

Thm 1 (Graph Counting Lemma) Let  $H$  be a graph with  $V(H) = [k]$ .  
For  $\varepsilon > 0$ , let  $G$  be a graph with vertex subsets  $V_1 \subseteq \dots \subseteq V_k \subseteq V(G)$  such that  $(V_i, V_j)$  is  $\varepsilon$ -regular whenever  $i, j \in E(H)$ .

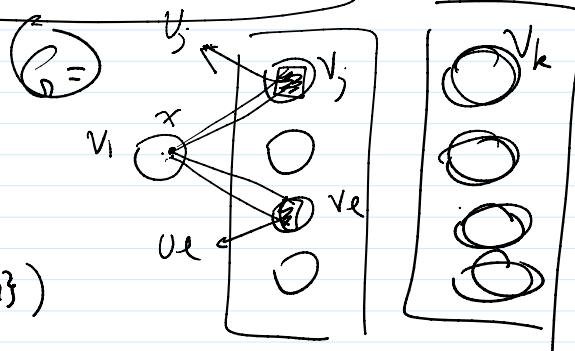
Then, the number of tuples  $(v_1, v_2, \dots, v_k) \in V_1 \times V_2 \times \dots \times V_k$  such that  $v_i v_j \in E(G)$  whenever  $i, j \in E(H)$  is

$$\left( \prod_{ij \in E(H)} d(v_i, v_j) \pm c_H \varepsilon \right) \cdot \prod_{i \in [k]} |V_i|,$$

where  $c_H$  is a constant depending on  $H$ .

$$\begin{aligned} t &= (a \pm \varepsilon)b \\ \Leftrightarrow \left| \frac{t}{b} - a \right| &\leq \varepsilon \end{aligned}$$

pf 1 (Sketch)



Fix  $x \in V_1$ . Define the following subsets  $U_j \subseteq V_j$  for  $j \neq 1$ .

If  $j \in N(1)$ , then  $U_j = V_j \cap N(x)$ ;

o.w.  $j \notin N(1)$ , then  $U_j = V_j$

$\Rightarrow (U_2, U_3, \dots, U_k)$  and we count the number of tuples  $(u_2, u_3, \dots, u_k) \in U_2 \times \dots \times U_k$  s.t.  $u_i u_j \in E(G)$  whenever  $i, j \in E(H')$  where  $H' = H \setminus \{1\}$

H/W EX (to extend this to a full proof).  $\square$

Thm 1 (Graph Counting Lemma) Let  $H$  be a graph with  $V(H) = [k]$ .

For  $\varepsilon > 0$ , let  $G$  be a graph with vertex subsets  $V_1 \subseteq \dots \subseteq V_k \subseteq V(G)$  such that  $(V_i, V_j)$  is  $\varepsilon$ -regular whenever  $i, j \in E(H)$ . Then

the number of tuples  $(v_1, \dots, v_k) \in V_1 \times V_2 \times \dots \times V_k$  such that  $v_i v_j \in E(G)$  whenever  $i, j \in E(H)$  is

$$\left( \prod_{ij \in E(H)} d(v_i, v_j) \pm c_H \varepsilon \right) \cdot \prod_{i \in [k]} |V_i|.$$

(\*)

$$\left( \prod_{ij \in E(H)} d(v_i, v_j) \pm e(H) \cdot \varepsilon \right) \cdot \prod_{i \in [k]} |V_i|.$$

pf 2 this can be rephrased into the following probabilistic form:  
choose  $v_1 \in V_1, \dots, v_k \in V_k$ , uniformly and independently at random;

If we let  $p = \Pr(v_i, v_j \in E^c(G) \text{ for all } ij \in E(H))$ ,

then  $(*)$  is the same as

$$\left| p - \prod_{ij \in E(H)} d(v_i, v_j) \right| \leq e(H) \cdot \varepsilon$$



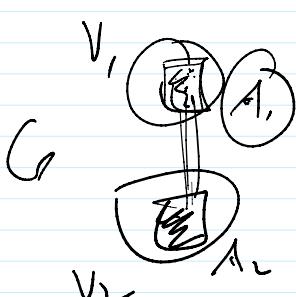
Assume  $v_2 \in E(H)$ .

Let  $p' = \Pr(v_i, v_j \in E^c(G) \text{ for all } ij \in E(H) \setminus \{v_2\})$ .

$$\text{We claim: } |p - d(V_1, V_2) p'| \leq \varepsilon \quad (1)$$

Pf of claim It suffices to show (1) holds whenever  $v_3, \dots, v_h$  are fixed arbitrarily and only  $v_1, v_2$  are random.

Define  $A_1 = \{v_i \in V_1 : v_i, v_j \in E^c(G) \text{ whenever } j \in N_H(v_2) \setminus \{v_1\}\}$   
 $A_2 = \{v_i \in V_2 : v_i, v_j \in E^c(G) \text{ whenever } j \in N_H(v_1) \setminus \{v_2\}\}$



$$\begin{cases} p = \frac{e(A_1, A_2)}{|V_1| \cdot |V_2|} \\ p' = \frac{|A_1| \cdot |A_2|}{|V_1| \cdot |V_2|} \end{cases}$$

Therefore, it suffices to show  $\left| \frac{e(A_1, A_2)}{|V_1| |V_2|} - d(V_1, V_2) \cdot \frac{|A_1| |A_2|}{|V_1| |V_2|} \right| \leq \varepsilon$ . (2)

If  $|A_1| \leq \varepsilon |V_1|$  or  $|A_2| \leq \varepsilon |V_2|$ ,

$$\text{then } \frac{e(A_1, A_2)}{|V_1| |V_2|} \leq \frac{|A_1| |A_2|}{|V_1| |V_2|} \leq \varepsilon \quad \& \quad d(V_1, V_2) \cdot \frac{|A_1| |A_2|}{|V_1| |V_2|} \leq \varepsilon$$

thus the inequality (2) holds. ✓

Otherwise, we have  $|A_1| > \varepsilon |V_1|$  and  $|A_2| > \varepsilon |V_2|$ .

By  $\varepsilon$ -regularity of  $(V_1, V_2)$ ,

$$\left| \frac{e(A_1, A_2)}{|V_1| |V_2|} - d(V_1, V_2) \cdot \frac{|A_1| |A_2|}{|V_1| |V_2|} \right| = \left| \frac{e(A_1, A_2)}{|A_1| \cdot |A_2|} - d(V_1, V_2) \cdot \frac{|A_1| |A_2|}{|V_1| |V_2|} \right| \leq \varepsilon$$

This proves the claim  $\square$

To complete the proof, we do induction on  $e(H)$ .

Let  $H' = H \setminus \{v_2\}$ . We have

$$\left| \beta - \prod_{ij \in E(H)} d(V_i, V_j) \right| \leq \left| \beta - d(V_1, V_2) \cdot \beta' + d(V_1, V_2) \cdot \prod_{ij \in E(H')} d(V_i, V_j) \right|$$

by claim  $\leq \varepsilon + d(V_1, V_2) \cdot (e(H') \cdot \varepsilon) \leq e(H) \cdot \varepsilon.$   $\square$

Theorem 2 (Graph removal Lemma) For any graph  $H$  and  $\varepsilon > 0$ ,

$\exists \delta = \delta(H, \varepsilon) > 0$  s.t. any  $n$ -vertex graph with less than

$\delta n^{|V(H)|}$  copies of  $H$  can be made  $H$ -free by deleting at most  $\varepsilon n^2$  edges.

pf.:  $H \in \mathcal{F}_X$ . (Similar to triangle removal lemma;  
use Graph Combing Lemma)  $\square$

Def. For two graphs  $G$  and  $H$  with the same number of vertices,  
the edit-distance  $d(G, H)$  is the minimum integer  $k$  such that  
 $G$  can be obtained from  $H$  by adding or deleting  $k$  edges.

$$\Leftrightarrow d(G, H) = |E(G) \Delta E(H)|$$

$\uparrow$  symmetric difference

Next lecture, we will show Erdős-Simonovits Stability Theorem!