

Domination Numbers of Undirected Toroidal Mesh

Xin XIE

Department of Mathematics, Huangshan University, Huangshan 245041, P. R. China
E-mail: xiexin@hsu.edu.cn

Jun Ming XU

*Department of Mathematics, University of Science and Technology of China,
Hefei 230026, P. R. China*
E-mail: xujm@ustc.edu.cn

Abstract The (d, m) -domination number $\gamma_{d,m}$ is a new measure to characterize the reliability of resources-sharing in fault tolerant networks, in some sense, which can more accurately characterize the reliability of networks than the m -diameter does. In this paper, we study the $(d, 4)$ -domination numbers of undirected toroidal mesh $C_{d_1} \times C_{d_2}$ for some special values of d , obtain that $\gamma_{d,4}(C_{d_1} \times C_3) = 2$ if and only if $d_4(G) - e_1 \leq d < d_4(G)$ for $d_1 \geq 5$, $\gamma_{d,4}(C_{d_1} \times C_4) = 2$ if $d_4(G) - (2e_1 - \lfloor \frac{d_1+e_1}{2} \rfloor) \leq d < d_4(G)$ for $d_1 \geq 24$, and $\gamma_{d,4}(C_{d_1} \times C_{d_2}) = 2$ if $d_4(G) - (e_1 - 2) \leq d < d_4(G)$ for $d_1 = d_2 \geq 14$.

Keywords Undirected toroidal mesh, reliability, m -diameter, domination number

MR(2000) Subject Classification 05C40, 68M10, 68M15, 68R10

1 Introduction

We quote from [1] the terminology and notations not defined here, use graphs to represent networks and denote the length of a path P by $|P|$.

The n -dimensional undirected toroidal mesh, denoted by $C(d_1, d_2, \dots, d_n)$, has the vertex-set $\{(x_1, \dots, x_n) | 0 \leq x_i < d_i (i = 1, 2, \dots, n)\}$. Each vertex (x_1, x_2, \dots, x_n) is adjacent to $2n$ other vertices: $(x_1 \pm 1, x_2, \dots, x_n)$, $(x_1, x_2 \pm 1, \dots, x_n)$, \dots , $(x_1, x_2, \dots, x_n \pm 1)$, where additions are performed modulo $d_i (1 \leq i \leq n)$. It is well known that $C(d_1, d_2, \dots, d_n)$ is $2n$ -regular and vertex-transitive. Its connectivity is $2n$. The toroidal mesh is widely used in network theory (see [2–4]).

In order to characterize the reliability of transmission delay in a real-time parallel processing system, Hsu and Lyuu [5] and Flandrin and Li [6] independently introduced m -diameter. For an m -connected graph G , the distance with width m from vertex x to y , denoted by $d_m(G; x, y)$, is the minimum number d for which there are m internally disjoint (x, y) -paths in G of length at most d . The diameter with width m , denoted by $d_m(G)$, is the maximum of $d_m(G; x, y)$ over all pairs (x, y) of vertices of G .

Received April 28, 2010, accepted July 23, 2010

Supported by National Natural Science Foundation of China (Grant No. 11071233) and Foundation of Huangshan University (Grant No. 2011xkj012)

Recently, Li and Xu in [7] defined a new parameter (d, m) -domination number in m -connected graphs, in some sense, which can more accurately characterize the reliability of networks than the wide-diameter does.

Definition For an m -connected graph G and a given integer d , a nonempty and proper subset S of the vertex set of G is called a (d, m) -dominating set of G if for any vertex x of G but not in S there are at least m internally disjoint (x, S) -paths of length at most d . Denote a set of all (d, m) -dominating sets of G by $S_{d,m}(G)$. The parameter

$$\gamma_{d,m}(G) = \min\{|S| : S \in S_{d,m}(G)\}$$

is called the (d, m) -domination number of G .

This notion not only generalizes that of the classical domination numbers of a graph but also gives a good measure of the problem of resources-sharing in fault tolerant networks. An important and practical problem is how to choose a (d, m) -dominating set S such that the number of vertices in S is as small as possible. Thus the (d, m) -domination numbers in conjunction with other well-known parameters, for example, the (d, m) -independence numbers (see [8]), can provide a more accurate analysis of fault tolerance for reliability and efficiency of networks of parallel architectures.

In general, to determine the (d, m) -domination number of a graph is NP-Complete since its special case of $(1, 1)$, the domination number of the graph, is NP-Complete (see [9]). Thus it is of interest to determine the (d, m) -domination number of some well-known networks for some special values of m and d . The $(1, 1)$ -domination numbers of some graphs have been studied (see, for example, [10]).

Of course, if $d \geq d_m(G)$, then we obtain that $\gamma_{d,m}(G) = 1$. And we have $\gamma_{d,m}(G) \geq 2$ for $d < d_m(G)$. So it is of interest to determine values of $\gamma_{d,m}(G)$ for $d < d_m(G)$. Lu and Zhang [11] proved the $(d, 2n)$ -domination number of $C(d_1, d_2, \dots, d_n)$ ($\neq C(3, 3, \dots, 3)$) is 2 for $d = \text{diam}(C(d_1, d_2, \dots, d_n))$ ($n \geq 3, d_i \geq 3, i \in \{0, 1, \dots, n\}$). Xie and Xu [12] proved the $(d, 2n)$ -domination number of $C(3, 3, \dots, 3)$ is 3 for $d = \text{diam}(C(3, 3, \dots, 3)) = n$. This paper obtains $\gamma_{d,4}(C_{d_1} \times C_3) = 2$ if and only if $d_4(G) - e_1 \leq d < d_4(G)$ for $d_1 \geq 5$, $\gamma_{d,4}(C_{d_1} \times C_4) = 2$ if $d_4(G) - (2e_1 - \lfloor \frac{d_1 + e_1}{2} \rfloor) \leq d < d_4(G)$ for $d_1 \geq 24$, and $\gamma_{d,4}(C_{d_1} \times C_{d_2}) = 2$ if $d_4(G) - (e_1 - 2) \leq d < d_4(G)$ for $d_1 = d_2 \geq 14$.

2 Preliminaries

The undirected toroidal mesh $C(d_1, d_2, \dots, d_n)$ can also be defined as the cartesian products $C_{d_1} \times C_{d_2} \times \dots \times C_{d_n}$, where C_{d_i} is an undirected cycle for each $i = 1, 2, \dots, n$. This graph is an n -dimensional orthogonal mesh with global edges. For $n = 2$, $C_{d_1} \times C_{d_2}$ is a mesh on a toroidal surface.

In this paper, we denote $\lfloor \frac{d_i}{2} \rfloor$ and $\lceil \frac{d_i}{2} \rceil$ by e_i and e'_i , define $\delta(x_i) = -1$ for $e_i < x_i \leq d_i - 1$ and $\delta(x_i) = 1$ for $0 < x_i \leq e_i$, respectively, where $i = 1, 2$.

Lemma 2.1 ([13]) For any integers $n \geq 2, d_i \geq 3 (1 \leq i \leq n)$,

$$\text{diam}(C_{d_1} \times C_{d_2} \times \dots \times C_{d_n}) = \sum_{i=1}^n \left\lceil \frac{d_i}{2} \right\rceil.$$

We will use another lemma from [13] that relates to the m -diameter of a general toroidal mesh.

Lemma 2.2 ([13]) *Let $G = C_{d_1} \times C_{d_2} \times \dots \times C_{d_n}$ for any $n \geq 2, d_1 \geq d_2 \geq \dots \geq d_n \geq 3$. Then*

$$d_{2n}(G) = \begin{cases} 2m + 1 = 2\text{diam}(G) - 1, & G = C_{2m+1} \times C_3, m \geq 2, \\ 2m + 2 = 2\text{diam}(G) - 2, & G = C_{2m+2} \times C_3, m \geq 2, \\ e_1 + 4 = \text{diam}(G) + 2, & G = C_d \times C_4, d \geq 9 \text{ or} \\ & G = C_d \times C_5, 9 \leq d \leq 13, \\ 1 + \sum_{i=1}^n e_i = \text{diam}(G) + 1, & \text{otherwise.} \end{cases}$$

3 Main Results

Theorem 3.1 *Let $G = C_{d_1} \times C_3, d_1 \geq 5$. We have $\gamma_{d,4}(G) = 2$ if and only if $d_4(G) - e_1 \leq d < d_4(G)$.*

Proof We easily know $d_4(G) - e_1 = d_1 - e_1 = e'_1 (\geq e_1)$ by Lemma 2.2.

First we prove $\gamma_{d,4}(G) = 2$ if $d_4(G) - e_1 \leq d < d_4(G)$. From the definition of the (d, m) -domination number, we have $\gamma_{d,4}(G) \geq 2$ for $d < d_4(G)$ and $\gamma_{d,4}(G) \leq \gamma_{d-1,4}(G)$, so we only need to prove $\gamma_{d,4}(G) \leq 2$ for $d = e'_1$. Let $S = \{o, e\}$, where $o = (0, 0), e = (e_1, 1)$. For any vertex $x = (x_1, x_2) \in V(G) - S, x_1 \in \{0, 1, \dots, d_1 - 1\}, x_2 \in \{0, 1, 2\}$, we will prove S can $(d, 4)$ -dominate vertex x , we only need to consider the following cases by vertex-transitivity and symmetry:

Case 1 $x_2 = 0, x = (x_1, 0) (\neq o)$.

Subcase 1a $1 \leq x_1 \leq e_1 - 2$ or $e_1 + 2 \leq x_1 \leq d_1 - 1$.

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

$$P_1 : x = (x_1, 0) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow \dots \rightarrow (\delta(x_1), 0) \rightarrow o;$$

$$P_2 : x = (x_1, 0) \rightarrow (x_1, 1) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow \dots \rightarrow (\delta(x_1), 1) \rightarrow (0, 1) \rightarrow o;$$

$$P_3 : x = (x_1, 0) \rightarrow (x_1, 2) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow \dots \rightarrow (\delta(x_1), 2) \rightarrow (0, 2) \rightarrow o;$$

$$P_4 : x = (x_1, 0) \rightarrow (x_1 + \delta(x_1), 0) \rightarrow \dots \rightarrow (e_1, 0) \rightarrow e.$$

We can see $|P_1| = x_1, |P_2| = |P_3| = x_1 + 2 \leq e_1$ and $|P_4| = e_1 - x_1 + 1 \leq e_1$ if $1 \leq x_1 \leq e_1 - 2$. When $e_1 + 2 \leq x_1 \leq d_1 - 1$, we have $|P_1| = d_1 - x_1, |P_2| = |P_3| = d_1 - x_1 + 2 \leq d_1 - e_1 = e'_1$ and $|P_4| = x_1 - e_1 + 1 \leq d_1 - e_1 = e'_1$.

Subcase 1b $x_1 = e_1 - 1$ or $x_1 = e_1 + 1$.

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

$$P_1 : x = (e_1 - \delta(x_1), 0) \rightarrow (e_1 - 2\delta(x_1), 0) \rightarrow \dots \rightarrow (\delta(x_1), 0) \rightarrow o;$$

$$P_2 : x = (e_1 - \delta(x_1), 0) \rightarrow (e_1, 0) \rightarrow e;$$

$$P_3 : x = (e_1 - \delta(x_1), 0) \rightarrow (e_1 - \delta(x_1), 1) \rightarrow e;$$

$$P_4 : x = (e_1 - \delta(x_1), 0) \rightarrow (e_1 - \delta(x_1), 2) \rightarrow (e_1, 2) \rightarrow e.$$

The length of any path of the four is at most $\max\{x_1, d_1 - x_1\} = d_1 - e_1 - 1 = e'_1 - 1$.

Subcase 1c $x_1 = e_1$.

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

$$P_1 : x = (e_1, 0) \rightarrow e;$$

- $P_2 : x = (e_1, 0) \rightarrow (e_1, 2) \rightarrow e;$
- $P_3 : x = (e_1, 0) \rightarrow (e_1 + 1, 0) \rightarrow (e_1 + 1, 1) \rightarrow e;$
- $P_4 : x = (e_1, 0) \rightarrow (e_1 - 1, 0) \rightarrow (e_1 - 1, 1) \rightarrow e.$

Of course, the lengths of paths are at most $3 (\leq e'_1)$.

Case 2 $x_2 = 1, x = (x_1, 1) (\neq e).$

It is similar to Case 1 by vertex-transitivity.

Case 3 $x_2 = 2, x = (x_1, 2).$

Subcase 3a $x_1 = 0.$

We can construct four paths between vertices o and x in the same way as Subcase 1c by vertex-transitivity.

Subcase 3b $1 \leq x_1 \leq d_1 - 1$ and $x_1 \neq e_1.$

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

- $P_1 : x = (x_1, 2) \rightarrow (x_1, 0) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow \dots \rightarrow (\delta(x_1), 0) \rightarrow o;$
- $P_2 : x = (x_1, 2) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow \dots \rightarrow (\delta(x_1), 2) \rightarrow (0, 2) \rightarrow o;$
- $P_3 : x = (x_1, 2) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow \dots \rightarrow (e_1, 2) \rightarrow e;$
- $P_4 : x = (x_1, 2) \rightarrow (x_1, 1) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow \dots \rightarrow e.$

We can see $|P_1| = |P_2| = x_1 + 1 \leq e_1$ and $|P_3| = |P_4| = e_1 - x_1 + 1 \leq e_1$ if $1 \leq x_1 \leq e_1 - 1$. When $e_1 + 1 \leq x_1 \leq d_1 - 1$, we have $|P_1| = |P_2| = d_1 - x_1 + 1 \leq d_1 - e_1 = e'_1$ and $|P_3| = |P_4| = x_1 - e_1 + 1 \leq d_1 - e_1 = e'_1$.

Subcase 3c $x_1 = e_1.$

It is similar to Subcase 1c by vertex-transitivity.

From above, we know there exist four internally disjoint (x, S) -paths of length at most e'_1 in G for any vertex x . So S is a $(d, 4)$ -dominating set of G , we have $\gamma_{d,4}(G) \leq 2$.

Next we prove $d_4(G) - e_1 \leq d < d_4(G)$ if $\gamma_{d,4}(G) = 2$. We only need to prove $\gamma_{d,4}(G) \geq 3$ if $d < d_4(G) - e_1 = e'_1$.

Assume $S = \{o, y\}$ is a $(d, 4)$ -dominating set of G , where $o = (0, 0), y = (y_1, y_2), 0 \leq y_1 \leq e_1$ and $0 \leq y_2 \leq 1$ without loss of generality.

If $y = (y_1, 0)$, then the distance between vertices y and $x = (y_1 + e'_1 - 1, 1)$ is $\min\{e_1 + 2, e'_1\} = e'_1 (> d)$. If there exist four internally disjoint paths of length at most d between vertices o and x , we can assume that P_1 contains the vertices $o, (d_1 - 1, 0)$ and x, P_2 contains $o, (0, 1)$ and x, P_3 contains $o, (0, 2)$ and x . So the path P_4 must contain vertices $o, (1, 0)$ and $(y_1 + e'_1 - 2, 1)$. Of course, the length of P_4 is greater than or equal to $y_1 + e'_1 \geq 1 + e'_1 (> d)$. It is a contradiction. S can not $(d, 4)$ -dominate the vertex x .

If $y = (y_1, 1)$, then the distance between vertices y and $x = (y_1 + e'_1 - 1, 2)$ is at least $e'_1 (> d)$. Similarly, we can easily know there do not exist four internally disjoint paths of length at most d between vertices o and $x = (y_1 + e'_1 - 1, 2), S$ can not $(d, 4)$ -dominate the vertex $(y_1 + e'_1 - 1, 2)$.

It is a contradiction. So we have $\gamma_{d,4}(G) \geq 3$ for $d < e'_1$.

Thus, we have proven that $\gamma_{d,4}(G) = 2$ if and only if $d_4(G) - e_1 \leq d < d_4(G)$. □

Theorem 3.2 *Let $G = C_{d_1} \times C_4, d_1 \geq 24$. We have $\gamma_{d,4}(G) = 2$ if $d_4(G) - (2e_1 - \lfloor \frac{d_1 + e_1}{2} \rfloor) \leq d < d_4(G)$.*

Proof By Lemma 2.2, we have $d_4(G) - (2e_1 - \lfloor \frac{d_1+e_1}{2} \rfloor) = d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$. We only need to prove $\gamma_{d,4}(G) \leq 2$ for $d = d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$.

Let $S = \{o, e\}$, where $o = (0, 0)$, $e = (e_1, 0)$. We will prove S can $(d, 4)$ -dominate any vertex $x = (x_1, x_2) \in V(G) - S$, where $x_1 \in \{0, 1, \dots, d_1 - 1\}$ and $x_2 \in \{0, 1, 2, 3\}$. We only consider the following cases by vertex-transitivity:

Case 1 $x_2 = 0, x = (x_1, 0) (\neq o, e)$.

Subcase 1a $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 4$ or $\lceil \frac{d_1+e_1}{2} \rceil + 4 \leq x_1 \leq d_1 - 1$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$$P_1 : x = (x_1, 0) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow \dots \rightarrow (\delta(x_1), 0) \rightarrow o; \quad (1)$$

$$P_2 : x = (x_1, 0) \rightarrow (x_1, 1) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow \dots \rightarrow (\delta(x_1), 1) \rightarrow (0, 1) \rightarrow o; \quad (2)$$

$$P_3 : x = (x_1, 0) \rightarrow (x_1, -1) \rightarrow (x_1 - \delta(x_1), -1) \rightarrow \dots \rightarrow (\delta(x_1), -1) \rightarrow (0, -1) \rightarrow o; \quad (3)$$

$$P_4 : x = (x_1, 0) \rightarrow (x_1 + \delta(x_1), 0) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow (x_1, 2) \rightarrow \dots \rightarrow (\delta(x_1), 2) \rightarrow (0, 2) \rightarrow (-\delta(x_1), 2) \rightarrow (-\delta(x_1), 1) \rightarrow (-\delta(x_1), 0) \rightarrow o. \quad (4)$$

We can see $|P_1| = x_1$, $|P_2| = |P_3| = x_1 + 2$ and $|P_4| = x_1 + 8 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ if $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 4$. When $\lfloor \frac{d_1+e_1}{2} \rfloor + 5 \leq x_1 \leq d_1 - 1$, we have $|P_1| = d_1 - x_1$, $|P_2| = |P_3| = d_1 - x_1 + 2$ and $|P_4| = d_1 - x_1 + 8 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$.

Subcase 1b $d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 3 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 2$ or $\lceil \frac{d_1+e_1}{2} \rceil - 2 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil + 3$.

We construct the three paths P_i ($1 \leq i \leq 3$) as (1), (2), (3) and the one other path P_4 as follows:

$$P_4 : x = (x_1, 0) \rightarrow (x_1 + \delta(x_1), 0) \rightarrow \dots \rightarrow (e_1, 0) = e. \quad (5)$$

We can see $|P_1| = x_1$, $|P_2| = |P_3| = x_1 + 2 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ and $|P_4| = e_1 - x_1 \leq e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil + 3$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 3 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 2$. When $\lceil \frac{d_1+e_1}{2} \rceil - 2 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil + 3$, we have $|P_1| = d_1 - x_1$, $|P_2| = |P_3| = d_1 - x_1 + 2 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ and $|P_4| = x_1 - e_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - e_1 + 3$.

Subcase 1c $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 3 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ or $\lceil \frac{d_1+e_1}{2} \rceil - 4 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 3$.

We construct the two paths P_1, P_2 as (1), (5) and the other two paths P_3, P_4 as follows:

$$P_3 : x = (x_1, 0) \rightarrow (x_1, 1) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow \dots \rightarrow (e_1, 1) \rightarrow e; \quad (6)$$

$$P_4 : x = (x_1, 0) \rightarrow (x_1, 3) \rightarrow (x_1 + \delta(x_1), 3) \rightarrow \dots \rightarrow (e_1, 3) \rightarrow e. \quad (7)$$

We can see $|P_1| = x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$, $|P_2| = e_1 - x_1$, $|P_3| = |P_4| = e_1 - x_1 + 2 \leq e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil - 1$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 3 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$. When $\lceil \frac{d_1+e_1}{2} \rceil - 4 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 3$, we have $|P_1| = d_1 - x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$, $|P_2| = x_1 - e_1$ and $|P_3| = |P_4| = x_1 - e_1 + 2 \leq \lceil \frac{d_1+e_1}{2} \rceil - e_1 - 1$.

Subcase 1d $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 5 \leq x_1 \leq e_1 - 1$ or $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 5$.

We construct the three paths P_1, P_2, P_3 as (5), (6), (7) and the one other path P_4 as follows:

$$P_4 : x = (x_1, 0) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow (x_1, 2) \rightarrow \dots \rightarrow (e_1 + \delta(x_1), 2) \rightarrow (e_1 + \delta(x_1), 1) \rightarrow (e_1 + \delta(x_1), 0) \rightarrow e.$$

We can see $|P_1| = e_1 - x_1$, $|P_2| = |P_3| = e_1 - x_1 + 2$, $|P_4| = e_1 - x_1 + 8 = e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil + 3$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 5 \leq x_1 \leq e_1 - 1$. When $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 5$, we have $|P_1| = x_1 - e_1$, $|P_2| = |P_3| = x_1 - e_1 + 2$, $|P_4| = x_1 - e_1 + 8 = \lceil \frac{d_1+e_1}{2} \rceil - e_1 + 3$.

Case 2 $x_2 = 1, x = (x_1, 1)$.

Subcase 2a $x = (0, 1)$ or $x = (e_1, 1)$.

We can easily construct four (x, S) -paths of length at most 3 in the same way as Subcase 1c of Theorem 3.1.

Subcase 2b $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 1$ or $\lceil \frac{d_1+e_1}{2} \rceil + 1 \leq x_1 \leq d_1 - 1$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$$P_1 : x = (x_1, 1) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow \cdots \rightarrow (\delta(x_1), 1) \rightarrow (0, 1) \rightarrow o; \quad (8)$$

$$P_2 : x = (x_1, 1) \rightarrow (x_1, 0) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow \cdots \rightarrow (\delta(x_1), 0) \rightarrow o; \quad (9)$$

$$P_3 : x = (x_1, 1) \rightarrow (x_1, 2) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow \cdots \rightarrow (\delta(x_1), 2) \rightarrow (0, 2) \rightarrow (-\delta(x_1), 2) \rightarrow (-\delta(x_1), 1) \rightarrow (-\delta(x_1), 0) \rightarrow o;$$

$$P_4 : x = (x_1, 1) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow (x_1 + \delta(x_1), 3) \rightarrow (x_1, 3) \rightarrow \cdots \rightarrow (\delta(x_1), 3) \rightarrow (0, 3) \rightarrow o.$$

We can see $|P_1| = |P_2| = x_1 + 1$ and $|P_3| = |P_4| = x_1 + 5 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ if $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil - 1$. When $\lceil \frac{d_1+e_1}{2} \rceil + 1 \leq x_1 \leq d_1 - 1$, we have $|P_1| = |P_2| = d_1 - x_1 + 1$, $|P_3| = |P_4| = d_1 - x_1 + 5 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$.

Subcase 2c $d_1 - \lceil \frac{d_1+e_1}{2} \rceil \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 3$ or $\lceil \frac{d_1+e_1}{2} \rceil - 3 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil$.

We construct the two paths P_1, P_2 as (8), (9) and the other two paths P_3, P_4 as follows:

$$P_3 : x = (x_1, 1) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow \cdots \rightarrow (e_1, 1) \rightarrow e; \quad (10)$$

$$P_4 : x = (x_1, 1) \rightarrow (x_1, 2) \rightarrow (x_1, 3) \rightarrow (x_1 + \delta(x_1), 3) \rightarrow \cdots \rightarrow (e_1, 3) \rightarrow e.$$

We can see $|P_1| = |P_2| = x_1 + 1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$, $|P_3| = e_1 - x_1 + 1$, $|P_4| = e_1 - x_1 + 3 \leq e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil + 3$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 3$. When $\lceil \frac{d_1+e_1}{2} \rceil - 3 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil$, we have $|P_1| = |P_2| = d_1 - x_1 + 1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$, $|P_3| = x_1 - e_1 + 1$, $|P_4| = x_1 - e_1 + 3 = \lceil \frac{d_1+e_1}{2} \rceil - e_1 + 3$.

Subcase 2d $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4 \leq x_1 \leq e_1 - 1$ or $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 4$.

We construct the one path P_1 as (10) and the other three paths P_2, P_3, P_4 as follows:

$$P_2 : x = (x_1, 1) \rightarrow (x_1, 0) \rightarrow (x_1 + \delta(x_1), 0) \rightarrow \cdots \rightarrow e;$$

$$P_3 : x = (x_1, 1) \rightarrow (x_1, 2) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow \cdots \rightarrow (e_1 + \delta(x_1), 2) \rightarrow (e_1 + \delta(x_1), 1) \rightarrow (e_1 + \delta(x_1), 0) \rightarrow e;$$

$$P_4 : x = (x_1, 1) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow (x_1 - \delta(x_1), 3) \rightarrow (x_1, 3) \rightarrow \cdots \rightarrow (e_1, 3) \rightarrow e.$$

We can see $|P_1| = |P_2| = e_1 - x_1 + 1$, $|P_3| = |P_4| = e_1 - x_1 + 5 = e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil + 1$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4 \leq x_1 \leq e_1 - 1$. When $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 4$, we have $|P_1| = |P_2| = x_1 - e_1 + 1$, $|P_3| = |P_4| = x_1 - e_1 + 5 = \lceil \frac{d_1+e_1}{2} \rceil - e_1 + 1$.

Case 3 $x_2 = 2, x = (x_1, 2)$.

Subcase 3a $x = (0, 2)$ or $x = (e_1, 2)$.

We can easily construct four (x, S) -paths of length at most 4.

Subcase 3b $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil$ or $\lceil \frac{d_1+e_1}{2} \rceil \leq x_1 \leq d_1 - 1$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$$P_1 : x = (x_1, 2) \rightarrow (x_1, 1) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow \cdots \rightarrow (\delta(x_1), 1) \rightarrow (0, 1) \rightarrow o;$$

$$P_2 : x = (x_1, 2) \rightarrow (x_1, 3) \rightarrow (x_1 - \delta(x_1), 3) \rightarrow \cdots \rightarrow (\delta(x_1), 3) \rightarrow (0, 3) \rightarrow o;$$

$$P_3 : x = (x_1, 2) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow (x_1 + \delta(x_1), 0) \rightarrow (x_1, 0) \rightarrow \cdots \rightarrow (\delta(x_1), 0) \rightarrow o;$$

$P_4 : x = (x_1, 2) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow \cdots \rightarrow (\delta(x_1), 2) \rightarrow (0, 2) \rightarrow (-\delta(x_1), 2) \rightarrow (-\delta(x_1), 1) \rightarrow (-\delta(x_1), 0) \rightarrow o$.

We can see $|P_1| = |P_2| = x_1 + 2$ and $|P_3| = |P_4| = x_1 + 4 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$ if $1 \leq x_1 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil$. When $\lceil \frac{d_1+e_1}{2} \rceil \leq x_1 \leq d_1 - 1$, we have $|P_1| = |P_2| = d_1 - x_1 + 2$, $|P_3| = |P_4| = d_1 - x_1 + 4 \leq d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 4$.

Subcase 3c $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 1 \leq x_1 \leq e_1 - 1$ or $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 1$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$P_1 : x = (x_1, 2) \rightarrow (x_1, 1) \rightarrow (x_1 + \delta(x_1), 1) \rightarrow \cdots \rightarrow (e_1, 1) \rightarrow e$;

$P_2 : x = (x_1, 2) \rightarrow (x_1, 3) \rightarrow (x_1 + \delta(x_1), 3) \rightarrow \cdots \rightarrow (e_1, 3) \rightarrow e$;

$P_3 : x = (x_1, 2) \rightarrow (x_1 + \delta(x_1), 2) \rightarrow \cdots \rightarrow (e_1 + \delta(x_1), 2) \rightarrow (e_1 + \delta(x_1), 1) \rightarrow (e_1 + \delta(x_1), 0) \rightarrow e$;

$P_4 : x = (x_1, 2) \rightarrow (x_1 - \delta(x_1), 2) \rightarrow (x_1 - \delta(x_1), 1) \rightarrow (x_1 - \delta(x_1), 0) \rightarrow (x_1, 0) \rightarrow \cdots \rightarrow (e_1 - \delta(x_1), 0) \rightarrow e$.

We can see $|P_1| = |P_2| = e_1 - x_1 + 2$ and $|P_3| = |P_4| = e_1 - x_1 + 4 = e_1 - d_1 + \lceil \frac{d_1+e_1}{2} \rceil + 3$ if $d_1 - \lceil \frac{d_1+e_1}{2} \rceil + 1 \leq x_1 \leq e_1 - 1$. When $e_1 + 1 \leq x_1 \leq \lceil \frac{d_1+e_1}{2} \rceil - 1$, we have $|P_1| = |P_2| = x_1 - e_1 + 2$, $|P_3| = |P_4| = x_1 - e_1 + 4 \leq \lceil \frac{d_1+e_1}{2} \rceil - e_1 + 3$.

Case 4 $x_2 = 3, x = (x_1, 3)$.

It is similar to Case 2 by vertex-transitivity.

Note that the length of any path of the four in the above cases is at most

$$\max \left\{ e_1 - d_1 + \left\lceil \frac{d_1 + e_1}{2} \right\rceil + 3, \left\lceil \frac{d_1 + e_1}{2} \right\rceil - e_1 + 3, d_1 - \left\lceil \frac{d_1 + e_1}{2} \right\rceil + 4 \right\} \leq d_1 - \left\lceil \frac{d_1 + e_1}{2} \right\rceil + 4.$$

We know S can $(d, 4)$ -dominate any vertex $x \in V(G) - S$. So we have $\gamma_{d,4}(G) \leq 2$.

Thus, we have $\gamma_{d,4}(G) = 2$ if $d_4(G) - (2e_1 - \lfloor \frac{d_1+e_1}{2} \rfloor) \leq d < d_4(G)$. □

Similarly, we can discuss $\gamma_{d,4}(C_{d_1} \times C_{d_2})$ for some special values $d_2 (< d_1)$. Next we consider $\gamma_{d,4}(C_{d_1} \times C_{d_2})$ for $d_1 = d_2 \geq 14$.

Theorem 3.3 *Let $G = C_{d_1} \times C_{d_2}$, $d_1 = d_2 \geq 14$. We have $\gamma_{d,4}(G) = 2$ if $d_4(G) - (e_1 - 2) \leq d < d_4(G)$.*

Proof By Lemma 2.2, we have $d_4(G) - (e_1 - 2) = e_1 + 3$. First assume d_1 and d_2 are odd, and we only need to prove $\gamma_{d,4}(G) \leq 2$ for $d = e_1 + 3$. Let $S = \{o, e\}$, where $o = (0, 0)$, $e = (e_1, e_1)$. For any vertex $x = (x_1, x_2) \in V(G) - S$, where $x_1, x_2 \in \{0, 1, \dots, d_1 - 1\}$, we consider the following cases by vertex-transitivity and symmetry:

Case 1 $0 \leq x_1 \leq e_1$ and $0 \leq x_2 \leq e_1$.

Subcase 1a $x_1 = 1$ or $2, x_2 = 0$.

We construct the four paths P_i ($1 \leq i \leq 4$) as (1), (2), (3), (4). We can see $|P_1| = x_1$, $|P_2| = |P_3| = x_1 + 2$ and $|P_4| = x_1 + 8 \leq 10 \leq e_1 + 3$.

Subcase 1b $3 \leq x_1 \leq e_1 - 3$ and $x_2 = 0$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$P_1 : x = (x_1, 0) \rightarrow (x_1, 1) \rightarrow (x_1 - 1, 1) \rightarrow \cdots \rightarrow (1, 1) \rightarrow (1, 0) \rightarrow o$;

$P_2 : x = (x_1, 0) \rightarrow (x_1 + 1, 0) \rightarrow (x_1 + 1, 1) \rightarrow (x_1 + 1, 2) \rightarrow (x_1, 2) \rightarrow \dots \rightarrow (0, 2) \rightarrow (0, 1) \rightarrow o;$

$P_3 : x = (x_1, 0) \rightarrow (x_1 - 1, 0) \rightarrow (x_1 - 1, -1) \rightarrow (x_1 - 2, -1) \rightarrow \dots \rightarrow (0, -1) \rightarrow o;$

$P_4 : x = (x_1, 0) \rightarrow (x_1, -1) \rightarrow (x_1, -2) \rightarrow (x_1 - 1, -2) \rightarrow \dots \rightarrow (0, -2) \rightarrow (-1, -2) \rightarrow (-1, -1) \rightarrow (-1, 0) \rightarrow o.$

We can see $|P_1| = |P_3| = x_1 + 2, |P_2| = |P_4| = x_1 + 6 \leq e_1 + 3.$

Subcase 1c $e_1 - 2 \leq x_1 \leq e_1 - 1$ and $x_2 = 0.$

We construct the three paths $P_i (1 \leq i \leq 3)$ as (1), (2), (3) and the one other path P_4 as follows:

$P_4 : x = (x_1, 0) \rightarrow (x_1 + 1, 0) \rightarrow \dots \rightarrow (-1, 0) \rightarrow o.$

We can see $|P_1| = x_1, |P_2| = |P_3| = x_1 + 2 \leq e_1 + 1, |P_4| = d_1 - x_1 \leq d_1 - e_1 + 2 = e_1 + 3.$

Subcase 1d $x = (e_1, 0).$

We can easily construct four paths of length at most $e_1 + 2$ in the same way as Subcase 1c of Theorem 3.1.

Subcase 1e $1 \leq x_1 \leq -x_2 + e_1 - 1,$ where $1 \leq x_2 \leq e_1 - 2.$

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

$P_1 : x = (x_1, x_2) \rightarrow (x_1, x_2 - 1) \rightarrow \dots \rightarrow (x_1, 0) \rightarrow (x_1 - 1, 0) \rightarrow \dots \rightarrow o;$ (11)

$P_2 : x = (x_1, x_2) \rightarrow (x_1 - 1, x_2) \rightarrow \dots \rightarrow (0, x_2) \rightarrow (0, x_2 - 1) \rightarrow \dots \rightarrow o;$ (12)

$P_3 : x = (x_1, x_2) \rightarrow (x_1 + 1, x_2) \rightarrow (x_1 + 1, x_2 - 1) \rightarrow \dots \rightarrow (x_1 + 1, 0) \rightarrow (x_1 + 1, -1) \rightarrow (x_1, -1) \rightarrow \dots \rightarrow (0, -1) \rightarrow o;$

$P_4 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow (x_1 - 1, x_2 + 1) \rightarrow \dots \rightarrow (0, x_2 + 1) \rightarrow (-1, x_2 + 1) \rightarrow (-1, x_2) \rightarrow \dots \rightarrow (-1, 0) \rightarrow o.$

We can see $|P_1| = |P_2| = x_1 + x_2, |P_3| = |P_4| = x_1 + x_2 + 4 \leq e_1 + 3.$

Subcase 1f $-x_2 + e_1 - 1 \leq x_1 \leq -x_2 + e_1,$ where $x_1 \geq 1$ and $x_2 \geq 1.$

We construct the two paths P_1, P_2 as (11), (12) and the other two paths P_3, P_4 as follows:

$P_3 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow \dots \rightarrow (x_1, e_1) \rightarrow (x_1 + 1, e_1) \rightarrow \dots \rightarrow e;$

$P_4 : x = (x_1, x_2) \rightarrow (x_1 + 1, x_2) \rightarrow \dots \rightarrow (e_1, x_2) \rightarrow (e_1, x_2 + 1) \rightarrow \dots \rightarrow e.$

We can see $|P_1| = |P_2| = x_1 + x_2 \leq e_1, |P_3| = |P_4| = 2e_1 - x_1 - x_2 \leq e_1 + 1.$

Case 2 $e_1 + 1 \leq x_1 \leq d_1 - 1$ and $0 \leq x_2 \leq e_1.$

Subcase 2a $e_1 + 1 \leq x_1 \leq d_1 - 1$ and $x_2 = 0.$

We can construct four paths between vertices x and o in the same way as Subcase 1a, Subcase 1b, Subcase 1c or Subcase 1d, respectively.

Subcase 2b $x_2 + e_1 + 2 \leq x_1 \leq d_1 - 1,$ where $1 \leq x_2 \leq e_1 - 2.$

We construct the four paths $P_i (1 \leq i \leq 4)$ as follows:

$P_1 : x = (x_1, x_2) \rightarrow (x_1 + 1, x_2) \rightarrow \dots \rightarrow (d_1 - 1, x_2) \rightarrow (0, x_2) \rightarrow (0, x_2 - 1) \rightarrow \dots \rightarrow (0, 1) \rightarrow o;$ (13)

$P_2 : x = (x_1, x_2) \rightarrow (x_1, x_2 - 1) \rightarrow \dots \rightarrow (x_1, 0) \rightarrow (x_1 + 1, 0) \rightarrow \dots \rightarrow (-1, 0) \rightarrow o;$ (14)

$P_3 : x = (x_1, x_2) \rightarrow (x_1 - 1, x_2) \rightarrow (x_1 - 1, x_2 - 1) \rightarrow \dots \rightarrow (x_1 - 1, 0) \rightarrow (x_1 - 1, d_1 - 1) \rightarrow (x_1, d_1 - 1) \rightarrow \dots \rightarrow (d_1 - 1, d_1 - 1) \rightarrow (0, d_1 - 1) \rightarrow o;$

$P_4 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow (x_1 + 1, x_2 + 1) \rightarrow \dots \rightarrow (d_1 - 1, x_2 + 1) \rightarrow (0, x_2 + 1) \rightarrow (1, x_2 + 1) \rightarrow (1, x_2) \rightarrow \dots \rightarrow (1, 0) \rightarrow o.$

We can see $|P_1| = |P_2| = d_1 - x_1 + x_2$, $|P_3| = |P_4| = d_1 - x_1 + x_2 + 4 \leq d_1 - e_1 + 2 = e_1 + 3$.

Subcase 2c $x_2 + e_1 \leq x_1 \leq x_2 + e_1 + 1$, where $x_1 \leq d_1 - 1$ and $x_2 \geq 1$.

We construct the two paths P_i ($1 \leq i \leq 2$) as (13), (14) and the other two paths P_i ($3 \leq i \leq 4$) as follows:

$$P_1 : x = (x_1, x_2) \rightarrow (x_1 - 1, x_2) \rightarrow \cdots \rightarrow (e_1, x_2) \rightarrow (e_1, x_2 + 1) \rightarrow \cdots \rightarrow e;$$

$$P_2 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow \cdots \rightarrow (x_1, e_1) \rightarrow (x_1 - 1, e_1) \rightarrow \cdots \rightarrow e;$$

We can see $|P_1| = |P_2| = d_1 - x_1 + x_2 \leq d_1 - e_1 = e_1 + 1$, $|P_3| = |P_4| = x_1 - x_2 \leq e_1 + 1$.

Case 3 $0 \leq x_1 \leq e_1$ and $e_1 + 1 \leq x_2 \leq d_1 - 1$.

It is similar to Case 2 by vertex-transitivity and symmetry.

Case 4 $e_1 + 1 \leq x_1 \leq d_1 - 1$ and $e_1 + 1 \leq x_2 \leq d_1 - 1$.

Subcase 4a $e_1 + 1 \leq x_1 \leq -x_2 + d_1 + e_1 - 2$, where $e_1 + 1 \leq x_2 \leq d_1 - 3$.

We construct the four paths P_i ($1 \leq i \leq 4$) as follows:

$$P_1 : x = (x_1, x_2) \rightarrow (x_1 - 1, x_2) \rightarrow \cdots \rightarrow (e_1, x_2) \rightarrow (e_1, x_2 - 1) \rightarrow \cdots \rightarrow (e_1, e_1 + 1) \rightarrow e; \tag{15}$$

$$P_2 : x = (x_1, x_2) \rightarrow (x_1, x_2 - 1) \rightarrow \cdots \rightarrow (x_1, e_1) \rightarrow (x_1 - 1, e_1) \rightarrow \cdots \rightarrow (e_1 + 1, e_1) \rightarrow e; \tag{16}$$

$$P_3 : x = (x_1, x_2) \rightarrow (x_1 + 1, x_2) \rightarrow (x_1 + 1, x_2 - 1) \rightarrow \cdots \rightarrow (x_1 + 1, e_1 - 1) \rightarrow (x_1, e_1 - 1) \rightarrow (e_1, e_1 - 1) \rightarrow e;$$

$$P_4 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow (x_1 - 1, x_2 + 1) \rightarrow \cdots \rightarrow (e_1 - 1, x_2 + 1) \rightarrow (e_1 - 1, x_2) \rightarrow \cdots \rightarrow (e_1 - 1, e_1) \rightarrow e.$$

We can see $|P_1| = |P_2| = x_1 + x_2 - 2e_1$, $|P_3| = |P_4| = x_1 + x_2 - 2e_1 + 4 \leq e_1 + 3$.

Subcase 4b $-x_2 + d_1 + e_1 - 1 \leq x_1 \leq -x_2 + d_1 + e_1$, where $x_1 \geq e_1 + 1$ and $x_2 \geq e_1 + 1$.

We construct the two paths P_i ($1 \leq i \leq 2$) as (15), (16) and the other two paths P_i ($3 \leq i \leq 4$) as follows:

$$P_3 : x = (x_1, x_2) \rightarrow (x_1 + 1, x_2) \rightarrow \cdots \rightarrow (d_1 - 1, x_2) \rightarrow (0, x_2) \rightarrow (0, x_2 + 1) \rightarrow \cdots \rightarrow (0, d_1 - 1) \rightarrow o;$$

$$P_4 : x = (x_1, x_2) \rightarrow (x_1, x_2 + 1) \rightarrow \cdots \rightarrow (x_1, d_1 - 1) \rightarrow (x_1, 0) \rightarrow (x_1 + 1, 0) \rightarrow \cdots \rightarrow (d_1 - 1, 0) \rightarrow o.$$

We can see $|P_1| = |P_2| = x_1 + x_2 - 2e_1 \leq d_1 - e_1 = e_1 + 1$, $|P_3| = |P_4| = 2d_1 - x_1 - x_2 \leq d_1 - e_1 + 1 = e_1 + 2$.

From the above cases, we know there exist four internally disjoint (x, S) -paths of length at most $e_1 + 3$ in G . So S can $(d, 4)$ -dominate any vertex x in $G - S$, S is a $(d, 4)$ -dominating set of G .

Next we assume d_1 and d_2 are even, and so we only need to consider the case of $0 \leq x_1 \leq e_1$ and $0 \leq x_2 \leq e_1$ in the same way as Case 1 by vertex-transitivity and symmetry. We easily see S can also $(d, 4)$ -dominate any vertex in $V(G) - S$.

So we have $\gamma_{d,4}(G) \leq 2$. □

4 Conclusion and Problems

For the undirected toroidal mesh $C(d_1, d_2, \dots, d_n)$, we prove that $\gamma_{d,4}(C(d_1, 3)) = 2$ if and only if $d_4(G) - e_1 \leq d < d_4(G)$ for $d_1 \geq 5$, $\gamma_{d,4}(C(d_1, 4)) = 2$ if $d_4(G) - (2e_1 - \lfloor \frac{d_1 + e_1}{2} \rfloor) \leq d < d_4(G)$

for $d_1 \geq 24$, and $\gamma_{d,4}(C(d_1, d_2)) = 2$ if $d_4(G) - (e_1 - 2) \leq d < d_4(G)$ for $d_1 = d_2 \geq 14$. To determine the values of $\gamma_{d,4}(C(d_1, 4))$ for $d < d_4(G) - (2e_1 - \lfloor \frac{d_1 + e_1}{2} \rfloor)$, $\gamma_{d,4}(C(d_1, d_2))$ for $d < d_4(G) - (e_1 - 2)$ and $\gamma_{d,2n}(C(d_1, d_2, \dots, d_n))$ for $d < \text{diam}(C(d_1, d_2, \dots, d_n))$ are worth studying further.

Acknowledgements The authors would like to thank the referees for their careful reading and helpful comments.

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