearly independent. They are all polynomials, and in the polynomial space which is spanned by

$$\left(\sum_{k=1}^n x_k^2\right)^2, \left(\sum_{k=1}^n x_k^2\right)x_i, x_i x_j, x_i, 1$$

whose dimension is $\binom{n+2}{2} + n + 1$. So $m = \dim \text{span}\{f_1, f_2, \dots, f_m\} \leq$ the dimension of the space above $= \binom{n+2}{2} + n + 1$. ■

Odd-Town Problem

- **Problem.** Suppose we have a town with n people, they are assigned to form m clubs A_1, A_2, \dots, A_m such that
 - (1) $|A_i|$ is odd.
 - (2) $|A_i \cap A_j|$ is even for $\forall i \neq j$.
 How many clubs can we have?
- **Examples.** Here are some examples where $m = n$.
 - (1) $A_i = \{i\}$
 - (2) $A_i = [n] - \{i\}$, n is even.
 - (3) $A_1 = [n] - \{1\}, A_2 = [n] - \{2\}$ and $A_i = \{1, 2, i\}$ for $i > 2$, n is even.

- **Theorem.** In the odd-town problem, $m \leq n$.
- **Proof.** For A_i , define 0-1 vector \mathbf{v}_i , $\mathbf{v}_i(j) = 1$ iff $j \in A_i$. We have
 $\langle \mathbf{v}_i, \mathbf{v}_i \rangle = |A_i|$ is odd.
 $\langle \mathbf{v}_i, \mathbf{v}_j \rangle = |A_i \cap A_j|$ is even for $\forall i \neq j$.
 Claim: $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_m$ are linearly independent over \mathbb{F}_2 .
 Proof: Exercise.
 Since $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_m \in \mathbb{F}_2^n$ are linearly independent, $m \leq n$.
 Or we can let

$$A = \begin{pmatrix} \mathbf{v}_1 \\ \mathbf{v}_2 \\ \vdots \\ \mathbf{v}_m \end{pmatrix}_{m \times n} \quad AA^T = \begin{pmatrix} 1 & 0 & \cdots & 0 \\ 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 1 \end{pmatrix}_{m \times m} \quad \text{over } \mathbb{F}_2$$

so $m = \text{rank}(AA^T) \leq \text{rank} A \leq n$. ■

Even-Town Problem

- **Problem.** Suppose we have distinct $A_1, A_2, \dots, A_m \subseteq [n]$, satisfying
 (1) $|A_i|$ is even.
 (2) $|A_i \cap A_j|$ is even for $\forall i \neq j$.
 How much can m be?
- **Example.** Pair the n elements into $n/2$ pairs, say $S = \{(1, 2), (3, 4), \dots, (n-1, n)\} = \{P_1, P_2, \dots, P_{n/2}\}$, there are $2^{n/2}$ subsets of S . That is, $m = 2^{n/2}$.
- **Theorem.** In the even-town problem, $m \leq 2^{n/2}$.
- **Proof.** Define \mathbf{v}_i as before, and let

$$A = \begin{pmatrix} \mathbf{v}_1 \\ \mathbf{v}_2 \\ \vdots \\ \mathbf{v}_m \end{pmatrix}_{m \times n}$$

$\text{rank} A = \dim \text{span}\{\mathbf{v}_i\}$, and $\text{rank} A + \dim \text{Ker} A = n$ (number of columns of A).

Case1: $\text{rank} A \leq \frac{n}{2}$

This shows that $\dim \text{span}\{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_m\} \leq \frac{n}{2}$, so $\text{span}\{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_m\} \subseteq \mathbb{F}_2^{n/2}$, $m \leq 2^{n/2}$.

Case2: $\text{rank} A \geq \frac{n}{2} + 1$

This implies that we have at least $\frac{n}{2} + 1$ linearly independent vectors, say $\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_{\frac{n}{2}+1}$, consider

$$A\mathbf{u}_i^T = \begin{pmatrix} \langle \mathbf{v}_1, \mathbf{u}_i \rangle \\ \langle \mathbf{v}_2, \mathbf{u}_i \rangle \\ \vdots \\ \langle \mathbf{v}_m, \mathbf{u}_i \rangle \end{pmatrix} = \begin{pmatrix} \text{even} \\ \text{even} \\ \vdots \\ \text{even} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix} \text{ over } \mathbb{F}_2$$

so $\{\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_{\frac{n}{2}+1}\} \subseteq \text{Ker} A$, $\dim \text{Ker} A \geq \frac{n}{2} + 1$. But it contradicts to $\text{rank} A + \dim \text{Ker} A = n$. ■