

Local Strongly Arc-Connectivity in Regular Bipartite Digraphs*

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Received October 10, 1991

We show that for any vertex x of a d -regular bipartite digraph there are a vertex y , in the other class of the bipartition, and $d(x, y)$ -paths and $d(y, x)$ -paths such that all $2d$ of them are pairwise arc-disjoint. This result generalizes a theorem of Hamidoune and Las Vergnas for graphs. © 1993 Academic Press, Inc.

1. INTRODUCTION

Hamidoune and Las Vergnas [1] proved that for any vertex x of a d -regular bipartite graph G there is a vertex y of G , in the other class of the bipartition, joined to x by d pairwise edge-disjoint paths. In the present note we generalize this to digraphs by using the same techniques used in [1].

Let $D = (V, A)$ be a digraph (possibly with multiple arcs and loops) with the vertex-set $V = V(D)$ and the arc-set $A = A(D)$. A digraph D is said to be d -regular if both the outdegree $d_D^+(x)$ and the indegree $d_D^-(x)$ are equal to d for any $x \in V(D)$. Our result can be stated as follows:

THEOREM. *Let D be a d -regular bipartite digraph. For all vertex x of D there are a vertex y of D , in the other class of the bipartition, and $d(x, y)$ -paths P_1, P_2, \dots, P_d and $d(y, x)$ -paths Q_1, Q_2, \dots, Q_d such that all $2d$ of them are pairwise arc-disjoint.*

2. PROOF OF THEOREM

For two disjoint nonempty and proper subsets X and Y of $V(D)$ we denote by $A_D(X, Y)$ the set of the arcs of D joining X to Y and by $D[X]$ the subdigraph of D induced by X . Let $\bar{X} = V \setminus X$.

* This work was partially supported by Natural Science Foundation of China.

LEMMA 1 (Kotzig [2], cf. Problem 6.47 in [3]). *Let x and y be two distinct vertices of a digraph D and the local strongly arc-connectivity be k between x and y . If the outdegree of each vertex of D is equal to its indegree, then there are k (x, y) -paths and k (y, x) -paths which are mutually arc-disjoint.*

LEMMA 2. *Let D be a digraph and X be a nonempty and proper subset of $V(D)$. If the outdegree of each vertex of D is equal to its indegree, then $|A_D(X, \bar{X})| = |A_D(\bar{X}, X)|$.*

Proof. In fact, since

$$\sum_{x \in X} d_{D[X]}^+(x) = \sum_{x \in X} d_{D[X]}^-(x),$$

and $d_D^+(x) = d_D^-(x)$ for each $x \in V(D)$ we have immediately that

$$\begin{aligned} |A_D(X, \bar{X})| &= \sum_{x \in X} d_D^+(x) - \sum_{x \in X} d_{D[X]}^-(x) \\ &= \sum_{x \in X} d_D^-(x) - \sum_{x \in X} d_{D[X]}^-(x) = |A_D(\bar{X}, X)|. \end{aligned}$$

Proof of Theorem. The proof is by induction on $|V(D)|$. For $|V(D)| = 2$ the theorem is obvious (note that a bipartite digraph has no loops).

Let $V(D) = S \cup T$ be the bipartition of $V(D)$ and, without loss of generality, $x \in S$. If $|A_D(X, \bar{X})| \geq d$ for every proper subset X of $V(D)$ containing x , then the theorem follows immediately from Menger's theorem and Lemma 1. Next we suppose that there is a proper subset X of $V(D)$ such that $x \in X$ and $|A_D(X, \bar{X})| < d$. We can furthermore suppose that X is chosen in such a way that $|A_D(X, \bar{X})|$ takes the minimum possible value. Then $X \cap S \neq \emptyset$ and $X \cap T \neq \emptyset$ by our choice of X . Let

$$\begin{aligned} A_D(X, T \setminus X) &= \{a_1, a_2, \dots, a_s\}, & A_D(X, S \setminus X) &= \{a_{s+1}, a_{s+2}, \dots, a_{s+t}\}, \\ A_D(S \setminus X, X) &= \{b_1, b_2, \dots, b_{s'}\}, & A_D(T \setminus X, X) &= \{b_{s'+1}, b_{s'+2}, \dots, b_{s'+t'}\}. \end{aligned}$$

Then, by Lemma 2,

$$s + t = |A_D(X, \bar{X})| = |A_D(\bar{X}, X)| = s' + t'.$$

Considering that $X = (X \cap S) \cup (X \cap T)$ is the bipartition of the bipartite digraph $D[X]$ and counting the number of the arcs in $D[X]$ with initial vertices in $X \cap S$ and terminal vertices in $X \cap Y$, we have

$$d |X \cap S| - s = d |X \cap T| - s'.$$

Noting that $s + t = s' + t' < d$, we obtain, from the above,

$$s = s', \quad t = t'.$$

Let D' be the digraph obtained from $D[X]$ by adding the arcs a'_i ($i = 1, 2, \dots, s + t$) joining the initial vertex of a_i ($i = 1, 2, \dots, s + t$) so the terminal vertex of b_i ($i = 1, 2, \dots, s' + t'$). Clearly D' is again a d -regular bipartite digraph and $|V(D')| < |V(D)|$. Hence by the induction hypothesis there are $y \in X \cap T$ and $d(x, y)$ -paths P'_1, P'_2, \dots, P'_d and $d(y, x)$ -paths Q'_1, Q'_2, \dots, Q'_d such that all $2d$ of them are pairwise arc-disjoint.

Let D'' be the digraph obtained from D by deleting all arcs in $D[X]$ and then identifying with x all vertices in X . Clearly the outdegree of each vertex of D'' is equal to its indegree. By the minimality of $|A_D(X, \bar{X})| = s + t$ and Lemma 2, the digraph D'' is strongly $(s + t)$ -arc-connected. Hence by Lemma 1, for any $z \in \bar{X}$ there are $(s + t)$ (x, z) -paths and $(s' + t')$ (z, x) -paths which are mutually arc-disjoint in D'' . Let $z \in \bar{X}$ and X_1, X_2, \dots, X_{s+t} and $Y_1, Y_2, \dots, Y_{s'+t'}$ be these (x, z) -paths and, respectively, (z, x) -paths labelled such that $a_i \in A(X_i)$ ($i = 1, 2, \dots, s + t$) and $b_i \in A(Y_i)$ ($i = 1, 2, \dots, s' + t'$). Then, for $j = 1, 2, \dots, d$, let P_j be the (x, y) -path of D obtained from P'_j by replacing each occurrence of an arc a'_i , $i = 1, 2, \dots, s$, in P'_j by $X_i \cap Y_i$. Similarly, for $j = 1, 2, \dots, d$, let Q_j be the (y, x) -path of D obtained from Q'_j by replacing each occurrence of an arc a_{s+i} , $i = 1, 2, \dots, t$ in Q'_j by $X_{s+i} \cup Y_{s+i}$. Clearly $P_1, P_2, \dots, P_d, Q_1, Q_2, \dots, Q_d$ have the required properties which proves the theorem.

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