

# Generalized Measures of Edge Fault Tolerance in $(n, k)$ -star Graphs

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**Abstract:** This paper considers a kind of generalized measure  $\lambda_s^{(h)}$  of fault tolerance in the  $(n, k)$ -star graph  $S_{n,k}$  for  $2 \leq k \leq n-1$  and  $0 \leq h \leq n-k$ , and determines  $\lambda_s^{(h)}(S_{n,k}) = \min\{(n-h-1)(h+1), (n-k+1)(k-1)\}$ , which implies that at least  $\min\{(n-k+1)(k-1), (n-h-1)(h+1)\}$  edges of  $S_{n,k}$  have to remove to get a disconnected graph that contains no vertices of degree less than  $h$ . This result shows that the  $(n, k)$ -star graph is robust when it is used to model the topological structure of a large-scale parallel processing system.

**Keywords:** Combinatorics, fault-tolerant analysis,  $(n, k)$ -star graphs, edge-connectivity,  $h$ -super edge-connectivity

## I. Introduction

It is well known that interconnection networks play an important role in parallel computing/communication systems. An interconnection network can be modeled by a graph  $G = (V, E)$ , where  $V$  is the set of processors and  $E$  is the set of communication links in the network. The connectivity  $\lambda(G)$  of a graph  $G$  is an important measurement for fault-tolerance of the network, and the larger  $\lambda(G)$  is, the more reliable the network is.

A subset of vertices  $B$  of a connected graph  $G$  is called an *edge-cut* if  $G - B$  is disconnected. The *edge connectivity*  $\lambda(G)$  of  $G$  is defined as the minimum cardinality over all edge-cuts of  $G$ . Because  $\lambda$  has many shortcomings, one proposes the concept of the  $h$ -super edge connectivity of  $G$ , which can measure fault tolerance of an interconnection network more accurately than the classical connectivity  $\lambda$ .

A subset of vertices  $B$  of a connected graph  $G$  is called an  *$h$ -super edge-cut*, or  *$h$ -edge-cut* for short, if  $G - B$  is disconnected and has the minimum degree at least  $h$ . The  *$h$ -super edge-connectivity* of  $G$ , denoted by  $\lambda_s^{(h)}(G)$ , is defined as the minimum cardinality over all  $h$ -edge-cuts of  $G$ . It is clear that, if  $\lambda_s^{(h)}(G)$  exists, then

$$\lambda(G) = \lambda_s^{(0)}(G) \leq \lambda_s^{(1)}(G) \leq \lambda_s^{(2)}(G) \leq \dots \leq \lambda_s^{(h-1)}(G) \leq \lambda_s^{(h)}(G).$$

For any graph  $G$  and integer  $h$ , determining  $\lambda_s^{(h)}(G)$  is quite difficult. In fact, the existence of  $\lambda_s^{(h)}(G)$  is an open problem so far when  $h \geq 1$ . Some results have been obtained on  $\lambda_s^{(h)}(G)$  for particular classes of graphs and small  $h$ 's (see Section 16.7 in [5]).

This paper is concerned about  $\lambda_s^{(h)}$  for the  $(n, k)$ -star graph  $S_{n,k}$ . For the  $h$ -super connectivity, several authors have done some work. For  $k = n-1$ ,  $S_{n,n-1}$  is isomorphic to a star graph  $S_n$ . Akers and Krishnamurthy [1] determined  $\lambda(S_n) = n-1$  for  $n \geq 2$  and  $\lambda_s^{(1)}(S_n) = 2n-4$  for  $n \geq 3$ . In this paper, we show the following result.

**Theorem:** If  $2 \leq k \leq n-1$  and  $0 \leq h \leq n-k$ , then



$$\lambda_s^{(h)}(S_{n,k}) = \begin{cases} (n-h-1)(h+1) & \text{for } h \leq k-2 \text{ and } h \leq \frac{n}{2}-1, \\ (n-k+1)(k-1) & \text{otherwise.} \end{cases}$$

This result implies that at least  $\min\{(n-k+1)(k-1), (n-h-1)(h+1)\}$  edges of  $S_{n,k}$  have to remove to get a disconnected graph that contains no vertices of degree less than  $h$ . The proof of this result is in Section 3. In Section 2, we recall the structure of  $S_{n,k}$  and some lemmas used in our proofs.

## II. Definitions and Lemmas

For integers  $n$  and  $k$  with  $1 \leq k \leq n-1$ , let  $I_n = \{1, 2, \dots, n\}$  and  $P(n, k) = \{p_1 p_2 \dots p_k : p_i \in I_n, p_i \neq p_j, 1 \leq i \neq j \leq k\}$ , the set of  $k$ -permutations on  $I_n$ . Clearly,  $|P(n, k)| = n!/(n-k)!$ .

**Definition 2.1:** The  $(n, k)$ -star graph  $S_{n,k}$  is a graph with vertex-set  $P(n, k)$ . The adjacency is defined as follows: a vertex  $p = p_1 p_2 \dots p_i \dots p_k$  is adjacent to a vertex

- (a)  $p_i p_2 \dots p_{i-1} p_1 p_{i+1} \dots p_k$ , where  $2 \leq i \leq k$  (swap  $p_1$  with  $p_i$ ).
- (b)  $\alpha p_2 p_3 \dots p_k$ , where  $\alpha \in I_n \setminus \{p_i : 1 \leq i \leq k\}$  (replace  $p_1$  by  $\alpha$ ).

The vertices of type (a) are referred to as *swap-neighbors* of  $p$  and the edges between them are referred to as *swap-edge* or  *$i$ -edges*. The vertices of type (b) are referred to as *unswap-neighbors* of  $p$  and the edges between them are referred to as *unswap-edges*. Clearly, every vertex in  $S_{n,k}$  has  $k-1$  swap-neighbors and  $n-k$  unswap-neighbors. Usually, if  $x = p_1 p_2 \dots p_k$  is a vertex in  $S_{n,k}$ , we call  $p_i$  the  $i$ -th bit for each  $i \in I_k$ .

The  $(n, k)$ -star graph  $S_{n,k}$  is proposed by Chiang and Chen [4]. Some nice properties of  $S_{n,k}$  are compiled by Cheng and Lipman (see Theorem 1 in [2]).

**Lemma 2.2:**  $S_{n,k}$  is  $(n-1)$ -regular  $(n-1)$ -connected.

**Lemma 2.3:** For any  $\alpha = p_1 p_2 \dots p_{k-1} \in P(n, k-1)$  ( $k \geq 2$ ), let  $V_\alpha = \{p\alpha : p \in I_n \setminus \{p_i : i \in I_{k-1}\}\}$ . Then the subgraph of  $S_{n,k}$  induced by  $V_\alpha$  is a complete graph of order  $n-k+1$ , denoted by  $K_{n-k+1}^\alpha$ .

Let  $S_{n-1, k-1}^{t:i}$  denote the subgraph of  $S_{n,k}$  induced by vertices with the  $t$ -th bit  $i$  for  $2 \leq t \leq k$ . The following lemma is a slight modification of the result of Chiang and Chen [4].

**Lemma 2.4:** For a fixed integer  $t$  with  $2 \leq t \leq k$ ,  $S_{n,k}$  can be decomposed into  $n$  subgraphs  $S_{n-1, k-1}^{t:i}$ , which is isomorphic to  $S_{n-1, k-1}$ , for each  $i \in I_n$ . Moreover, there are  $\frac{(n-2)!}{(n-k)!}$  independent swap-edges between  $S_{n-1, k-1}^{t:i}$  and  $S_{n-1, k-1}^{t:j}$  for any  $i, j \in I_n$  with  $i \neq j$ .

Since  $S_{n,1} \cong K_n$ , we only consider the case of  $k \geq 2$  in the following discussion.

**Lemma 2.5:** If  $2 \leq k \leq n-1$  and  $0 \leq h \leq n-k$ , then

$$\lambda_s^{(h)}(S_{n,k}) \leq \begin{cases} (n-h-1)(h+1) & \text{for } h \leq \frac{n}{2}-1, \\ (n-k+1)(k-1) & \text{otherwise.} \end{cases}$$

**Proof:** By our hypothesis of  $2 \leq k \leq n-1$ , for any  $\alpha \in P(n, k-1)$ , we can choose a subset  $X \subseteq V(K_{n-k+1}^\alpha)$  such that  $|X| = h+1$ . Then the subgraph of  $K_{n-k+1}^\alpha$  induced by  $X$  is a complete graph  $K_{h+1}$ . Let  $B$  be the set of incident edges with and not within  $X$ . Since  $S_{n,k}$  is  $(n-1)$ -regular and  $K_{h+1}$  is  $h$ -regular, we have that

$$|B| = (n-h-1)(h+1).$$

Clearly,  $B$  is an edge-cut of  $S_{n,k}$ . Let  $x$  be any vertex in  $S_{n,k} - X$ , and  $d(x)$  denote the number of edges incident with  $x$  in  $S_{n,k} - X$ . In order to prove that  $B$  is an  $h$ -edge-cut, we only need to show  $d(x) \geq h$ . Note that  $X$  is contained in  $S_{n-1, k-1}^i$  and edges between  $S_{n-1, k-1}^i$  and  $S_{n-1, k-1}^j$  are independent for any  $i, j \in I_n$  with  $i \neq j$  by

**Lemma 2.4.** If  $x$  is in  $S_{n-1,k-1}^i - K_{n-k+1}^\alpha$  or is in  $S_{n-1,k-1}^j$  with  $i \neq j$ , then  $d(x) \geq n - 2 \geq n - k \geq h$ . For  $x \in V(K_{n-k+1}^\alpha - X)$ , if exists, then  $d(x) = n - 1 - |X| = n - h - 2 \geq h$  for  $h \leq \frac{n}{2} - 1$ . Therefore,  $B$  is an  $h$ -edge-cut of  $S_{n,k}$ , and so

$$\lambda_s^{(h)}(S_{n,k}) \leq |B| = (n - h - 1)(h + 1) \text{ for } h \leq \frac{n}{2} - 1.$$

For  $h \geq \frac{n}{2}$ , we choose  $X = V(K_{n-k+1}^\alpha)$ . Then  $|B| = (n - k + 1)(k - 1)$ . For any  $x$  in  $S_{n-1,k-1}^i - X$  or  $S_{n-1,k-1}^j$  with  $i \neq j$ , we have  $d(x) \geq n - 2 \geq n - k \geq h$ . Thus,  $B$  is an  $h$ -edge-cut of  $S_{n,k}$ , and so

$$\lambda_s^{(h)}(S_{n,k}) \leq |B| = (n - k + 1)(k - 1) \text{ for } h \geq \frac{n}{2}.$$

The lemma follows.

**Corollary 2.6:**  $\lambda_s^{(h)}(S_{n,2}) = n - 1$  for  $0 \leq h \leq n - 2$ .

**Proof:** On the one hand,  $\lambda_s^{(h)}(S_{n,2}) \leq n - 1$  by Lemma 2.5 when  $k = 2$ . On the other hand,  $\lambda_s^{(h)}(S_{n,2}) \geq \lambda(S_{n,2}) = n - 1$  by Lemma 2.2.

The following lemma shows the relations between  $(n - h - 1)(h + 1)$  and  $(n - k + 1)(k - 1)$ .

**Lemma 2.7:** For  $2 \leq k \leq n - 1, 0 \leq h \leq n - k$ , let

$$\psi(h, k) = \min\{(n - h - 1)(h + 1), (n - k + 1)(k - 1)\}.$$

If  $h \leq \frac{n}{2} - 1$ , then

$$\psi(h, k) = \begin{cases} (n - h - 1)(h + 1) & \text{if } 0 \leq h \leq k - 2; \\ (n - k + 1)(k - 1) & \text{if } h \geq k - 1. \end{cases}$$

**Proof:** Let  $f(x) = (n - x)x$ , then  $\psi(h, k) = \min\{f(h + 1), f(k - 1)\}$ . It can be easily checked that  $f(x)$  is a convex function on the interval  $[0, n]$ , the maximum value is reached at  $x = \frac{n}{2}$ . Thus,  $f(x)$  is an increasing function on the interval  $[0, \frac{n}{2}]$ .

If  $0 \leq h \leq k - 2$ , then  $h + 1 \leq k - 1$ . Since  $h \leq n - k, h + 1 \leq n - k + 1$  and  $\min\{k - 1, n - k + 1\} \leq \frac{n}{2}$ . Thus, when  $h \leq \frac{n}{2} - 1, f(h + 1) \leq f(k - 1) = f(n - k + 1)$ , and so  $\psi(h, k) = f(h + 1) = (n - h - 1)(h + 1)$ .

If  $h \geq k - 1$ , then  $k - 1 < h + 1 \leq \frac{n}{2}, f(k - 1) < f(h + 1)$ , so  $\psi(n, k) = f(k - 1) = (n - k + 1)(k - 1)$ .

The lemma follows.

To state and prove our main results, we need some notations. Let  $B$  be a minimum  $h$ -edge-cut of  $S_{n,k}$ . Clearly,  $S_{n,k} - B$  has exactly two connected components. Let  $X$  and  $Y$  be two vertex-set of two connected components of  $S_{n,k} - B$ . For a fixed  $t \in I_k \setminus \{1\}$  and any  $i \in I_n$ , let

$$\begin{aligned} X_i &= X \cap V(S_{n-1,k-1}^{t:i}), \\ Y_i &= Y \cap V(S_{n-1,k-1}^{t:i}), \\ B_i &= B \cap E(S_{n-1,k-1}^{t:i}) \text{ and} \\ B_{ij} &= B \cap E(S_{n-1,k-1}^{t:i}, S_{n-1,k-1}^{t:j}), \end{aligned}$$

and let

$$\begin{aligned} J &= \{i \in I_n : X_i \neq \emptyset\}, \\ J' &= \{i \in J : Y_i \neq \emptyset\} \text{ and} \\ T &= \{i \in I_n : Y_i \neq \emptyset\}. \end{aligned}$$

**Lemma 2.8:** Let  $B$  be a minimum  $h$ -edge-cut of  $S_{n,k}$  and  $X$  be the vertex-set of a connected component of  $S_{n,k} - B$ . If  $3 \leq k \leq n - 1$  and  $1 \leq h \leq n - k$  then, for any  $t \in I_k \setminus \{1\}$ ,

(a)  $B_i$  is an  $(h - 1)$ -edge-cut of  $S_{n-1,k-1}^{t:i}$  for any  $i \in J'$ ,



$$(b) \lambda_s^{(h)}(S_{n,k}) \geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1}).$$

**Proof.** (a) By the definition of  $J'$ ,  $B_i$  is an edge-cut of  $S_{n-1,k-1}^{t:i}$  for any  $i \in J'$ . For any vertex  $x$  in  $S_{n-1,k-1}^{t:i} - B_i$ , since  $x$  has degree at least  $h$  in  $S_{n,k} - S$  and has exactly one neighbor outsider  $S_{n-1,k-1}^{t:i}$ ,  $x$  has degree at least  $h-1$  in  $S_{n,k}^{t:i} - B_i$ . This fact shows that  $B_i$  is an  $(h-1)$ -edge-cut of  $S_{n-1,k-1}^{t:i}$  for any  $i \in J'$ .

(b) By the assertion (a), we have  $|B_i| \geq \lambda_s^{(h-1)}(S_{n-1,k-1})$ , and so

$$\lambda_s^{(h)}(S_{n,k}) = |B| \geq \sum_{i \in J'} |B_i| \geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1}).$$

The lemma follows.

### III. Proof of Theorem

By Lemma 2.5 and Lemma 2.7, we only need to prove that, for  $2 \leq k \leq n-1$  and  $0 \leq h \leq n-k$ ,

$$\lambda_s^{(h)}(S_{n,k}) \geq \begin{cases} (n-h-1)(h+1) & \text{for } h \leq k-2 \text{ and } h \leq \frac{n}{2}-1, \\ (n-k+1)(k-1) & \text{otherwise.} \end{cases}$$

Let  $\omega(h,k) = \max\{(n-h-1)(h+1), (n-k+1)(k-1)\}$ .

We proceed by induction on  $k \geq 2$  and  $h \geq 0$ . The inequality is true for  $k=2$  and any  $h$  with  $0 \leq h \leq n-2$  by Corollary 2.6. The inequality is also true for  $h=0$  and any  $k$  with  $2 \leq k \leq n-1$  since  $\lambda_s^{(0)}(S_{n,k}) = \lambda(S_{n,k}) = n-1$ . Assume the induction hypothesis for  $k-1$  with  $k \geq 3$  and for  $h-1$  with  $h \geq 1$ , that is,

$$\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq \begin{cases} (n-h)h & \text{for } h \leq k-3 \text{ and } h \leq \frac{n-1}{2}, \\ (n-k+2)(k-2) & \text{otherwise.} \end{cases}$$

Let  $B$  be a minimum  $h$ -edge-cut of  $S_{n,k}$  and  $X$  be the vertex-set of a minimum connected component of  $S_{n,k} - B$ . By Lemma 2.5, we have

$$|B| \leq \omega(h,k). \quad (1)$$

Use notations defined in Section II. Choose  $t \in I_k \setminus \{1\}$  such that  $|J|$  is as large as possible. For each  $i \in I_n$ , we write  $S_{n-1,k-1}^i$  for  $S_{n-1,k-1}^{t:i}$  for short.

We first show  $|J|=1$ . Suppose to the contrary  $|J| \geq 2$ . We will deduce contradictions by considering three cases depending on  $|J'|=0$ ,  $|J'|=1$  or  $|J'| \geq 2$ .

*Case 1.*  $|J'|=0$ .

In this case,  $X_i \neq \emptyset$  and  $Y_i = \emptyset$  for each  $i \in J$ , that is,  $J \cap T = \emptyset$ . By  $|J| \geq 2$  and the minimality of  $X$ ,  $|T| \geq 2$ . Assume  $\{i_1, i_2\} \subseteq J$  and  $\{i_3, i_4\} \in T$ . By Lemma 2.4, there are  $\frac{(n-2)!}{(n-k)!}$  independent swap-edges between  $S_{n-1,k-1}^{i_1}$  (resp.  $S_{n-1,k-1}^{i_2}$ ) and  $S_{n-1,k-1}^{i_3}$  (resp.  $S_{n-1,k-1}^{i_4}$ ), all of which are contained in  $B$ . Since  $J \cap T = \emptyset$ , we have that

$$|B| \geq 4 \frac{(n-2)!}{(n-k)!}.$$

For  $k=3$ ,

$$|B| \geq 4 \frac{(n-2)!}{(n-k)!} \geq 4(n-2) > 2(n-2)$$

Combining Lemma 2.5 with Lemma 2.7 yields  $|B| \leq \lambda_s^{(h)}(S_{n,3}) \leq 2(n-2)$ , a contradiction. For  $k \geq 4$ , it is easy to check that

$$\begin{aligned}
 |B| &\geq 4 \frac{(n-2)!}{(n-k)!} \geq 4(n-2)(n-3) = (2n-4)(2n-6) \\
 &> \max\{(n-h-1)(h+1), (n-k+1)(k-1)\} \\
 &= \omega(h, k),
 \end{aligned}$$

which contradicts the inequality (1).

*Case 2.*  $|J'| = 1$ .

Without loss of generality, assume  $J' = \{1\}$ . By Lemma 2.8 (a),  $B_1$  is an  $(h-1)$ -edge-cut of  $S_{n-1, k-1}^1$ . By  $|J| \geq 2$ , there exists an  $i \in J - J'$  such that  $X_i = V(S_{n-1, k-1}^i)$ . By the minimality of  $X$ , there exists some  $j \in T - J'$  such that  $Y_j = V(S_{n-1, k-1}^j)$ . By Lemma 2.4, there are  $\frac{(n-2)!}{(n-k)!}$  independent swap-edges between  $S_{n-1, k-1}^i$  and  $S_{n-1, k-1}^j$ , thus  $|B_{ij}| = \frac{(n-2)!}{(n-k)!} \geq n-2$ . We consider the following two cases.

If  $\lambda_s^{(h-1)}(S_{n-1, k-1}) \geq (n-h)h$ , then

$$\begin{aligned}
 |B| &\geq |B_1| + |B_{ij}| \\
 &\geq (n-h)h + (n-2) \\
 &> (n-h-1)h + (n-h-1) \\
 &= (n-h-1)(h+1),
 \end{aligned}$$

If  $\lambda_s^{(h-1)}(S_{n-1, k-1}) \geq (n-k+2)(k-2)$ , then

$$\begin{aligned}
 |B| &\geq |B_1| + |B_{ij}| \\
 &\geq (n-k+2)(k-2) + (n-2) \\
 &> (n-k+1)(k-2) + (n-k+1) \\
 &= (n-k+1)(k-1).
 \end{aligned}$$

Therefore, we have  $|B| > \omega(h, k)$ , which contradicts the inequality (1).

*Case 3.*  $|J'| \geq 2$ .

By Lemma 2.8 (b), if  $\lambda_s^{(h-1)}(S_{n-1, k-1}) \geq (n-h)h$  then

$$\begin{aligned}
 |B| &\geq |J'| \lambda_s^{(h-1)}(S_{n-1, k-1}) \\
 &\geq 2(n-h)h \geq (n-h)h + (n-h) \\
 &> (n-h-1)(h+1),
 \end{aligned}$$

if  $\lambda_s^{(h-1)}(S_{n-1, k-1}) \geq (n-k+2)(k-2)$  then

$$\begin{aligned}
 |B| &\geq |J'| \lambda_s^{(h-1)}(S_{n-1, k-1}) \\
 &\geq 2(n-k+2)(k-2) \\
 &\geq (n-k+2)(k-2) + (n-k+2) \\
 &> (n-k+1)(k-1).
 \end{aligned}$$

Therefore, we have  $|B| > \omega(h, k)$ , which contradicts the inequality (1).

Thus, we have  $|J| = 1$ . By the choice of  $t$ , the  $i$ -th bits of all vertices in  $X$  are same for each  $i = 2, 3, \dots, k$ , and so  $X$  is a complete graph. Thus, we have that

$$\lambda_s^{(h)}(S_{n, k}) = |B| = (n - |X|)|X|$$

Since  $h+1 \leq |X| \leq n-k+1$  and  $f(x) = (n-x)x$  is a convex function on the interval  $[0, n]$ , we have that

$$\lambda_s^{(h)}(S_{n, k}) = |B| = (n - |X|)|X| \geq \psi(h, k),$$

where  $\psi(h, k)$  is defined in Lemma 2.7.

If  $h \leq \frac{n}{2} - 1$ , using Lemma 2.7, we have



$$\lambda_s^{(h)}(S_{n,k}) \geq \psi(h, k) = \begin{cases} (n-h-1)(h+1) & \text{if } 0 \leq h \leq k-2; \\ (n-k+1)(k-1) & \text{if } h \geq k-1. \end{cases}$$

If  $h \geq \frac{n}{2}$ , then  $X = V(K_{n-k+1})$ . Otherwise, there exists some  $x \in V(K_{n-k+1} - X)$  such that

$$h \leq d(x) = n-1 - |X| \leq n-h-2,$$

which implies  $h \leq \frac{n}{2} - 1$ , a contradiction. Therefore, we have  $|X| = n-k+1$ , and

$$\lambda_s^{(h)}(S_{n,k}) = |B| = (n-|X|)|X| = (n-k+1)(k-1) \text{ for } h \geq \frac{n}{2}.$$

By the induction principle, the theorem follows.

As we have known, when  $k = n-1$ ,  $S_{n,n-1}$  is isomorphic to the star graph  $S_n$ . Akers and Krishnamurthy [1] determined  $\lambda(S_n)$  and  $\lambda_s^{(1)}(S_n)$ , which can be obtained from our result by setting  $k = n-1$  and  $h = 0, 1$ , respectively.

**Corollary:**  $\lambda(S_n) = n-1$  for  $n \geq 2$  and  $\lambda_s^{(1)}(S_n) = 2n-4$  for  $n \geq 3$ .

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