

# Extremal and Probabilistic Graph Theory

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## 1 Lecture 10. Induced Removal Lemma

**Lemma 1.1** (Strong regularity lemma). *For any sequence of constants  $\varepsilon_0 \geq \varepsilon_1 \geq \dots > 0$ , there exists an integer  $M$  such that every graph has two equipartitions  $P, Q$  satisfying  $P$  is  $\varepsilon_0$ -regular,  $Q$  is  $\varepsilon_{|P|}$ -regular,  $Q$  refines  $P$ ,  $|Q| \leq M$  and  $q(Q) \leq q(P) + \varepsilon_0$ .*

*Proof.* The proof of regularity lemma can also show the following statement: for  $\varepsilon > 0$ , there is a constant  $M(\varepsilon)$  such that for any equipartition  $P$  of  $V(G)$ , we can get a refinement  $P'$  of  $P$  where each part of  $P$  refines to at most  $M(\varepsilon)$  parts and  $P'$  is  $\varepsilon$ -regular.

By repeatedly applying the above statement, we obtain a sequence of equipartitions  $P_0, P_1, \dots$  of  $V(G)$ , satisfying each  $P_{i+1}$  refines  $P_i$  and  $P_{i+1}$  is  $\varepsilon_{|P_i|}$ -regular with  $|P_{i+1}| \leq |P_i| M(\varepsilon_{|P_i|})$ . Since  $q(P_i)$  is non-decreasing and  $0 \leq q(\cdot) \leq 1$ , we can find an  $i \leq 1/\varepsilon$  satisfies  $q(P_{i+1}) \leq q(P_i) + \varepsilon_0$ . Thus  $P = P_i$  and  $Q = P_{i+1}$  are the desired sets.  $\blacksquare$

Recall the definition of the potential function  $q(\cdot)$ . Given an  $n$ -vertex graph  $G$  and  $A, B \subseteq V(G)$ , the potential function is  $q(A, B) = \frac{|A||B|}{n^2} d^2(A, B)$ . For a partition  $P = \{V_1, \dots, V_k\}$  of  $V(G)$ , let  $q(P) = \sum_{i,j} q(V_i, V_j)$ . Here we note that  $q(V_i, V_i) = d(V_i, V_i) = 0$ .

Now let  $\mathcal{A} = \{A_1, \dots, A_s\}$  be a partition of  $A$  and  $\mathcal{B} = \{B_1, \dots, B_t\}$  be a partition of  $B$ . Define  $q(\mathcal{A}, \mathcal{B}) = \sum_{1 \leq i \leq s} \sum_{1 \leq j \leq t} q(A_i, B_j)$ . Choose  $x \in A$  uniformly at random, and choose  $y \in B$  uniformly at random, independently of each other. Let  $A_x \in \mathcal{A}$  be the part containing  $x$ , and  $B_y \in \mathcal{B}$  be the part containing  $y$ . Let  $z = d(A_x, B_y)$  be the random variable. Then

$$\mathbb{E}[z] = \sum_{i \in [s]} \sum_{j \in [t]} \frac{|A_i| |B_j|}{|A| |B|} d(A_i, B_j) = \sum_{i \in [s]} \sum_{j \in [t]} \frac{e(A_i, B_j)}{|A||B|} = \frac{e(A, B)}{|A||B|} = d(A, B),$$

and

$$\mathbb{E}[z^2] = \sum_{i \in [s]} \sum_{j \in [t]} \frac{|A_i| |B_j|}{|A| |B|} d^2(A_i, B_j) = \frac{n^2}{|A||B|} q(\mathcal{A}, \mathcal{B}).$$

**Lemma 1.2.** *For any sequence of constant  $\varepsilon_0 \geq \varepsilon_1 \geq \dots > 0$ , there exists a  $\delta_{1.2} > 0$  such that every  $n$ -vertex graph  $G$  has an equipartition  $V_1, V_2, \dots, V_k$  and  $W_i \subseteq V_i$  for  $i \in [k]$  satisfying*

- $|W_i| \geq \delta_{1.2} n$ .
- $(W_i, W_j)$  is  $\varepsilon_k$ -regular for all  $1 \leq i < j \leq k$ .
- All but at most  $\varepsilon_0 \binom{k}{2}$  of the pairs  $\{i, j\}$  with  $1 \leq i < j \leq k$  satisfy  $|d(V_i, V_j) - d(W_i, W_j)| \leq \varepsilon_0$ .

*Proof.* By the regularity lemma, there exist two equipartitions  $P$  and  $Q$  of  $V(G)$  such that the following hold.  $P = \{V_1, V_2, \dots, V_k\}$  is  $\varepsilon_0^3$ -regular,  $Q$  is  $\varepsilon_k$ -regular,  $Q = \{W_i^j\}_{i \in [k], j \in [t]}$ , where  $\{W_i^j\}_{j \in [t]}$  is a partition of  $V_i$ ,  $|W_i^j| \geq \delta n$  and  $q(Q) \leq q(P) + \varepsilon_0^3/4$ .

Choose  $x \in V(G)$  uniformly at random, and choose  $y \in V(G)$  uniformly at random, independently of each other. Let  $z_P = d(V_i, V_j)$  where  $V_i \ni x$  and  $V_j \ni y$ . Similarly, let  $z_Q = d(W_i^\alpha, W_j^\beta)$  where  $W_i^\alpha \ni x$  and  $W_j^\beta \ni y$ . So we get  $\mathbb{E}[z_P^2] = q(P, P) = 2q(P)$  and  $\mathbb{E}[z_Q^2] = q(Q, Q) = 2q(Q)$ . Also we have

$$\mathbb{E}[z_P z_Q] = \sum_{i \in [k]} \sum_{j \in [k]} \frac{|V_i|}{n} \frac{|V_j|}{n} d(V_i, V_j) \mathbb{E}[d(W_i^\alpha, W_j^\beta)] = \sum_{i \in [k]} \sum_{j \in [k]} \frac{|V_i|}{n} \frac{|V_j|}{n} d^2(V_i, V_j) = q(P, P) = \mathbb{E}[z_P^2],$$

then

$$\mathbb{E}[(z_Q - z_P)^2] = \mathbb{E}[z_Q^2] + \mathbb{E}[z_P^2] - 2\mathbb{E}[z_Q z_P] = \mathbb{E}[z_Q^2] - \mathbb{E}[z_P^2] = 2q(Q) - 2q(P) \leq \varepsilon_0^3/2.$$

This implies that  $|d(V_i, V_j) - d(W_i^\alpha, W_j^\beta)| \leq \varepsilon_0$  holds for all  $(W_i^\alpha, W_j^\beta)$  but at most  $\varepsilon_0(tk)^2/2$  pairs.

Now for each  $1 \leq i \leq k$ , choose  $W_i^\alpha$  uniformly at random, independently of each other. So we get  $k$  disjoint sets. Since there are at most  $\varepsilon_0(tk)^2/2$  pairs  $(W_i^\beta, W_j^\theta)$  with  $|d(V_i, V_j) - d(W_i^\beta, W_j^\theta)| \geq \varepsilon_0$ , the expected number of such pairs in the sets we chosen is no more than  $\varepsilon_0(tk)^2 t^{k-2}/t^k = \varepsilon_0 k^2/2$ . So with probability at least  $1/2$ , the sets we chosen contain no more than  $\varepsilon_0 k^2$  such pairs. Meanwhile, we know  $G$  has at most  $\varepsilon_k(tk)^2$  non- $\varepsilon_k$ -regular  $(W_i^\alpha, W_j^\beta)$  pairs. Thus with probability more than  $1/2$ , all  $(W_i^\alpha, W_j^\beta)$  pairs we chosen are  $\varepsilon_k$ -regular, since  $\varepsilon_k(tk)^2 t^{k-2}/t^k < 1/2$  where  $\varepsilon_k < 1/(2k^2)$ . Together, there exists a choice  $k$  sets, say  $W_1, W_2, \dots, W_k$ , such that all pairs are  $\varepsilon_k$ -regular and all but at most  $\varepsilon_0 \binom{k}{2}$  pairs satisfy  $|d(V_i, V_j) - d(W_i, W_j)| \leq \varepsilon_0$ , as desired.  $\blacksquare$

**Lemma 1.3.** *For every  $l, \gamma$  there exist  $\delta_{1.3} = \delta(t, \gamma)$  such that for every graph  $G$  there is a disjoint vertex sets  $W_1, W_2, \dots, W_\ell$  satisfying*

- $|W_i| \geq \delta_{1.3}n$ .
- All  $\binom{\ell}{2}$  pairs are  $\gamma$ -regular.
- Either all pairs are with densities at least  $1/2$ , or all pairs are with densities less than  $1/2$ .

*Proof.* By regularity lemma, we can get an equipartition  $\{V_1, \dots, V_k\}$  of the of  $G$  such that all but at most  $(r-1)^{-1} \binom{k}{2}$  pairs of them are  $\gamma$ -regular, with  $k$  to be chosen later. Then apply Turán's Theorem, there exist  $r$  sets  $V_i$  such that all pairs of them are  $\gamma$ -regular. Finally use Ramsey Theorem, we can take  $\ell$  sets among these  $r$  sets such that all pairs are with densities at least  $1/2$ , or all pairs are with densities less than  $1/2$ . Obviously, we know such  $k, r$  exists.  $\blacksquare$

We are ready to prove the induced graph removal lemma.

**Theorem 1.4** (Induced removal lemma). *For any graph  $H$  and  $\varepsilon > 0$ , there is a constant  $\delta > 0$  such that any  $n$ -vertex graph has at most  $\delta n^{|V(H)|}$  induced copies of  $H$  can be made induced  $H$ -free by deleting or adding  $\varepsilon n^2$  edges.*

*Proof.* Let  $h = |V(H)|$ . Apply Lemma 1.2 to  $G$  with  $\varepsilon_i = \varepsilon/3$ , we can get an equipartition  $V_1, \dots, V_2$  and  $W_i \subseteq V_i$  for  $i \in [k]$  with the following properties. All pairs  $(W_i, W_j)$  are  $\varepsilon/3$ -regular and all but at most  $\varepsilon \binom{k}{2}/3$  of the them satisfy  $|d(V_i, V_j) - d(W_i, W_j)| \leq \varepsilon/3$ . Next we use Lemma 1.3 on the induced subgraph  $G[W_i]$  to obtain sets  $W_{i,1}, W_{i,2}, \dots, W_{i,h}$ , where  $\ell = h$  and  $\gamma = \varepsilon/3$ .

Next we construct a new graph  $G_1$  from  $G$  by add and removing the following edges.

- For  $1 \leq i < j \leq k$  such that  $|d(V_i, V_j) - d(W_i, W_j)| \geq \varepsilon/3$ , remove all edges in  $(W_i, W_j)$  if  $d(W_i, W_j) < 1/2$  and add all missing edges while  $d(W_i, W_j) > 1/2$ .
- For the rest  $1 \leq i < j \leq k$ , if  $d(W_i, W_j) \leq 2\varepsilon/3$ , remove all edges in  $(W_i, W_j)$ , and if  $d(V_i, V_j) \geq 1 - 2\varepsilon/3$ , add all missing edges in  $(V_i, V_j)$ .
- For a fixed  $i$ , if all densities of pairs from  $W_{i,1}, W_{i,2}, \dots, W_{i,h}$  are less than  $1/2$ , all edges in  $G[V_i]$  are removed. Otherwise, all those densities are at least  $1/2$ , then we add all missing edges in  $G[V_i]$ .

By doing this, the total number of edges removed and added is at most  $\varepsilon n^2$ . Now if the resulting graph  $G_1$  is induced  $H$ -free, then we are done. So we may assume  $G_1$  has an induced copy of  $H$ , and the vertices of it are from  $h$  parts with repetition, say  $V_{i_1}, V_{i_2}, \dots, V_{i_h}$ . Let  $G_2$  be the induced subgraph of  $G$  with vertices  $W_{i_1,1}, W_{i_2,2}, \dots, W_{i_h,h}$ , clearly these sets are disjoint and pairwise  $\varepsilon$ -regular. Also we can check that they satisfy the density condition required for graph counting lemma, so there are at least  $\delta n^h$  induced copies of  $H$  in  $G$  where  $\delta = (\delta_{1,2}\delta_{1,3}\varepsilon/3)^h$ , a contradiction. ■