

An infinite family of 4-tight optimal double loop networks

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Abstract An infinite family of 4-tight optimal double loop networks is given in this paper.

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Because of their symmetry, simplicity and extensionality, the double loop networks (DLNs) have been widely used in the topological design of local networks, multi-module memory organizations, data alignments in parallel memory systems, and supercomputer architecture^[1]. The graphical model of a DLN is a digraph (also called circulant digraph) $G(n; s)$ with the vertex set $\{0, 1, 2, \dots, n-1\}$ and the edge set $\{i \rightarrow i+1 \pmod{n}, i \rightarrow i+s \pmod{n} : i = 0, 1, 2, \dots, n-1\}$, where s is a given integer with $1 < s < n$. From the definition, it is clear that $G(n; s)$ is strongly connected and the diameter is only determined by n and s . Denote by $d(n; s)$ the diameter of $G(n, s)$ and let $d(n) = \min\{d(n; s) : 1 < s < n\}$. A network $G(n; s)$ is optimal if $d(n; s) = d(n)$. Wong and Coppersmith^[2] have shown that $d(n) \geq lb(n) = \lceil \sqrt{3n} \rceil - 2$, where the symbol $\lceil m \rceil$ denotes the smallest integer not less than the real number m .

An important problem is to determine the value of $d(n)$ and find s such that $G(n; s)$ is optimal for any given $n \geq 4$. The problem has attracted many authors' interest^[2-10], although it seems impossible to express the function $d(n)$ in a closed form.

Let Z be the infinite set of all nonnegative integers. For $k \in Z$, an optimal $G(n; s)$ is said to be k -tight if $d(n; s) = lb(n) + k$. Generally, 0- or 1-tight optimal DLNs are called tight optimal and near-tight optimal, respectively^[7,8]. We say that $\{G(n(t); s(t)) : t \in Z, t \geq t_0\}$ is an infinite family of k -tight optimal DLNs if $G(n(t); s(t))$ is k -tight optimal for any $t \in Z$ and $t \geq t_0$, where $G(n(t_0); s(t_0))$ is the initial element. We say that $\{n(t) : t \in Z, t \geq t_0\}$ contains no k -tight optimal DLN if $d(n(t); s(t)) > lb(n(t)) + k$ for any $s(t)$ and $t \geq t_0$. The functions $n(t)$ and $s(t)$ are polynomials in $t \in Z$ with integral coefficients.

Li, Xu et al.^[7] have presented a systematic method to construct optimal DLNs and listed 102 infinite families of optimal DLNs, of which 69 are tight and 33 near-tight, to show that $d(n) \leq lb(n) + 1$ for $n \leq 300$. Xu^[9] has found 3 infinite families of 2-tight optimal DLNs. Erdős and Hsu^[4] reported an exhaustive computer search. Chen showed that $d(n) \leq lb(n) + 4$ for $n \leq 75000$, there exist only 3 n 's, 53749, 64729 and 69283, for which the equality holds; the corresponding 4-tight optimal DLN's are $G(53749; 985)$, $G(64729; 394)$ and $G(69283; 1764)$, and their diameters are 404, 443 and 458, respectively. With computer, we found that the fourth

4-tight optimal DLN is $G(94921; 515)$, and its diameter is 536. However, as far as we know, no infinite family of 4-tight optimal DLNs has been found so far. In this paper, combining geometric method with number theoretic technics, we will construct an infinite family of 4-tight optimal DLNs with the initial element $G(69283; 1764)$.

In the following, we consider $n(t) = 3t^2 + 6t - 26$, $t \in \mathbb{Z}$. It is easy to verify $lb(n(t)) = 3t + 1$ for $t \geq 14$. Please refer to ref. [1] or ref. [7] for terminology and notation not defined and explained in this paper.

1 Some lemmas

Lemma 1^[7]. Let $L = L(n; l, h, x, y)$ be an L-tile. If $|y - x| \geq z_0 \geq 1$, then $d(n) \geq \sqrt{3n - \frac{3}{4}z_0^2} + \frac{1}{2}z_0 - 2$.

Lemma 2^[8]. Let $n(t) = 3t^2 + 3t - 26$. Then $\{n(t) : t \in \mathbb{Z}, t \geq 29\}$ contains no tight optimal DLN. Moreover, if $L(n, l, h, x, y)$ is near-tight optimal, then $|x - y| \leq 1$.

Lemma 3^[7]. Let $n(t) = 3t^2 + 3t - 26$ and $L = L(n; l, h, x, y)$ be an L-tile, where $l = 2t + a$, $h = 2t + b$, $x = t + a + b - j$, $z = |x - y|$, a, b, x and y are polynomials in $t \in \mathbb{Z}$ with integral coefficients. Then L is k -tight if and only if

$$(a + b - j)(a + b - j + z) - ab + (6 + z - 2j) - 26 = 0 \quad (1)$$

holds for any $j = 3 + k$ ($k \in \mathbb{Z}$). Moreover, if there exist integers α and β such that $\alpha y + \beta(h - y) \equiv 1$, then there exists only one k -tight optimal $G(n(t); s(t))$, where $s \equiv \alpha l - \beta(x - l) \pmod{n}$.

2 Main results

Theorem 1. For $n \in \mathbb{Z}$, a necessary condition¹⁾ that there exist $s, m \in \mathbb{Z}$ such that $n = s^2 + 3m^2$ is that if n has a prime divisor p with $p = 2$ or $p \equiv 5 \pmod{6}$, then p has an even power in the prime decomposition of n .

Proof. We proceed by induction on $n \geq 3$. If $n = 3$ there is nothing to prove, so we suppose that the theorem holds for any n less than an integer k with $k \geq 4$. Let $n = k$ and suppose that there exist $s, m \in \mathbb{Z}$ such that $k = s^2 + 3m^2$ and k has a prime divisor p with $p = 2$ or $p \equiv 5 \pmod{6}$.

If $p = 2$, then s and m are of the same parity since 2 is a divisor of $k = s^2 + 3m^2$. If $s = 2i$ and $m = 2j$, then $k = 2^2(i^2 + 3j^2)$, so the theorem holds by the induction hypothesis. If $s = 2i + 1$ and $m = 2j + 1$, then

$$k = s^2 + 3m^2 = 4i^2 + 4i + 1 + 3(4j^2 + 4j + 1) = 2^2[i(i + 1) + 3j(j + 1) + 1].$$

The theorem holds by the induction hypothesis since the number in the square bracket is odd.

We now suppose $p \equiv 5 \pmod{6}$. If p is not a divisor of sm , since $k = s^2 + 3m^2 \equiv 0 \pmod{p}$, then $p \equiv 1 \pmod{6}$, a contradiction. Thus, p is a divisor of s or m . Since p is a divisor of $k = s^2 + 3m^2$, it follows that p is a common divisor of s and m . Let $s = pi$ and $m = pj$. Then $k = p^2(i^2 + 3j^2)$, so the theorem holds by the induction hypothesis.

Corollary. Let $n = s^2 + 3m^2$ ($s, m \in \mathbb{Z}$) and suppose that 3 is not a divisor of n . Then $n \equiv 4 \pmod{6}$ if 2 is a divisor of n , and $n \equiv 1 \pmod{6}$ otherwise.

1) It is easy to prove that this condition is also sufficient.

Theorem 2. A necessary condition¹⁾ for eq. (1) in a and b to have an integral solution is that for any $z, j, t \in Z$, if

$$H_{z,j} = (2j - z)^2 - 3[j(j - z) + (6 + z - 2j)t - 26] \quad (2)$$

has a prime divisor p with $p = 2$ or $p \equiv 5 \pmod{6}$, then p has an even power in the prime decomposition of n .

Proof. Suppose that eq. (1) in a and b has an integral solution and rewrite it as

$$a^2 + [b - (2j - z)]a + b^2 - (2j - z)b + C = 0, \quad (3)$$

where $C = j(j - z) + (6 + z - 2j)t - 26$. Then eq. (3) in a has an integral solution by our assumption. Thus, there exists an $m \in Z$ such that

$$[b - (2j - z)]^2 - 4[b^2 - (2j - z)b + C] = m^2.$$

Express it as an equation in b :

$$3b^2 - 2(2j - z)b + [4C + m^2 - (2j - z)^2] = 0. \quad (4)$$

Then eq. (4) in b has an integral solution. Thus, there exists an $n \in Z$ such that

$$4(2j - z)^2 - 12[4C + m^2 - (2j - z)^2] = n^2.$$

This implies that n is even. Let $n = 2s$ in the above expression. We have $4(2j - z)^2 - 12C = s^2 + 3m^2$, that is, $4H_{z,j} = s^2 + 3m^2$. The theorem follows by Theorem 1.

Theorem 3. Let $n(t) = 3t^2 + 6t - 26$, $t(f) = 28f^2 + 132f + 151$, $f = 22 \cdot 85^2e$. Then $\{n(t(f(e))) : e \in Z\}$ contains no k -tight optimal DLN for each $k = 0, 1, 2, 3$.

Proof. (a) By Lemma 2, $\{n(t) : t \in Z, t \geq 29\}$ contains no tight optimal DLN.

(b) If $\{n(t) : t \in Z, t \geq 29\}$ contains a near-tight optimal DLN, then there exists a near-tight tile $L = L(n; l, h, x, y)$. Let $z = |x - y|$. Then $z \leq 1$ by Lemma 2. Counting $H_{0,4}$ and $H_{1,4}$ in expression (2), we have

$$\begin{aligned} H_{0,4} &= 6t + 94 = 6 \cdot 28 \cdot 22^2 \cdot 85^4 e^2 + 6 \cdot 132 \cdot 22 \cdot 85^2 e + 1000 \\ &= 2^3(3 \cdot 7 \cdot 22^2 \cdot 85^4 e^2 + 3 \cdot 33 \cdot 22 \cdot 85^2 e + 125); \\ H_{1,4} &= 3t + 91 = 3 \cdot 28 \cdot 22^2 \cdot 85^4 e^2 + 3 \cdot 132 \cdot 22 \cdot 85^2 e + 544 \\ &= 17(3 \cdot 28 \cdot 22^2 \cdot 85^3 \cdot 5e^2 + 3 \cdot 132 \cdot 22 \cdot 85 \cdot 5e + 32). \end{aligned}$$

Note that the number in the brackets of the expression of $H_{0,4}$ is odd, and, that of the expression of $H_{1,4}$ is not a multiple of 17. It follows that neither $H_{0,4}$ nor $H_{1,4}$ satisfies the condition in Theorem 2; that is, eq. (1) in a and b has no integral solution. This implies that $\{n(t(f(e))) : e \in Z\}$ contains no near-tight optimal DLN.

(c) If $\{n(t(f(e))) : e \in Z\}$ contains a 2-tight optimal DLN, then there exists a 2-tight tile $L = L(n; l, h, x, y)$. Let $z = |x - y|$. Then for $z \geq 5$ and $t \geq 151$, we have

$$3n(t) - \frac{3}{4}5^2 = \left(3t + \frac{5}{2}\right)^2 + 3t - 97 > \left(3t + \frac{5}{2}\right)^2.$$

By Lemma 1, a contradiction can be deduced as follows:

$$3t + 3 = d(n(t)) > 3t + \frac{5}{2} + \frac{5}{2} - 2 = 3t + 3.$$

1) It is easy to prove that this condition is also sufficient.

Therefore $z \leq 4$. For $0 \leq z \leq 4$ and $j = 5$, the values of $H_{z,5}$ in expression (2) can be counted as follows:

$$H_{4,5} = 99 = 3^2 \cdot 11;$$

$$H_{3,5} = 3t + 97 = 2(3 \cdot 14f^2 + 3 \cdot 66f + 275);$$

$$H_{2,5} = 6t + 97 = 17(6 \cdot 28 \cdot 22^2 \cdot 85^2 \cdot 5e^2 + 6 \cdot 132 \cdot 22 \cdot 85 \cdot 5e + 59);$$

$$H_{1,5} = 9t + 99 = 2(9 \cdot 14f^2 + 9 \cdot 66f + 729);$$

$$H_{0,5} = 12t + 103 = 5(12 \cdot 28 \cdot 22^2 \cdot 17 \cdot 85^3 e^2 + 12 \cdot 132 \cdot 22 \cdot 17 \cdot 85e + 383).$$

It is easy to verify that no value of $H_{z,5}$ satisfies the condition in Theorem 2. Therefore, eq. (1) in a and b has no integral solution, implying that $\{n(t(f(e))) : e \in Z\}$ contains no 2-tight optimal DLN.

(d) If $\{n(t(f(e))) : e \in Z\}$ contains a 3-tight optimal DLN, then there exists a 3-tight tile $L = L(n; l, h, x, y)$. Let $z = |x - y|$. Then for $z \geq 7$ and $t \geq 151$,

$$3n(t) - \frac{3}{4}7^2 = \left(3t + \frac{5}{2}\right)^2 + 3t - 115 > \left(3t + \frac{5}{2}\right)^2.$$

By Lemma 1, a contradiction can be deduced as follows:

$$3t + 4 = d(n(t)) > 3t + \frac{5}{2} + \frac{7}{2} - 2 = 3t + 4.$$

Therefore $z \leq 6$. For $0 \leq z \leq 6$ and $j = 6$, the values of $H_{z,6}$ in expression (2) can be counted as follows:

$$H_{6,6} = 114 = 2 \cdot 57;$$

$$H_{5,6} = 3t + 109 = 2(3 \cdot 14f^2 + 3 \cdot 66f + 281);$$

$$H_{4,6} = 6t + 106 = 11(6 \cdot 28 \cdot 22 \cdot 2 \cdot 85^4 e^2 + 6 \cdot 12 \cdot 22 \cdot 85^2 e + 92);$$

$$H_{3,6} = 9t + 105 = 3 \cdot 2^3(3 \cdot 7 \cdot 11 \cdot 22 \cdot 85^4 e^2 + 3 \cdot 33 \cdot 22 \cdot 85e + 61);$$

$$H_{2,6} = 12t + 106 = 2(6t + 53);$$

$$H_{1,6} = 15t + 109 = 2(15 \cdot 14f^2 + 15 \cdot 66f + 1187);$$

$$H_{0,6} = 18t + 114 = 3(6t + 38).$$

It is easy to verify that for each $z = 1, 2, \dots, 6$, $H_{z,6}$ does not satisfy the condition in Theorem 2; that is, eq. (1) has no integral solution. For $H_{0,6}$, if eq. (1) has an integral solution, then, by Theorem 2, $6t + 38$ can be expressed as the form $s^2 + 3m^2$. Note that 3 is not a divisor of $(6t + 38)$, but 2 is. By Corollary of Theorem 1, we should have $6t + 38 \equiv 4 \pmod{6}$. But $6t + 38 \equiv 2 \pmod{6}$, a contradiction.

To sum up, $\{n(t(f(e))) : e \in Z\}$ contains no k -tight optimal DLN for each $k = 0, 1, 2, 3$.

Theorem 4. Let $n(t) = 3t^2 + 6t - 26$, $t = t(g) = 14812g^2 + 3036g + 151$, $s(g) = 14308392g^4 + 6176604g^3 + 984630g^2 + 68625g + 1764$. Then $\{G(n(t(g))); s(g)) : g = g(e) = 22 \cdot 85^2 e, e \in Z\}$ is an infinite family of 4-tight optimal DLNs, diameter $3t + 5$, and initial element $G(69283; 1764)$.

Proof. Let $n(t) = 3t^2 + 6t - 26$, $t(f) = 28f^2 + 132f + 151$, $f = 23g$, $g = 22 \cdot 85^2 e$. Then $t = 14812g^2 + 3036g + 151$. By Theorem 6, $\{n(t(g(e))) : e \in Z\}$ contains no k -tight optimal DLN for each $k = 0, 1, 2, 3$. To complete the proof, it suffices to show that it contains a 4-tight optimal DLN. To this end, let $z = x - y = 1$ and $j = 7$. By (1) we have

$$(a + b - 6)(a + b - 7) - ab - 7t - 26 = 0. \quad (5)$$

We find that $(a, b) = (1, -322g - 27)$ is a solution of (5). It follows that

$$l(g) = 2t + a = 29624g^2 + 6072g + 303;$$

$$h(g) = 2t + b = 29624g^2 + 5750g + 275;$$

$$x(g) = t + a + b - 7 = 14812g^2 + 2714g + 119;$$

$$y(g) = x - 1 = 14812g^2 + 2714g + 118;$$

$$h'(g) = h - y = 14182g^2 + 3036g + 157;$$

$$l'(g) = l - x = 14182g^2 + 3358g + 184.$$

Choose $\alpha(g) = 322g^2 + 73g + 4$, $\beta(g) = -322g^2 - 66g - 3$. Then $\alpha(g)y(g) + \beta(g)h'(g) = 1$ and $\text{g.c.d.}(y(g), h'(g)) = 1$ for $g = 22 \cdot 85^2e$, $e \in \mathbb{Z}$. By Lemma 3, the L -tile $L(n; l, h, x, y)$ is $(1, s)$ -realizable, where $s(g) = \alpha(g)l(g) - \beta(g)l'(g) = 14308392g^4 + 6176604g^3 + 984630g^2 + 68625g + 1764$.

Corollary. $G(69283; 1764)$ is a 4-tight optimal DLN; its diameter is 458.

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