

Extremal and Probabilistic Graph Theory

Lecture 1

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First we review some Min-Max Theorems.

Theorem 1.1 (König Theorem). *For a bipartite graph G , the minimum number of a vertex-cover in G = the maximum size of a matching in G*

Definition 1.2. A vertex-cover A is a set of $V(G)$, s.t. $V(G) \setminus A$ is a stable set.

Theorem 1.3 (Menger's Theorem). *For any graph G and any subset $A, B \subset V(G)$, the minimum size of a cut S separating A from B in G = the maximum number of vertex-disjoint paths from A to B in G*

We will investigate some Min-Max relations on packings and coverings in digraphs.

Definition 1.4. A (proper) coloring of digraph D is just a (proper) coloring of its underlying graph G .

Definition 1.5. The chromatic number $\chi(D)$ of digraph D is the chromatic number $\chi(G)$ of graph G .

Definition 1.6. A path or cycle in digraph denotes a directed path or cycle.

Theorem 1.7 (Gallai-Roy Theorem). *Every digraph D has a path with $\chi(D)$ vertices.*

Proof. We say a path in digraph is a u -path if the first vertex of it is u .

Let k be the number of vertices in a longest path in D , so it will sufficient to show that there exists a k -coloring of D .

Let D' be a maximum acyclic subdigraph of D , and for any $v \in V(D)$, define $c(v)$ to be the number of vertices in a longest v -path in D' , then it is easy to see: $1 \leq c(v) \leq k$.

Then we show that this function $c : V(D) \rightarrow [k]$ is a proper coloring of D .

Consider any arc $(u, v) \in A(D)$, we claim that $c(u) \neq c(v)$.

Case 1: $(u, v) \in A(D)$

Consider the longest v -path P in D' , then $u \notin P$, otherwise we get a cycle! So we have $c(u) \geq c(v) + 1$, done!

Case 2: $(u, v) \notin A(D)$

Then $(u, v) \in D'$ contains a cycle, for D' is the maximum graph. So D' contains a path P from v to u . Let Q denotes the longest u -path in D' , for the similar reason, P and Q are disjoint, except at u . Then $P \cup Q$ is longer than Q , so we have $c(v) > c(u)$, done! ■

Exercises:

For digraph D , let $\lambda(D)$ = the number of vertices in a longest path in D . Then for graph G , we have $\chi(D) = \min\{\lambda(D) : D \text{ is an orientation of } G\}$.

Definition 1.8. A path partition \mathcal{P} of a digraph is a collection of vertex-disjoint paths, s.t. the union of the vertex set of those path is $V(D)$. A single vertex can be viewed as a trivial path.

Definition 1.9. Let $\pi(D) = \min|\mathcal{P}|$, where the minimum is taken for all path partitions, and let $\alpha(D) = \max|\text{stable set in } D|$.

Then we have:

Theorem 1.10 (Gallai-Milgram Theorem). *For any digraph D , $\pi(D) \leq \alpha(D)$.*

Proof. We prove a stronger statement:

We say a stable set S is orthogonal to a path partition \mathcal{P} , if every $P \in \mathcal{P}$, has exactly one common vertex in S .

We also need the following lemma:

Let \mathcal{P} be a path partition, suppose no stable set in D orthogonal to \mathcal{P} , then there exists a path partition \mathcal{Q} in D , s.t. $|\mathcal{Q}| = |\mathcal{P}| - 1$, $i(\mathcal{Q}) \subset i(\mathcal{P})$, & $t(\mathcal{Q}) \subset t(\mathcal{P})$. Here, $i(\mathcal{P})$ and $t(\mathcal{P})$ denotes the sets of first vertices and last vertices in the paths of \mathcal{P} respectively.

Proof of the lemma: By induction on the vertices of D .

Base case is trivial when $|V(D)| = 1$.

General cases: By the hypothesis, $t(\mathcal{P})$ is not stable. So there exists $y, z \in t(\mathcal{P})$, & $(y, z) \in A(D)$.

If $|V(P_z)| = 1$, then \mathcal{Q} is obtained from \mathcal{P} by deleting the path P_z and replace P_y by $P_y \cap (y, z)$, done! Here P_y denotes the path with the end points y .

If $|V(P_z)| \geq 2$, let x be the predecessor of z in P_z . Consider $D' = D \setminus \{z\}$, $P'_z = P_z \setminus \{z\}$, & $\mathcal{P}' = (\mathcal{P} \setminus \{P_z\}) \cup \{P'_z\}$.

We should note that $t(P') = (t(P) \setminus \{z\}) \cup \{x\}$, $i(\mathcal{P}') = i(\mathcal{P})$

And if there is a stable set S' in D' orthogonal to \mathcal{P}' , then clearly this is also a stable set in D and also orthogonal to \mathcal{P} , which leads to a contradiction!

So, there are no such stable set in D' . By induction on D' , we can find a new path partition \mathcal{Q}' in D' , with $|\mathcal{Q}'| = |\mathcal{P}'| - 1$, $i(\mathcal{Q}') \subset i(\mathcal{P}')$, & $t(\mathcal{Q}') \subset t(\mathcal{P}')$, then either x or $y \in t(\mathcal{Q}')$.

If $x \in t(\mathcal{Q}')$, add the arc (x, z) is ok. Otherwise, add the arc (y, z) is done!

In both cases, we have $|\mathcal{Q}| = |\mathcal{P}| - 1$, $i(\mathcal{Q}) \subset i(\mathcal{P})$, & $t(\mathcal{Q}) \subset t(\mathcal{P})$.

Now we turn to prove the main theorem:

Choose \mathcal{P} to be the minimum partition of D , then by our lemma before, there must be a stable set S orthogonal to \mathcal{P} , so $\pi(D) = |\mathcal{P}| \leq |S| \leq \alpha(D)$. ■

Summary:

Gallai-Roy Theorem says that any digraph D has a path with $\chi(D)$ vertices, or equivalently, Minimum number of disjoint stable sets whose union in $V(D) \leq$ maximum of the number of vertices in a path of D .

Gallai-Milgram Theorem says that for any digraph D we have $\pi(D) \leq \alpha(D)$, or equivalently, Minimum number of disjoint paths whose union in $V(D) \leq$ maximum of the number of vertices in a stable set of D .

Now we can see that Gallai-Roy Theorem and Gallai-Milgram Theorem have the following *dual* relation: one can be transformed from another by just interchanging the role of *paths* and *stable sets*.