



Meshes: Definitions & Terminologies

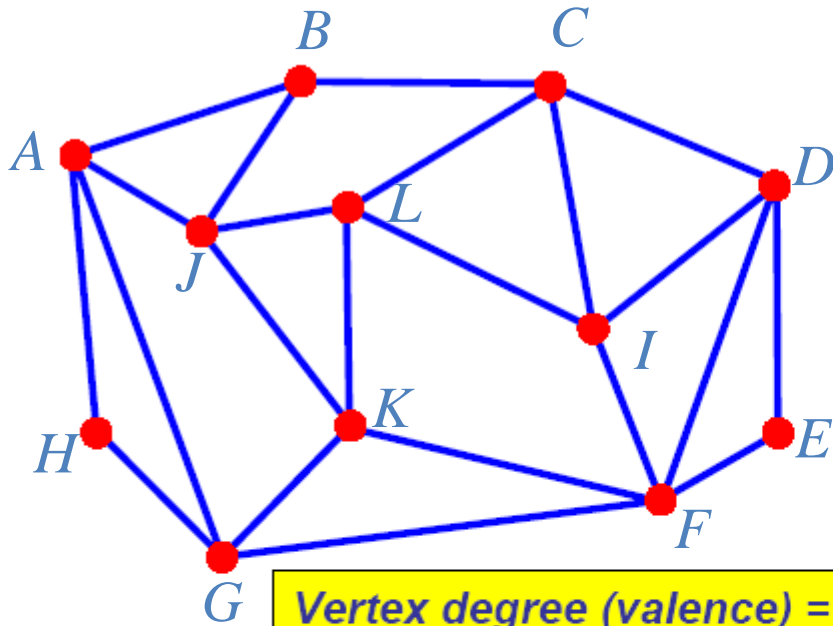
Ligang Liu

Graphics&Geometric Computing Lab

USTC

<http://staff.ustc.edu.cn/~lgliu>

Standard Graph Definition



$G = \langle V, E \rangle$

$V =$ vertices =

$\{A, B, C, D, E, F, G, H, I, J, K, L\}$

$E =$ edges =

$\{(A, B), (B, C), (C, D), (D, E), (E, F), (F, G),$
 $(G, H), (H, A), (A, J), (A, G), (B, J), (K, F),$
 $(C, L), (C, I), (D, I), (D, F), (F, I), (G, K),$
 $(J, L), (J, K), (K, L), (L, I)\}$

Vertex degree (valence) = number of edges incident on vertex

$\deg(J) = 4, \deg(H) = 2$

***k*-regular graph** = graph whose vertices all have degree *k*

Face: cycle of vertices/edges which cannot be shortened

$F =$ faces =

$\{(A, H, G), (A, J, K, G), (B, A, J), (B, C, L, J), (C, I, J), (C, D, I),$
 $(D, E, F), (D, I, F), (L, I, F, K), (L, J, K), (K, F, G)\}$

Connectivity

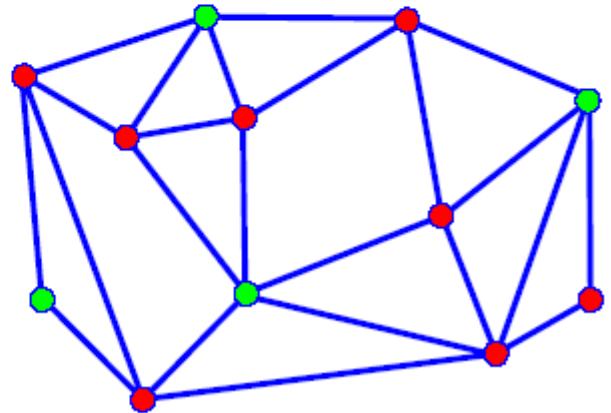
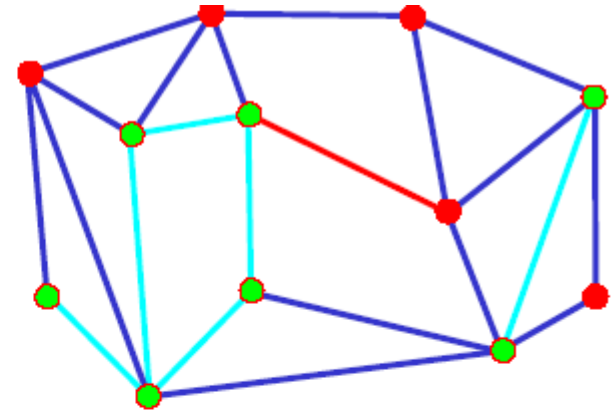
Graph is *connected* if there is a path of edges connecting every two vertices

Graph is *k-connected* if between every two vertices there are k edge-disjoint paths

Graph $G' = \langle V', E' \rangle$ is a *subgraph* of graph $G = \langle V, E \rangle$ if V' is a subset of V and E' is the subset of E incident on V'

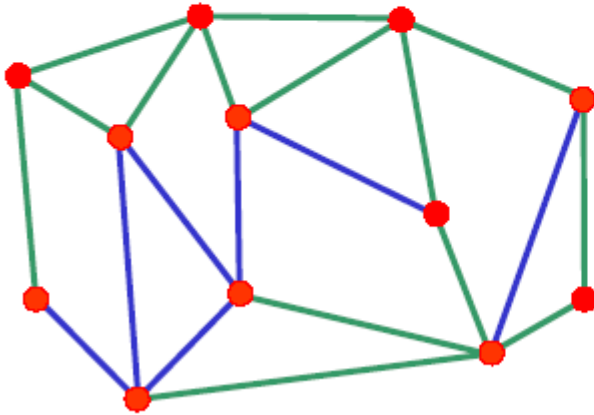
Connected component of a graph: maximal connected subgraph

Subset V' of V is an *independent set* in G if the subgraph it induces does not contain any edges of E

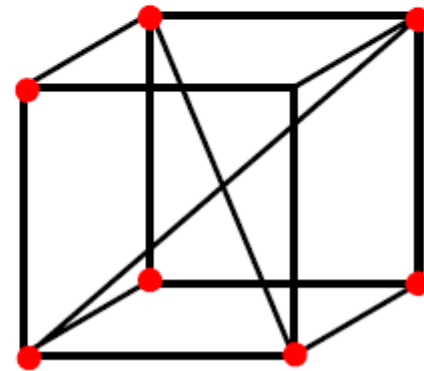


Graph Embedding

Graph is *embedded* in \mathbb{R}^d if each vertex is assigned a position in \mathbb{R}^d



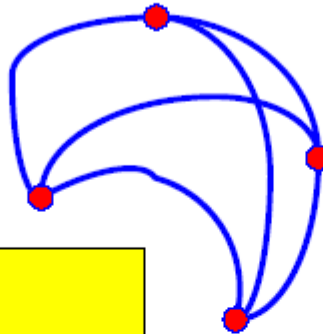
Embedding in \mathbb{R}^2



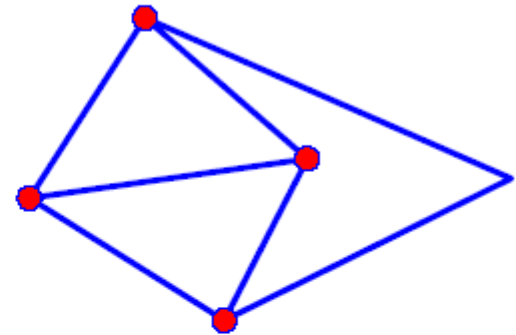
Embedding in \mathbb{R}^3

Planar Graphs

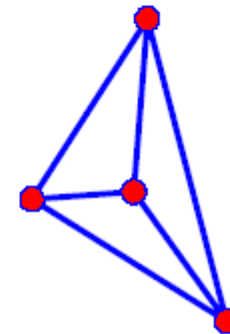
Planar Graph



Plane Graph



Straight Line Plane Graph



Planar graph: graph whose vertices and edges can be embedded in \mathbb{R}^2 such that its edges do not intersect

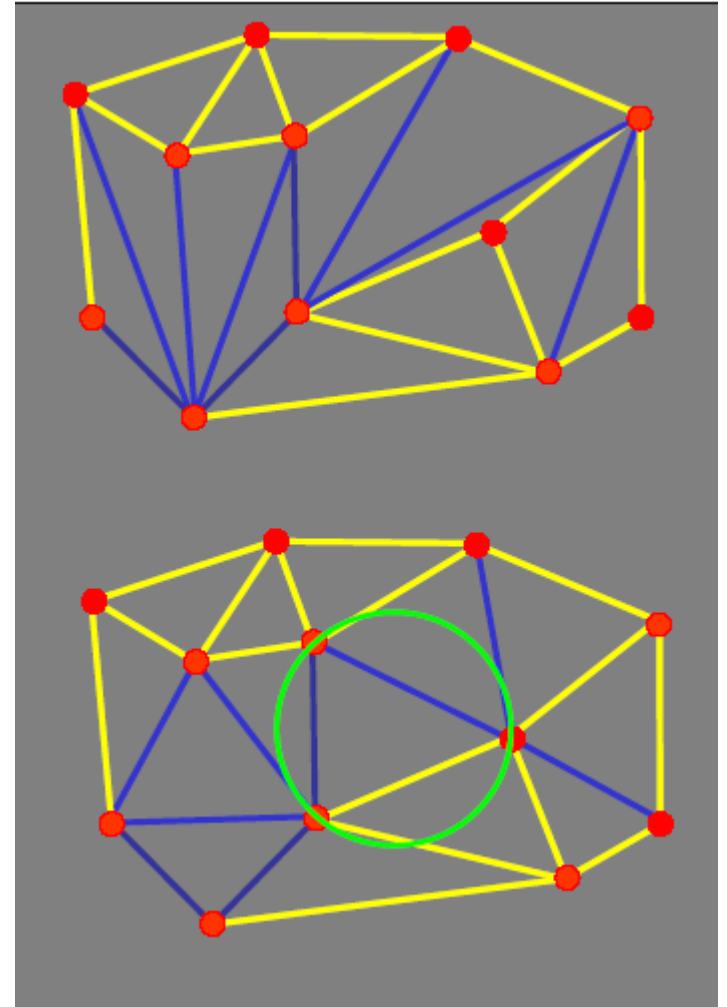
Every planar graph can be drawn as a ***straight-line plane graph***

Triangulation

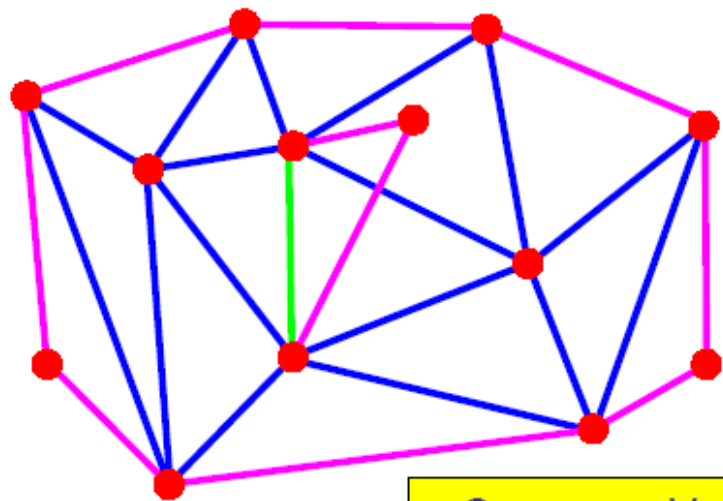
Triangulation: straight line plane graph all of whose faces are triangles

Delaunay triangulation of a set of points: unique set of triangles such that the circumcircle of any triangle does not contain any other point

Delaunay triangulation avoids long and skinny triangles



Meshes



Mesh: straight-line graph embedded in \mathbb{R}^3

Boundary edge: adjacent to exactly *one* face

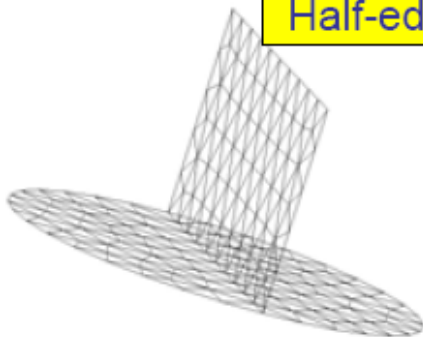
Regular edge: adjacent to exactly *two* faces

Singular edge: adjacent to more than two faces

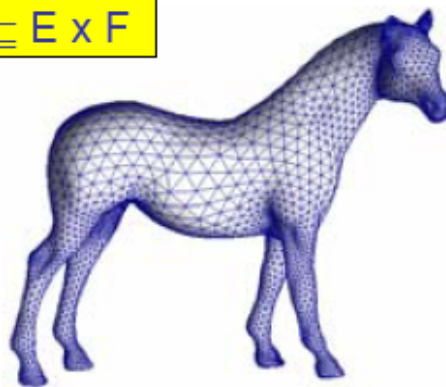
Closed mesh: mesh with no boundary edges

Manifold mesh: mesh with no singular edges

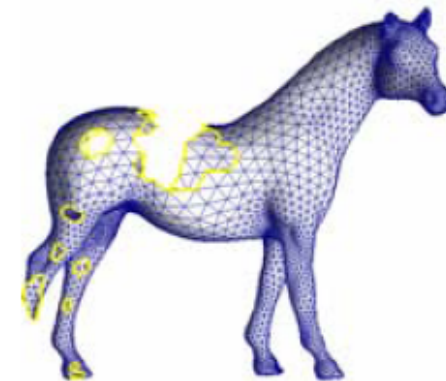
Corners $\subseteq V \times F$
Half-edges $\subseteq E \times F$



Non-Manifold

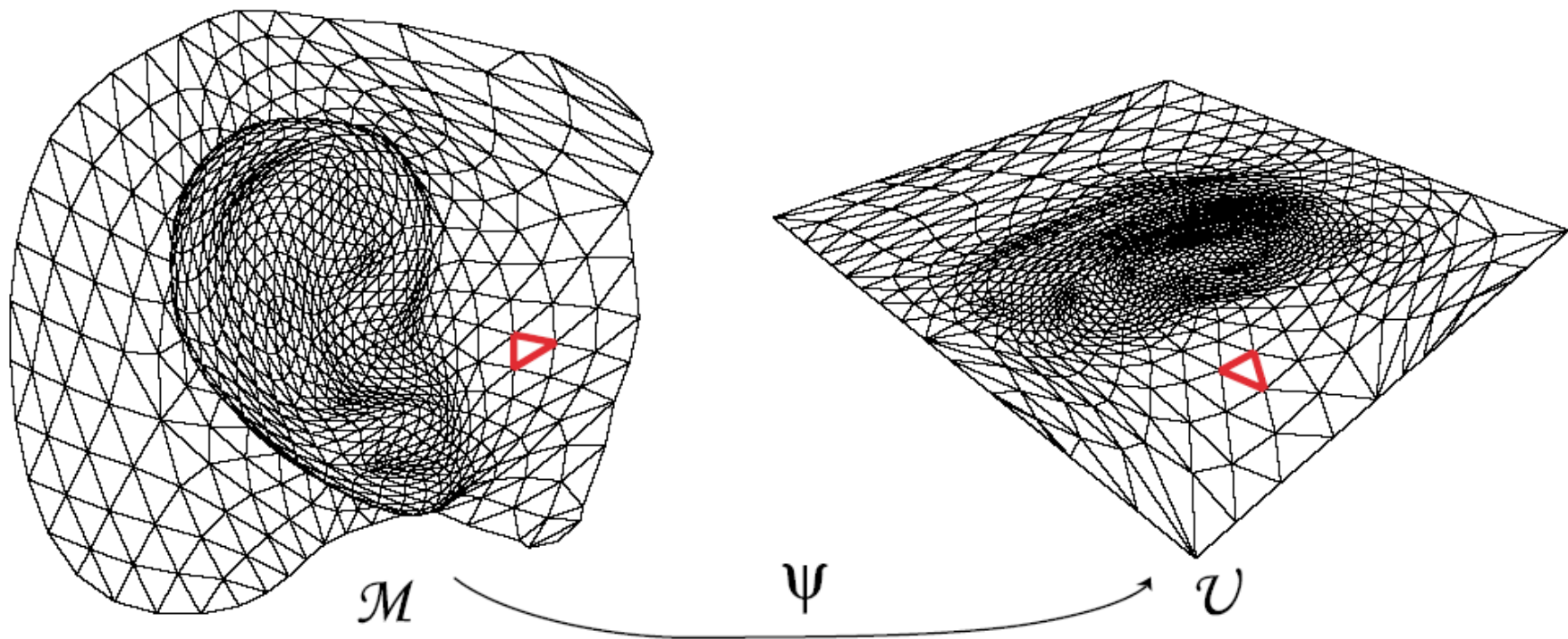


Closed Manifold

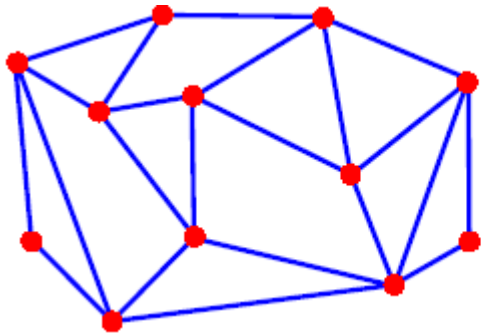


Open Manifold

Planar Graphs and Meshes



Topology



$v = 12$
 $f = 14$
 $e = 25$
 $c = 1$
 $g = 0$
 $b = 1$

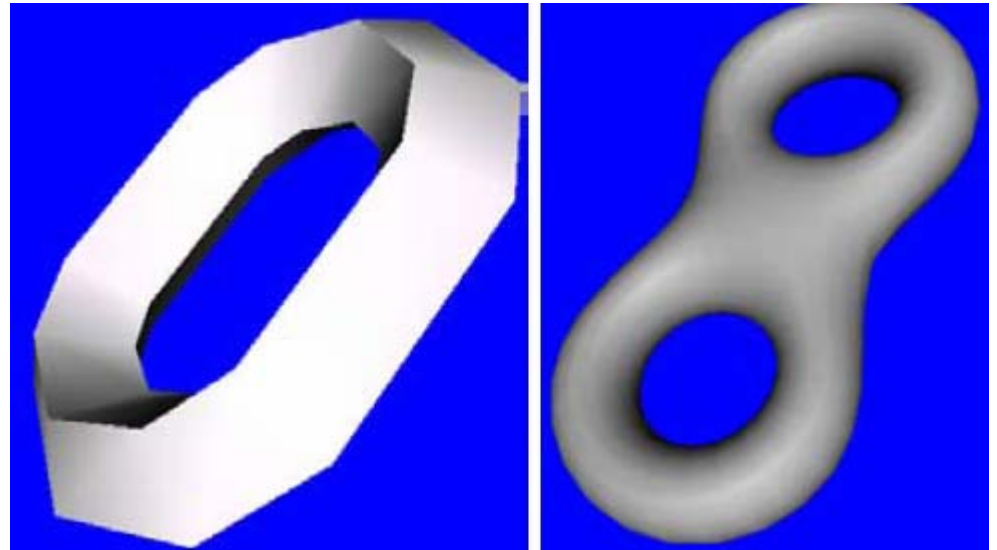
Genus of graph: *half* of maximal number of closed paths that do *not* disconnect the graph (number of “holes”)

Genus(sphere) = 0
Genus(torus) = 1

Euler-Poincare Formula

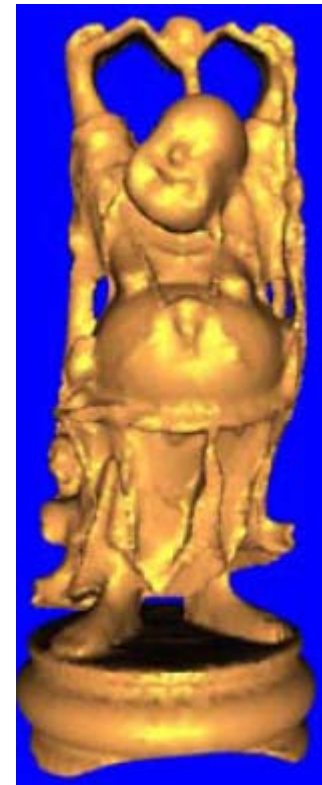
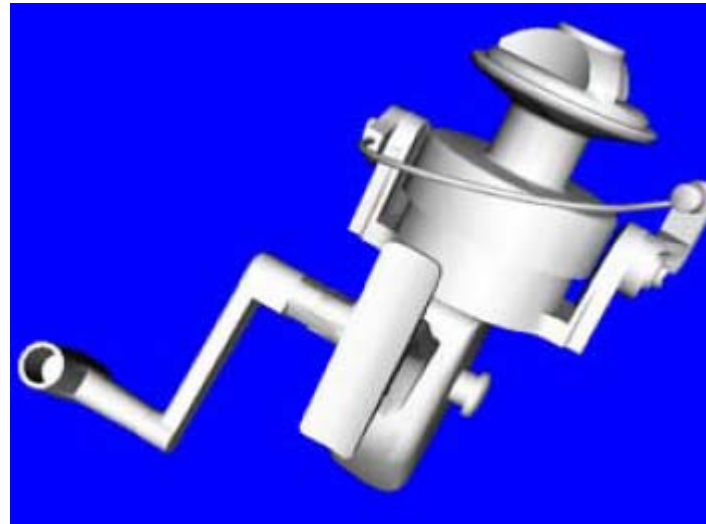
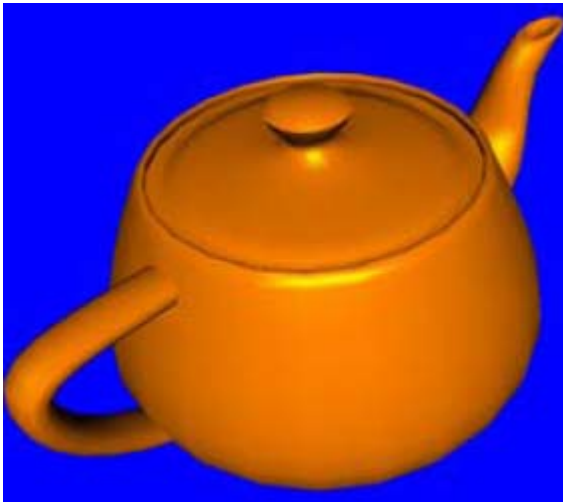
$$v + f - e = 2(c - g) - b$$

v = # vertices c = # conn. comp
 f = # faces g = genus
 e = # edges b = # boundaries

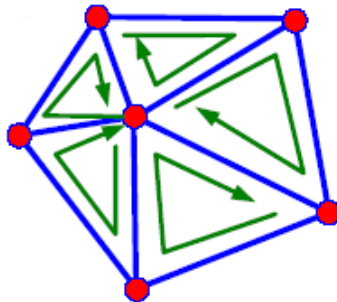


Topology Quiz

What can you say about the genus of these meshes ?



Orientability



Orientation of a face is clockwise or anticlockwise order in which its vertices and edges are listed

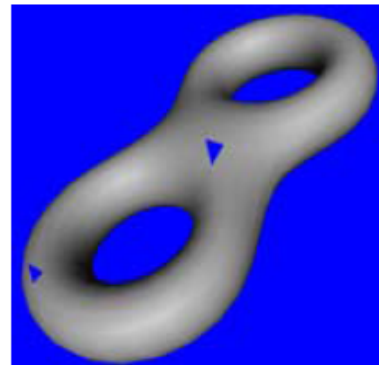
This defines the direction of face *normal*

Oriented
 $F = \{(L, J, B), (B, C, L), (L, C, I), (I, K, L), (L, K, J)\}$
Not Oriented
 $F = \{(B, J, L), (B, C, L), (L, C, I), (L, I, K), (L, K, J)\}$

Straight line graph is *orientable* if orientations of its faces can be chosen so that each edge is oriented in *both* directions



Not Backface Culled

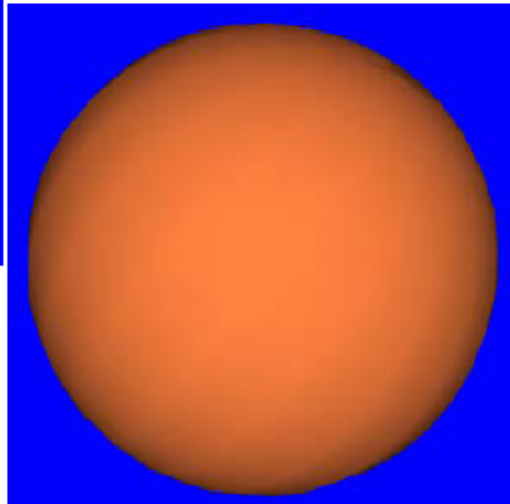
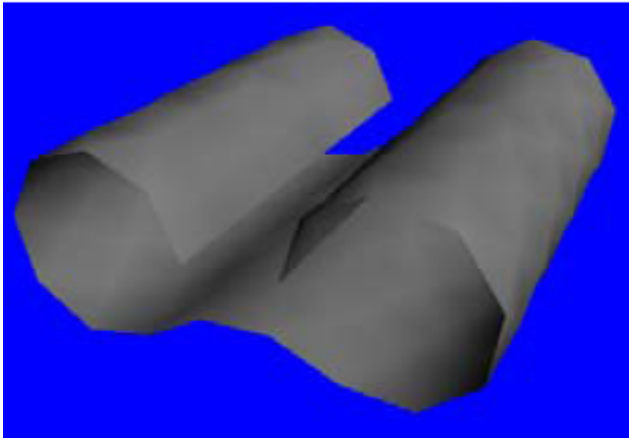


Backface Culled

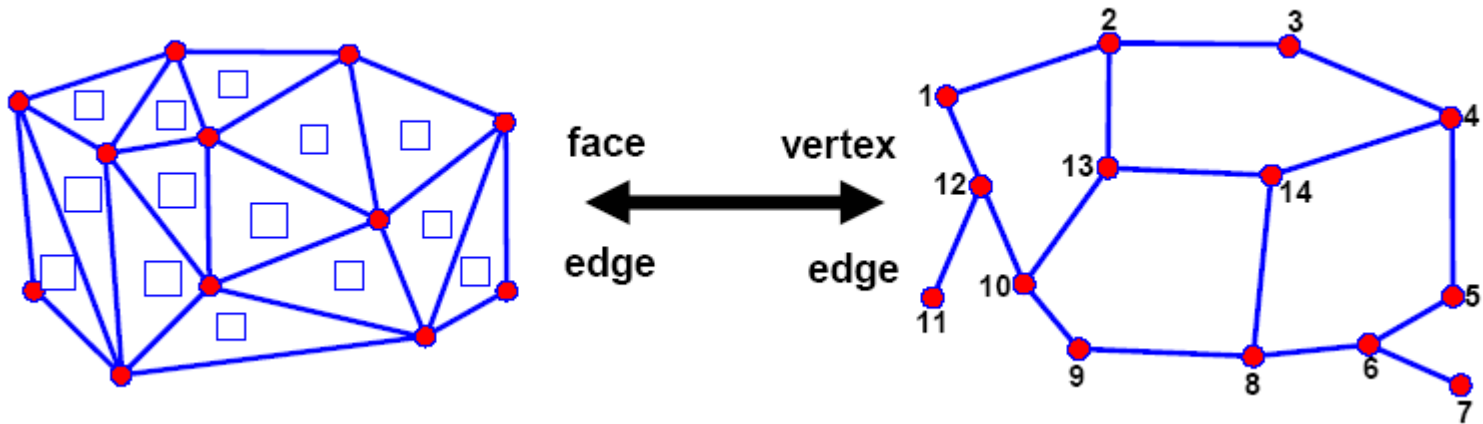
Mobius strip or Klein bottle - not orientable

Developability

Mesh is *developable* if it may be embedded in \mathbb{R}^2 without distortion



Duality



- Delaunay Triangulation vs. Voronoi Graph

Q&A