

On Edge Addition of Altered Graphs^{*}

NAJIM Alaa A, XU Jun-ming

(Department of Mathematics, University of Science and Technology of China, Hefei 230026, China)

Abstract: This paper proves that, for any integers $t \geq 6$ and $d \geq 2$, the upper bound of minimum diameter of a connected graph obtained from a single path of length d by adding t extra edges is $\left\lceil \frac{d-2}{t+1} \right\rceil + 2$ for $d \in I'(t, k) = \{2k(t+1) + 1, 2k(t+1) + 2, 2k(t+1) - t + 1\} \cup \{2k(t+1) - t + h : h = 6, 7, \dots, t\}$ and $\left\lceil \frac{d-2}{t+1} \right\rceil + 1$ otherwise for any integer $k \geq 1$, which improves the known results. The bound $\left\lceil \frac{d-2}{t+1} \right\rceil + 1$ is best possible.

Key words: diameter; altered graph; edge addition

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0 Introduction

We follow ref. [1] for graph-theoretical terminology and notation not defined here. Let $G = (V, E)$ be a simple undirected graph, where $V = V(G)$ and $E = E(G)$ are the vertex-set and the edge-set of G , respectively. Let t_d denote the maximum number of edges that can be added to a path of length d , and $P(t, d)$ the minimum diameter of a graph obtained by adding t extra edges to a path of length d . Clearly, the symbol $P(t, d)$ means $t \leq t_d$. Determining $P(t, d)$ for given d and t , proposed by ref. [2], is of important interest in designing and analyzing interconnection networks.

For some small t 's and special d 's, the values of $P(t, d)$ have been determined. It is easy to verify that $P(1, d) = \left\lceil \frac{d+1}{2} \right\rceil$ for $d \geq 2$; Schoone *et al.* [3] determined $P(2, d) = \left\lceil \frac{d+1}{3} \right\rceil$ for $d \geq 3$ and $P(3, d) = \left\lceil \frac{d+2}{4} \right\rceil$ for $d \geq 5$; Deng and Xu [4] determined

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Biography: NAJIM Alaa A, male, born in 1965, PhD candidate. Research field: graphs and combinatorics.

E-mail: alaaamer6@hotmail.com

$P(t, (2k-1)(t+1)+1) = 2k$ for any positive integer k , $\left\lfloor \frac{d}{t+1} \right\rfloor \leq P(t, d) \leq \left\lceil \frac{d}{t+1} \right\rceil + 1$ for $t = 4, 5$ and $d \geq 4$, and, in general, $\left\lfloor \frac{d}{t+1} \right\rfloor \leq P(t, d) \leq \left\lfloor \frac{d-2}{t+1} \right\rfloor + 3$. For any integer $k \geq 1$, let

$$I'(t, k) = \{2k(t+1)+1, 2k(t+1)+2, 2k(t+1)-t+1\} \cup \{2k(t+1)-t+h : h = 6, 7, \dots, t\}.$$

In this paper, we improve the upper bound of $P(t, d)$ to

$$P(t, d) \leq \begin{cases} \left\lfloor \frac{d-2}{t+1} \right\rfloor + 2 & \text{if } d \in I'(t, k) \text{ for some } t \text{ and } k, \\ \left\lfloor \frac{d-2}{t+1} \right\rfloor + 1 & \text{otherwise} \end{cases}$$

for the integers $t \geq 6, d \geq 2$ and $k \geq 1$. This upper bound is tight for some t and d .

1 Several lemmas

Lemma 1.1^[3] For any integers k and h with $2 \leq h \leq t+1$,

$$P(t, (t+1)(2k-1)+h) \leq 2k+1, \text{ when } t = 4, 5.$$

The following lemma is simple, but useful.

Lemma 1.2 Let t_d be the maximum number of edges added to a path of length $d > 1$, then

$$t_d = \frac{d(d-1)}{2}.$$

Lemma 1.3^① $P(t, d) \leq 2k$, where $d \geq 2, d-2 \leq t < t_d$ for $k = 1$, and $\left\lfloor \frac{d-1}{2k-1} \right\rfloor - 1 \leq t \leq \left\lfloor \frac{d}{2k-2} \right\rfloor - 2$ for any integer $k \geq 2$.

Proof To prove the lemma, we construct an altered graph G from a single path $P = x_1 x_2 \cdots x_{d+1}$ by adding t extra edges such that the diameter of G is at most $2k$.

When $k = 1$, choose $r = 1$ or 2 . We add t extra edges $x_r x_j$ to P , where $j = r+2, r+3, \dots, d+1$. Since every vertex can reach the vertex x_r within one step, then the distance between any two vertices is at most 2. Thus $P(t, d) \leq 2$ for $t = d-2$. Since $P(t, d) \geq P(t', d)$ if $t \leq t' < t_d$, so $P(t, d) \leq 2$ when $d-2 \leq t < t_d$.

Now, assume $k > 1$. Put $m = \left\lfloor \frac{d-1}{2k-1} \right\rfloor - 1$. We add m edges $x_{k+1} x_{d-(k-2)-i(2k-1)}$ for $i = 0, 1, \dots, m-1$. The end-vertices of these edges divide P into $m+2$ segments $L_1, L_2, \dots, L_{m+1}, L_{m+2}$, where

① Clearly, the lemma is invalid for some d and k when $\left\lfloor \frac{d-1}{2k-1} \right\rfloor - 1 > \left\lfloor \frac{d}{2k-2} \right\rfloor - 2$ i. e. $2k+1 \leq P(t, d)$ or $P(t, d) \leq 2k-1$ for any $t \geq 1$.

$$\begin{aligned}
 L_1 &= P(x_1, x_{k+1}), \\
 L_2 &= P(x_{k+1}, x_{d-(k-2)-(m-1)(2k-1)}), \\
 L_{m+1-i} &= P(x_{d-(k-2)-(i+1)(2k-1)}, x_{d-(k-2)-i(2k-1)}), \quad i = 0, 1, \dots, m-2, \\
 L_{m+2} &= P(x_{d-(k-2)}, x_{d+1}).
 \end{aligned}$$

See Fig. 1 for an example. Since $m = \left\lceil \frac{d-1}{2k-1} \right\rceil - 1 \geq \frac{d-1}{2k-1} - 1$, we have

$$\begin{aligned}
 d - (k-2) - (m-1)(2k-1) - (k+1) &\leq \\
 d - (k-2) - \left(\frac{d-1}{2k-1} - 2 \right) (2k-1) - (k+1) &= 2k.
 \end{aligned}$$

Indeed,

$$\begin{aligned}
 d(L_1) &= k, \\
 d(L_2) &\leq 2k, \\
 d(L_{m+1-i}) &= 2k-1, \quad i = 0, 1, \dots, m-2, \text{ if } m \geq 2, \\
 d(L_{m+2}) &= k-1.
 \end{aligned}$$

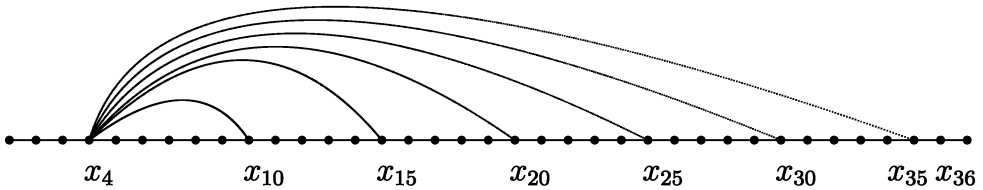


Fig. 1 Construction of Lemma 1. 3 for $k = 3, d = 36$ and $m = t = 6$

We need to prove that the diameter of G is at most $2k$. Since any vertex can reach the vertex x_{k+1} within k steps, the distance of any two vertices in G is at most $2k$, which means $P(m, d) \leq 2k$. Since $t \geq m$ and $P(t, d) \leq P(t', d)$ if $t \geq t'$, so $P(t, d) \leq 2k$ when $d \geq 2$ and $\left\lceil \frac{d-1}{2k-1} \right\rceil - 1 \leq t \leq \left\lceil \frac{d}{2k-2} \right\rceil - 2$ for any integer $k \geq 2$. □

Lemma 1. 4 Lemma 1. 3 is equivalent to the statement that $P(t, d) \leq 2k$, where $\frac{1 + \sqrt{9 + 8t}}{2} \leq d \leq t + 2$ for $k = 1$, and $d = 2k(t + 1) - t - h, h = 0, 1, \dots, t + 1$ for any integer $k \geq 2$.

Proof Clearly, the lemma satisfies when $k = 1$, since $d - 2 \leq t < t_d$ from lemma 1. 3. So we only need to prove the lemma for $k \geq 2$. Let $t = \left\lceil \frac{d-1}{2k-1} \right\rceil - 1$, then $t \geq \frac{d-1-(2k-1)}{2k-1}$, that is, $d \leq 2k(t+1) - t$. Now let $t \leq \left\lceil \frac{d}{2k-1} \right\rceil - 2$, we prove $d \geq 2k(t+1) - 2t - 1$ by contradiction. Suppose $d \leq 2k(t+1) - 2t - 2$. Then $d \leq (2k-2) \cdot (t+1)$, which implies $t \geq \left\lceil \frac{d}{2k-2} \right\rceil - 1$, a contradiction. Thus $d \geq 2k(t+1) - 2t - 1$. □

The following lemma extends Lemma 1. 1 for any $t \geq 6$.

Lemma 1. 5 $P(t, 2k(t+1) - t + 6 - h') \leq 2k + 1$ for the integers $t \geq 6, k \geq 1$, and

$$1 \leq h' \leq 5.$$

Proof Like the proof of Lemma 1.3, we construct an altered graph G from a single path $P = x_1 x_2 \cdots x_{d+1}$ by adding t extra edges such that the diameter of G is at most $2k+1$.

Let $d = 2k(t+1) - t + 5$. We add t edges

$$e_{2i-1} = x_{2k(3-r)+(1-r)} x_{d-(i-1)(2k-1)+1}, \quad \text{for } i = 1, 2, \dots, (t-r)/2 - 2$$

$$e_{2i} = x_{2k(2+r)+r} x_{d-(2i-1)(2k-1)-1}, \quad \text{for } i = 1, 2, \dots, (t+r)/2 - 3$$

$$e_{t-4} = x_{4k} x_{12k+1}, \quad e_{t-3} = x_{6k+1} x_{10k+1},$$

$$e_{t-2} = x_{4k} x_{8k+1}, \quad e_{t-1} = x_{2k} x_{6k+1}, \quad e_t = x_1 x_{4k},$$

$$\text{where } r = \begin{cases} 0 & \text{if } t \text{ is even,} \\ 1 & \text{if } t \text{ is odd.} \end{cases}$$

Now the end-vertices of these edges divide P into $t+1$ segments, since

$$d - (t-5)(2k-1) + 1 = 2k(t+1) - t + 5 - (t-5)(2k-1) + 1 = 12k + 1,$$

we have

$$L_i = P(x_{d-i(2k-1)+1}, x_{d-(i-1)(2k-1)+1}), \quad \text{for } i = 1, 2, \dots, t-5$$

$$L_{t-4} = P(x_{10k+1}, x_{12k+1}), \quad L_{t-3} = P(x_{8k+1}, x_{10k+1}),$$

$$L_{t-2} = P(x_{6k+1}, x_{8k+1}), \quad L_{t-1} = P(x_{4k}, x_{6k+1}),$$

$$L_t = P(x_{2k}, x_{4k}), \quad L_{t+1} = P(x_1, x_{2k}),$$

See Fig. 2 for an example.

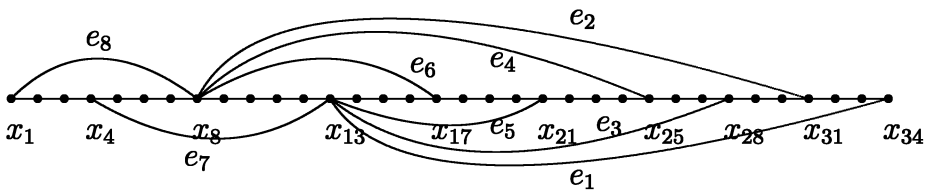


Fig. 2 Construction of Lemma 1.5 for $k = 2, t = 8$ and $d = 33$.

Let $X = \{x_1, x_2, \dots, x_{12k+1}\}$ and $X' = \{x_{2k+1}, x_{2k+2}, \dots, x_d\}$. From the proof of Lemma 1.1 we have the distance between any two vertices $x, y \in X$ is less than or equal to $2k+1$. So we only need to prove the lemma when $x, y \in X'$ or $x \in X$ and $y \in X'$ is less than or equal to $2k+1$. For $j = 1, 2, \dots, t-5$ we define $\left(\frac{t(t+1)}{2} - 15\right)$ cycles $C_j^1, C_j^2, \dots, C_j^{t-j+1}$ as

$$C_j^1 = L_j \cup L_{j+1} + e_j + e_{j+2},$$

$$C_j^i = L_j \cup L_{i+j} + e_j + e_{j+1} + e_{i+j} + e_{i+j+1}, \text{ for } i = 2, 3, \dots, t-j-3,$$

$$C_j^{t-j-2} = L_j \cup L_{t-2} + e_j + e_{j+1} + e_{t-2},$$

$$C_j^{t-j-1} = L_j \cup L_{t-1} + e_j + e_{j+1},$$

$$C_j^{t-j} = L_j \cup L_t + e_j + e_{j+1} + e_{t-1},$$

$$C_j^{t-j+1} = L_j \cup L_{t+1} + e_j + e_{j+1} + e_{t-1} + e_t.$$

Their lengths are

$$\varepsilon(C_j^1) = 4k, \quad \text{if } j < t-5;$$

$$\begin{aligned} \epsilon(C_{t-5}^1) &= 4k + 1; \\ \epsilon(C_j^i) &= 4k + 2, \quad \text{for } i = 2, 3, \dots, t - j - 5, j < t - 6, t > 7; \\ \epsilon(C_j^{\alpha_j-4}) &= 4k + 3, \quad \text{if } j < t - 5, t \neq 6; \\ \epsilon(C_j^{\alpha_j-3}) &= 4k + 3; \\ \epsilon(C_j^i) &= 4k + 2, \text{ for } t - j - 2 \leq i \leq t - j + 1. \end{aligned}$$

It is easy to see that any two vertices $x, y \in X'$ or $x \in X$ and $y \in X'$ are contained in some cycle C_j^i defined above. The fact

$$\max\{d(C_j^i) : 1 \leq j \leq t - 5, 1 \leq i \leq t - j + 1\} \leq \left\lfloor \frac{4k + 3}{2} \right\rfloor = 2k + 1$$

means

$$P(t, 2k(t + 1) - t + 5) \leq d(G) \leq 2k + 1.$$

Since $P(t, d) \leq P(t, d')$ if $d \leq d', P(t, 2k(t + 1) - t + 6 - h') \leq 2k + 1$ when $1 \leq h' \leq 5$. □

2 Proof of main result

In this section, we only consider $t \geq 6$. For a given t , let $d(k) = 2k(t + 1) - 2t - 1$. Then $d(k + 1) = 2k(t + 1) + 1 = d(k) + 2t + 2$. We give an upper bound of $P(t, d)$ when d is any integer in the interval $I(t, k) = [2k(t + 1) - 2t - 1, 2k(t + 1)]$ for any $k \geq 1$ and $t \geq 6$. To state our theorem, let $I'(t, k) = \{2k(t + 1) + 1, 2k(t + 1) + 2, 2k(t + 1) - t + 1\} \cup \{2k(t + 1) - t + h : h = 6, 7, \dots, t\}$.

Theorem 2.1 For any $t \geq 6$ and $k \geq 1$, if $d \in I(t, k)$ then

$$P(t, d) \leq \begin{cases} \left\lceil \frac{d-2}{t+1} \right\rceil + 2 & \text{if } d \in I'(t, k) \text{ for some } t \text{ and } k; \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{otherwise.} \end{cases}$$

Proof First, we consider $k = 1$. In this case, $\frac{1 + \sqrt{9 + 8t}}{2} \leq d \leq t + 2$ and $d \geq 3$ since $t \geq 1$. For $d \leq t + 2$ we have $d = t + 2 - i$ for some $i = 0, 1, \dots, t - 1$ and, thus,

$$\frac{d-2}{t+1} = \frac{(t+2-i)-2}{t+1} = \frac{t-i}{t+1},$$

which means $\left\lceil \frac{d-2}{t+1} \right\rceil = 1$. From Lemma 1.4 we have

$$P(t, d) \leq 2 = \left\lceil \frac{d-2}{t+1} \right\rceil + 1.$$

Secondly, we consider $k \geq 2$ and $d = 2k(t + 1) - t - h, h = 0, 1, \dots, t + 1$. In this case, we have

$$\frac{d-2}{t+1} = \frac{2k(t+1) - t - h - 2}{t+1} = 2k - 1 - \frac{h+1}{t+1},$$

which implies that

$$2k = \frac{d-2}{t+1} + 1 + \frac{h+1}{t+1}. \tag{1}$$

For $h = 0, 1, \dots, t$, we have $\frac{h+1}{t+1} \leq 1$. When $h = t$ the equality holds and, hence,

$$2k - 2 = \frac{d-2}{t+1} = \left\lceil \frac{d-2}{t+1} \right\rceil. \quad (2)$$

Also when $h = t+1, d = 2k(t+1) - 2t - 1$. From (1) we have

$$2k - 2 = \frac{d-2}{t+1} + \frac{1}{t+1} = \left\lceil \frac{d-2}{t+1} \right\rceil. \quad (3)$$

If $0 \leq h \leq t-1$, then from (1) we have

$$2k - 1 = \frac{d-2}{t+1} + \frac{h+1}{t+1} = \left\lceil \frac{d-2}{t+1} \right\rceil. \quad (4)$$

Thus, from Lemma 1.4 and (1) we have

$$P(t, 2k(t+1) - t - h) \leq 2k = \frac{d-2}{t+1} + 1 + \frac{h+1}{t+1}. \quad (5)$$

From (2), (3), (4) and (5) we have

$$P(t, d) \leq \begin{cases} \left\lceil \frac{d-2}{t+1} \right\rceil + 2 & \text{for } d = 2k(t+1) - 2t - 1 \text{ or } 2k(t+1) - 2t; \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{for } 2k(t+1) - 2t + 1 \leq d \leq 2k(t+1) - t. \end{cases}$$

which implies for $k \geq 1$

$$P(t, d) \leq \begin{cases} \left\lceil \frac{d-2}{t+1} \right\rceil + 2 & \text{for } d = 2k(t+1) + 1 \text{ or } 2k(t+1) + 2; \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{for } 2k(t+1) + 3 \leq d \leq 2k(t+1) + t + 2. \end{cases}$$

Thirdly, for $k \geq 1$ we consider $d = 2k(t+1) - t + 6 - h'$ with $1 \leq h' \leq 5$. In this case, we have

$$\frac{d-2}{t+1} = \frac{2k(t+1) - t + 6 - h' - 2}{t+1} = 2k - \frac{t-4+h'}{t+1},$$

that is,

$$2k + 1 = \frac{d-2}{t+1} + \frac{t-4+h'}{t+1} + 1. \quad (6)$$

If $h' = 5$ then, from (6), we have

$$2k + 1 = \frac{d-2}{t+1} + 2 = \left\lceil \frac{d-2}{t+1} \right\rceil + 2. \quad (7)$$

If $1 \leq h' \leq 4$ then, from (6), we have

$$2k + 1 = \frac{d-2}{t+1} + \frac{t-4+h'}{t+1} + 1 = \left\lceil \frac{d-2}{t+1} \right\rceil + 1. \quad (8)$$

From Lemma 1.5 and the equalities (6), (7) and (8) we have

$$P(t, d) \leq \begin{cases} \left\lceil \frac{d-2}{t+1} \right\rceil + 2 & \text{for } d = 2k(t+1) - t + 1; \\ \left\lceil \frac{d-2}{t+1} \right\rceil + 1 & \text{for } 2k(t+1) - t + 2 \leq d \leq 2k(t+1) - t + 5. \end{cases}$$

Finally, since for any $k \geq 1$,

$$P(t, 2k(t+1) + 1) \leq \left\lceil \frac{d-2}{t+1} \right\rceil + 2.$$

and $P(t, d') \leq P(t, d'')$ when $d' \leq d''$, then we have

$$P(t, d) \leq \left\lceil \frac{d-2}{t+1} \right\rceil + 2, \quad \text{for } 2k(t+1) - t + 6 \leq d \leq 2k(t+1). \quad \square$$

Corollary 2.1^[4] $P(t, 2k(t+1) - t) = 2k$ for any positive integer k .

Proof Let $d = 2k(t+1) - t$. On the one hand, by $P(t, d) \geq \left\lfloor \frac{d}{t+1} \right\rfloor$, due to ref. [4]

and stated in Introduction, we have

$$P(t, d) \geq \left\lfloor \frac{d}{t+1} \right\rfloor = \left\lfloor \frac{2k(t+1) - t}{t+1} \right\rfloor = 2k.$$

On the other hand, $(t+1) \nmid (d-2)$, by Theorem 2.1, we have

$$P(t, d) \leq \left\lceil \frac{d-2}{t+1} \right\rceil + 1 = \left\lceil \frac{2k(t+1) - t - 2}{t+1} \right\rceil + 1 = 2k.$$

So, $P(t, 2k(t+1) - t) = 2k$ for any positive integer k . □

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变更图的边添加

NAJIM Alaa A, 徐俊明

(中国科学技术大学数学系, 安徽合肥 230026)

摘要:证明了: 对任何整数 $t \geq 6$ 和 $d \geq 2$, 从一条长为 d 的简单路通过添加 t 条边后得到的图的最小直径上界为 $\left\lceil \frac{d-2}{t+1} \right\rceil + 2$, 如果 $d \in I'(t, k) = \{2k(t+1) + 1, 2k(t+1) + 2, 2k(t+1) - t + 1\} \cup \{2k(t+1) - t + h : h = 6, 7, \dots, t\}$; 其他情形为 $\left\lceil \frac{d-2}{t+1} \right\rceil + 1$. 这个证明改进了已知结果, 而且 $\left\lceil \frac{d-2}{t+1} \right\rceil + 1$ 是最好的上界.

关键词: 直径; 变更图; 边添加