

~~About Symplectic Group~~ Exlecture Note

1. definition: ~~标准正交群 $SO(2n)$ $J = \begin{pmatrix} 0 & I \\ -I & 0 \end{pmatrix}$ $\det(J) = 1$ $J^2 = -I$ $J^{-1} = J^T = -J$~~

~~$Sp(2n, \mathbb{R}) = \{ S \mid S^T J S = S J S^T = J \}$~~

~~S symplectic $\Rightarrow (S^{-1})^T J S^{-1} = -S^{-1} J^T (S^{-1})^{-1} = -(S J S^T)^{-1} = J$~~

~~$S^{-1} J S^{-T} = -\frac{S^{-1} J S^{-1}}{(S^T J^T S)^{-1}} = -\frac{S^{-1} J S^{-1}}{(S^T J S)^{-1}} = -\frac{S^{-1} J S^{-1}}{(S^T J S)^{-1}} = J$~~

~~Proof:~~

~~$S \in Sp(2n, \mathbb{R}) \Leftrightarrow S^T J S = J \Leftrightarrow S J S^T = J$~~

~~pf: (\Leftarrow) : $S^T J S = J \Rightarrow S^{-1} J S^{-T} = J \Rightarrow J S^{-T} = S J$
 $\Rightarrow J = S J S^T$~~

~~$S = \begin{bmatrix} 0 & I \\ -I & 0 \end{bmatrix}$ $Sp_{2n} = \{ P \in GL_{2n}(\mathbb{R}) \mid P^T S P = S \}$~~

① $Sp_{2n}(\mathbb{R})$ is a group: $\{$ Linear Groups (Some Calculation)

~~$P^T S P = S$ (i)~~

~~$\Rightarrow S = P^{-T} S P^{-1}$ $P^{-1} = P^{-T} \Rightarrow P^{-1} \in Sp_{2n}$~~

Multiplication-closed is trivial.

② $P \in Sp_{2n}(\mathbb{R}) \Leftrightarrow P^T \in Sp_{2n}(\mathbb{R})$

Take inverse in (i), we have $P^{-1} S^{-1} P^{-T} = S^{-1}$

But $S^{-1} = S^T = -S$, as you notice.

Thus $S = P S P^T$ Hence $P \in Sp_{2n}$.

③ $P \in Sp_{2n}(\mathbb{R}) \Rightarrow \det P = 1$.

It's not hard to prove in the case when $n=1$.

For general n , we refer to [线性代数讲义](#)

④ $Sp_2 = SL_2$ $Sp_4 \neq SL_4$

Notice that SL_2 is generated by $\left\{ \begin{pmatrix} 1 & a \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \begin{pmatrix} a & 0 \\ 0 & a^{-1} \end{pmatrix} \right\}$ $a \neq 0$ here.

We need to check these 3 class of matrices are in Sp_2 .



$$\begin{pmatrix} 1 & a \\ & 1 \end{pmatrix} \begin{pmatrix} 1 & -1 \\ & 1 \end{pmatrix} \begin{pmatrix} 1 & \\ a & 1 \end{pmatrix} = \begin{pmatrix} a & -1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & \\ a & 1 \end{pmatrix} = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 1 & -1 \\ & 1 \end{pmatrix} \begin{pmatrix} 1 & \\ a & 1 \end{pmatrix} \begin{pmatrix} 1 & a \\ & 1 \end{pmatrix} = \begin{pmatrix} 1 & -1 \\ & 1 \end{pmatrix}$$

$$\begin{pmatrix} a & -1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & -1 \\ & 1 \end{pmatrix} \begin{pmatrix} 1 & a \\ a & 1 \end{pmatrix} = \begin{pmatrix} a^{-1} & -a \\ & 1 \end{pmatrix} \begin{pmatrix} a & -1 \\ a & 1 \end{pmatrix} = \begin{pmatrix} 1 & -1 \\ & 1 \end{pmatrix}$$

Remark: $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \xrightarrow{a \neq 0} \begin{pmatrix} a & b \\ c & d - \frac{bc}{a} \end{pmatrix} \xrightarrow{\text{multiply } \begin{pmatrix} 1 & \\ & a^{-1} \end{pmatrix} \text{ if necessary}} \begin{pmatrix} a & b \\ & a^{-1} \end{pmatrix}$

By the same tone we check that $SL_2(\mathbb{Z}) = \langle \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & \\ 0 & -1 \end{pmatrix} \rangle$ In fact: notice that $ab-cd=1 \Rightarrow$ coprime \Rightarrow easy to do with $n=2$

since $(a, a^{-1}) = \mathbb{Z} \Rightarrow a = \pm 1$. (P32/1.3.24.)

The reason why we want to ~~find~~ find generators as small as possible is presented above.

• $SP_4(\mathbb{R}) \neq SL_4(\mathbb{R})$

~~It's~~ It's not such hard as it seems to find an example, right! Just try $\begin{pmatrix} 1 & 1 \\ & 1 \end{pmatrix}$. The reason I put an -1 here is

to make the det = 1

Having got some taste of calculation in matrix groups, we begin our discussion of a nice theorem.

Theorem: $F : \text{or } |F| \geq 4$. (& finite order: q)

a). The only proper normal subgroup of $SL_2(F)$ is its center $Z = \{\pm I\}$

b). $PSL_2(F)$ is a simple group

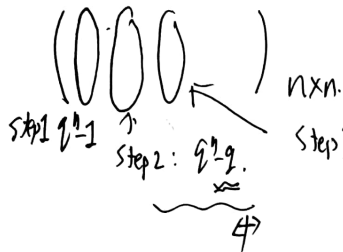
Lemma. $q = |F|$; Then $SL_2(\mathbb{F}_q) = q^3 - q$. (in fact here we can calculate SL_n .
if q is not power of 2. $PSL_2(\mathbb{F}_q) = \frac{1}{2}q^3 - q$.
if q is a ---, --- $PSL_2(\mathbb{F}_q) = SL_2(\mathbb{F}_q)$)

Pf: Recall a famous group homomorphism: $GL_n \rightarrow F^*$ with kernel

SL_n . $|F^*| = q-1$ thus we are reduced to calculate $\#(GL_n)$



Here is a trick. "GL" means the column vectors of matrix A should be linear independent with each other. Thus:



\leftrightarrow : why is " q " here? Note that $\#\{\lambda v\} = q$, where $\lambda \in \mathbb{F}_q$.

$$\text{Thus we have done: } (q^{n-1}) \cdots (q^n - q^{n-1}) = q^{\frac{n(n-1)}{2}} (q^{n-1}) \cdots (q-1)$$

$$\text{Hence } \#SL_2(\mathbb{F}_q) = \frac{(q^2-1)(q^2-q)}{q-1} = q(q^2-1)$$

$PSL_2(\mathbb{F}_q)$ doesn't lose half of its weight since $1 = -1$ in that case.
 when $q=2^m$ $m \geq 1$

It is left as an exercise to check that $Z = \{\pm I\}$ where Z is the center of SL_2

Lemma: F field, $\#F \geq 5$. Then $\exists r : \frac{r^2}{r^2-1} \neq 0, 1, -1$.

$$\text{pf: } r^2=0 \Rightarrow r=0 \quad r^2=1 \Rightarrow (r+1)(r-1)=0 \Rightarrow r=1 \text{ or } -1$$

$$a^2=b^2=-1 \Rightarrow (a+b)(a-b)=0 \leftarrow \text{Cauchy Law Holds in Field.} \Rightarrow a=-b \text{ or } a=b$$

at most 2.

$$2+2+1=5.$$

(This Lemma should be put forward when it is needed.)

We prove the case for $\#F \geq 5$; there we can ~~use~~ ^{enjoy} the lemma above.

We take an arbitrary $A \neq \pm I$ and $r \in F$, $r \neq 0, \pm 1$, then we try to "generate" the whole group SL_2 . Or more precisely, to generate the generators of SL_2 .

What kind of generators? $\left\{ \begin{pmatrix} 1 & x \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ x & 1 \end{pmatrix}, x \in F \right\}$



To see why this is true:

$$\begin{pmatrix} a & c \\ b & d \end{pmatrix} \xrightarrow[\text{similarly.}]{\text{for } ct0} \begin{pmatrix} a & \\ d & a^{-1} \end{pmatrix} \Rightarrow \begin{pmatrix} a & \\ & a^{-1} \end{pmatrix} \Rightarrow \begin{pmatrix} a & 1 \\ & a^{-1} \end{pmatrix} \Rightarrow \begin{pmatrix} 1 & 1 \\ a^{-1} & a^{-1} \end{pmatrix} \Rightarrow \begin{pmatrix} 1 & \\ & * \end{pmatrix}$$

Next, How to generate? We use A , and then conjugation operation (action) \rightarrow since we are given a normal subgroup. (i.e. similarity) multiplication and inverse. in $SL_2(\mathbb{R})$ (Not $GL_2(\mathbb{R})$), and then matrix.

There're some tricks here: we think about eigenvalues.

We claim 1 that at least one matrix with s, s^{-1} as e-v is in SL_2 .

And then we claim 2 that all (s, s^{-1}) can be generated by

conjugation in SL_2 . Then we observe that:

$$\begin{pmatrix} s^{-1} & \\ & s \end{pmatrix} \begin{pmatrix} s & \\ & s^{-1} \end{pmatrix} = \begin{pmatrix} 1 & \\ & 1 \end{pmatrix}, \text{ is what we want.}$$

Thus, it remains to prove claim 1 and 2.

1. (Use trick). choose vector $v_1 \in \mathbb{R}^2$ & v_1 is not eigenvector of A .

Then (v_1, Av_1) is invertible ~~and det = 1~~ ~~WLOG~~.

This can be achieved since $A \neq \pm I$.

Let $P = B \begin{pmatrix} r & \\ & r^{-1} \end{pmatrix} B^{-1}$ (recall $r^2 = s \neq \pm 1, 0$) $\Rightarrow P \in SL_2$

Then $P(v_1, Av_1) = PB = B \begin{pmatrix} r & \\ & r^{-1} \end{pmatrix} B^{-1} = B \begin{pmatrix} r & \\ & r^{-1} \end{pmatrix} = (rv_1, r^{-1}Av_1)$.

Then $C = A(PA^{-1}P^{-1})$: ~~WLOG~~
 $C(Av_1) = APA^{-1}P^{-1}Av_1 = rAPA^{-1}Av_1 = rAPv_1 = r^2Av_1 = s(Av_1)$

C is with eigenvalue s and s^{-1} .

2. Note that any matrix M with s, s^{-1} can be conjugated to $\begin{pmatrix} s & \\ & s^{-1} \end{pmatrix}$

with $Q \in GL_2(\mathbb{R})$.

$M = Q^{-1} \begin{pmatrix} s & \\ & s^{-1} \end{pmatrix} Q$. But $\begin{pmatrix} s & \\ & s^{-1} \end{pmatrix}$ can commute with

any diagonal matrix in $GL_2(\mathbb{R})$ thus Q can be choose with $\det = 1$.

Q



§. $\text{Aut}(\frac{\mathbb{Z}}{n\mathbb{Z}})$.

We have learned in class that $\text{Aut}(\frac{\mathbb{Z}}{n\mathbb{Z}}) \cong \begin{cases} \{\pm 1\} & n=2 \\ (\frac{\mathbb{Z}}{n\mathbb{Z}})^\times & \text{otherwise} \end{cases}$

where $(\frac{\mathbb{Z}}{n\mathbb{Z}})^\times$ is the unit group of the ring $\frac{\mathbb{Z}}{n\mathbb{Z}}$. We develop ~~more~~ finer discussions about the group.

Prop: $(\frac{\mathbb{Z}}{p\mathbb{Z}})^\times \cong \frac{\mathbb{Z}}{(p-1)\mathbb{Z}}$ where p is a prime.

Here is a trick.

Let $d \mid p-1$ and assume $r \in (\frac{\mathbb{Z}}{p\mathbb{Z}})^\times$ s.t. $\text{order}(r) = d$.

That is, ~~$r^d = 1$ and r is a root of $f(x) = x^d - 1$~~ ,

~~where $f(x) \in \mathbb{F}_p[x]$. It is easily seen (or you will see it in 2 months)~~

~~f can't claim more than d roots. Thus:~~

$$f_d = \#\{r \in (\frac{\mathbb{Z}}{p\mathbb{Z}})^\times, \text{order}(r) = d\} \leq d$$

$r^d = 1 \Rightarrow r$ is a root of $f(x) = x^d - 1 \in \mathbb{F}_p[x]$.

f has at most d roots and $\langle r \rangle = \mu_d$ ^{all elements of} μ_d are distinguished roots of $f(x)$. Hence these are ~~the~~ all the elements that satisfy $x^d = 1$.

What's more, the only ~~root~~ elements of order d are $\mu_d \setminus \mu_{d'}^*$, the generators of μ_d . And $\#\mu_d^\times = \varphi(d)$.

Hence, let $f_d = \#\{r \in (\frac{\mathbb{Z}}{p\mathbb{Z}})^\times, \text{order}(r) = d\}$

Then $f_d = 0$ or $\varphi(d)$.

But $\sum_{d \mid p-1} f_d = \#\{(\frac{\mathbb{Z}}{p\mathbb{Z}})^\times\} = p-1 = \sum_{d \mid p-1} \varphi(d)$, this makes $f_d = \varphi(d) \forall d \mid p-1$.

In particular, $f_{p-1} = \varphi(p-1) > 0$, which means that $\frac{\mathbb{Z}}{p\mathbb{Z}}$ is cyclic.



Prop 2. Let $p > 2$ and m be a primitive root ~~mod~~ modulo p .

Then: ~~(0)~~

either m or $m+p$ is a primitive root modulo p^2 .

In particular, $(\frac{\mathbb{Z}}{p^2\mathbb{Z}})^{\times}$ is cyclic.

Pf: ~~m~~ m exists due to previous prop.

Then $m^{p-1} \equiv 1 \pmod{p}$ and $m^k \not\equiv 1 \pmod{p}$ for $k < p-1$.

$\Rightarrow \text{ord}_{p^2}(m) \geq p-1$. But $\text{ord}_{p^2}(m) \mid p(p-1)$, where

we know that $\#(\frac{\mathbb{Z}}{p^2\mathbb{Z}})^{\times} = p(p-1)$

Thus either $\text{ord}_{p^2}(m) = p-1$ or $p(p-1)$ by Lagrange Theorem.

If $m^{p-1} \not\equiv 1 \pmod{p^2}$, then we are done.

Otherwise, $m+p$:

first notice that $m+p \equiv m \pmod{p} \xrightarrow{\text{(why?) think about it.}}$ $\text{ord}_{p^2}(p+m) \geq p-1$.

But $(p+m)^{p-1} = \underbrace{m^{p-1}} + (p-1)pm + \dots \equiv m^{p-1} + (p-1)pm \pmod{p^2}$.

$pm \not\equiv 1$ since $p^2 \equiv 0 \pmod{p^2}$. Thus we are done.

Prop 3. Aut $\frac{\mathbb{Z}}{8\mathbb{Z}} \cong \frac{\mathbb{Z}}{2\mathbb{Z}} \times \frac{\mathbb{Z}}{2\mathbb{Z}}$

Pf: Aut $\frac{\mathbb{Z}}{8\mathbb{Z}} = (\frac{\mathbb{Z}}{8\mathbb{Z}})^{\times} = (\{1, 3, 5, 7, \times \pmod{8}\}) = \frac{\mathbb{Z}}{2\mathbb{Z}} \times \frac{\mathbb{Z}}{2\mathbb{Z}}$

Prop: $(\frac{\mathbb{Z}}{2^m\mathbb{Z}})^{\times} \cong \frac{\mathbb{Z}}{2\mathbb{Z}} \times \frac{\mathbb{Z}}{2^{m-2}\mathbb{Z}}$



§. Semi-product

Motivation: Given two generators a, b , what information should be known to get the whole structure of the group? We hope $G = a^m b^n$

① We have learn that let $H, N \triangleleft G$, $H \cap N = \{e\}$, then $H \cong H/N$, $N \cong N/H$, then $H \times N \cong HN$ as in 2.

It still make sense if H fails to be normal (~~$H \triangleleft G$~~)

to write HN . HN is still a subgroup of G .

Notice that $\langle (12) \rangle \cdot \langle (123) \rangle = S_3$ but

$S_3 \neq \mu_2 \times \mu_3$ since S_3 is not abelian.

Can we learn more of structure of G using HN ?

We develop my observations.

① G "uniquely factor" into elements of H and N .

② $|G| = |N| |H|$ and we have: $G = HN$ as sets

③ we hope: $(h_1 n_1)(h_2 n_2) = h_1 n_1 h_2 n_2 = h' n' = (h', n')$

Then we have to commute $n_1 h_2$.

But $n \in N$, normal subgroup! Hence the above becomes

$$h_1 n_1 h_2 n_2 = h_1 h_2 (h_2^{-1} n_1 h_2) n_2 = (h_1 h_2, n_1 h_2 n_2)$$

④ Then Recall: $\varphi: H \rightarrow \text{Aut}(N)$

$$h \mapsto \varphi_h(n) = n^h = h^{-1} n h$$

We have: $\text{Hom}_{\text{Group}}(H, \text{Aut}(N)) \cong \text{Semi}(H, N)$

" φ " is "functorial" On the other hand, given N and H .

We can construct a group G with each $\varphi: H \rightarrow \text{Aut}(N)$



It is checked that: $\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/n\mathbb{Z} \triangleleft H \ltimes \mathbb{Z}/n\mathbb{Z}$.
 left is an exercise to

Conclusion: $G \cong (H \rtimes_{\text{Inn}} \mathbb{Z}/n\mathbb{Z})$. Useful: $(n, h) \leftrightarrow n$ elements in G ,
 but the multiplication is not ~~the~~ the trivial one.

A Basic Example: D_n , the dihedral group of order $2n$.
 (二面体群).

It is convenient to write a reflection by τ , and a rotation by δ .
 (or whatever tradition you like). To understand the group:

① order: $\tau^2 = 1, \delta^n = 1$ and $\langle \tau, \delta \rangle = D_n$.

② How to compute? We would like elements to be the form of $\delta^m \tau^n$, thus we try to find out $\tau \delta = \delta^m \tau^n$. If this min
 min of:
 is got, we can compute it using δ and τ freely.

Fortunately, we have the following facts:

• $\tau \delta \tau = \tau \delta \tau^{-1} = \delta^{-1} = \delta^{n-1}$

the relation between δ and τ
 or "definition"
 $\tau \delta = \delta^{-1} \tau$

• $\langle \delta \rangle \triangleleft D_n$, as the above formula shows (or $[D_n : \langle \delta \rangle] = 2$).
 Anal: thus $D_n = H \rtimes N$, where $H = \langle \tau \rangle, N = \langle \delta \rangle$.

Then we give a ~~finer~~ finer structure of D_n . We try to decide:

- ① the order of every elements.
- ② Every Subgroup / Normal subgroup structure.
- ③ Every conjugate class.

We first explain ③: Given an element $g \in G$, we push it to $h^{-1}gh$
 $\forall h \in G$ to get $\{h^{-1}gh : h \in G\}$. This will lead to a partition (划分) of G .

The equivalent relation correspondingly to that is: $x \sim y \iff \exists z : x = z^{-1}yz$.
 You see that $C_g = \{x : x \text{ "fix" } g, \text{ namely: } g = x^{-1}gx\}$ is the subgroup



of elements that commutes with g . You see that

$$\sum_{\substack{g \text{ over a} \\ \text{representative set}}} \#K(g) = \#G \quad \text{and:}$$

$$\#K(g) = \frac{\#G}{\#C(g)}$$

$K(g)$ means the conjugacy class of g .

In fact, that's ~~the~~ most important action on a set (in fact, ~~groups~~ group on the group itself) in group theory. In the proof of Sylow's theorem, a variant type of it (act on a subgroup class) plays a key role.

On the other hand, computing the conjugacy class of a group is standard if you want to attain a "character table" of G (see GTM/62 or Artin for further).

For ① every element can be written as $\delta^m \tau$ or δ^n .

$$(\delta^m \tau)^2 = \delta^m \tau \delta^m \tau = \delta^m (\tau \delta^m \tau^{-1})^m = \delta^m \delta^m = 1, \quad \text{done.}$$

For ②. Forget it at this step. The key is: how does generators behave?

For ③ Observation: $\tau \delta \tau = \delta^{-1}$, we have to consider odd or even.

Case n is odd. $\delta \neq \delta^{-k} \quad \forall k$. You ~~only~~ ^{first} need to consider how τ or δ push δ^k . $\tau \delta^k \tau^{-1} = \delta^{-k}$, $\delta \delta^k \delta^{-1} = \delta^k$. Thus the conjugacy class of δ^k is $\{\delta^k, \delta^{-k}\}$, as you can check using δ^{-1} .

For ~~δ^k~~ τ and other order 2 elements:

$$\begin{cases} \tau \tau \tau = \tau \\ \tau \tau \tau^{-1} = \tau \end{cases} \quad \text{if } n \text{ is odd}$$

You then see an orbit: $\tau \xrightarrow{\delta} \delta^{-2} \tau \xrightarrow{\delta} \delta^2 \tau \xrightarrow{\delta} \dots \xrightarrow{\delta} \delta^{-2n} \tau = \tau$, n elements here.

They form a conjugacy class. Select a remaining element, $\delta \tau$, for example, you see similar: $\delta \tau \xrightarrow{\delta} \delta^3 \tau \xrightarrow{\delta} \dots \xrightarrow{\delta} \delta^{-2n+1} \tau = \tau$.



For n even, it is left as an exercise.

For ②, at least you can find all normal subgroups using the above observations.

After you learn Sylow's theorem, it is suggested that you do the same ~~central~~ calculations for PQ groups, where p and q are prime.

§ Factorization of a Group: Normal ~~Series~~ Towers.

Of course: $12 = 3 \times 2^2$.

You can do a similar thing for groups: $S_4 \triangleright A_4 \triangleright K_4 \triangleright M_2 \triangleright \{1\}$.

It's hopeless to write $S_4 = \underbrace{S_4}_{A_4} \underbrace{\left(\begin{smallmatrix} \times \\ \circ \end{smallmatrix} \right)}_{K_4} \times \underbrace{K_4}_{M_2} \times \underbrace{M_2}_1$. But there are

still interesting things here: all "n" above are simple groups, namely:

M_2, M_3, M_2, M_2 . In what way do they play the role of

prime decomposition? Is it unique (in some way)?

On the other hand, it's famous that Galois combined the normal series of Galois Group with the extension field to prove the insolubility of polynomials with degree ≥ 5 . Also you will see many "series" here ^{one} after _{after}.

Definition:

① Tower: $G = G_0 \triangleright G_1 \triangleright \dots \triangleright G_m$. Normal tower: " \triangleright " \Rightarrow \triangleright .

② Abelian / Cyclic tower: Normal & G_i/G_{i+1} is cyclic.

③ refinement: just insert some into a given tower.

④ solvable: abelian tower & the ending element is trivial.



⑤ $G = G_1 \supset G_2 \supset \dots \supset G_r = \{e\}$ equivalent if:
 $G = H_1 \supset H_2 \supset \dots \supset H_s = \{e\}$

$r=s$ and $\{G_i\} \exists \tau: \{1, \dots, r\} \rightarrow \{1, \dots, r\}$ s.t.
 $\frac{G_i}{G_{i+1}} \cong \frac{H_{\tau(i)}}{H_{\tau(i)+1}}$

Theorem: G = group. Two normal towers of subgroups ending with the trivial group have equivalent refinements (Schrier)

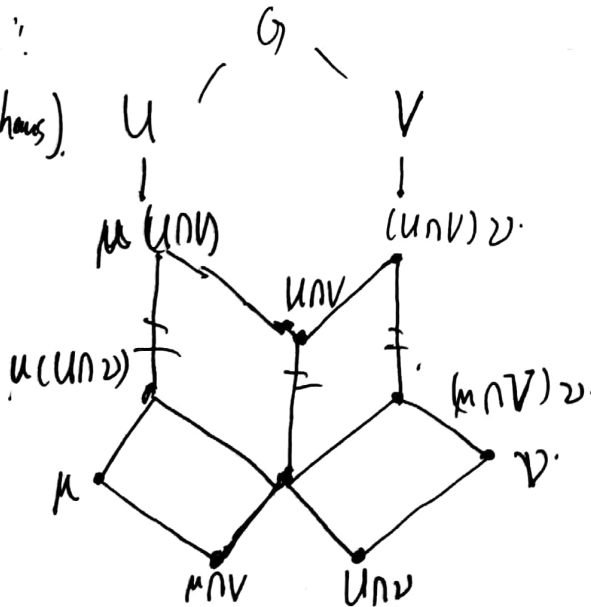
Theorem: $G = G_1 \supset G_2 \supset \dots \supset G_r = \{e\}$, normal tower $\frac{G_i}{G_{i+1}}$ is simple $G_i \neq G_{i+1}$. Any other "factorizations" of G is equivalent to it.

Lemma 1. $f: G \rightarrow G'$; $G' = G'_1 \supset G'_2 \supset \dots \supset G'_m$

Then: Let $G_i = f^{-1}(G'_i)$, where $G = G_0 \supset G_1 \supset \dots \supset G_m$.

"Pull back".

Lemma 2. (Zassenhaus)



$U \triangleleft U \cap V$
 $V \triangleleft (U \cap V) \cap V$

Then: $\mu(U \cap V) \triangleleft \mu(U \cap V)$
 $(\mu(U \cap V)) \cap V \triangleleft (U \cap V) \cap V$
 and: $\frac{\mu(U \cap V)}{\mu(U \cap V)} \cong \frac{(U \cap V) \cap V}{(U \cap V) \cap V}$



Proof: A left point : $(u \cap v)(U \cap v)$. Check : normal.

Check: $(u \cap v)(U \cap v) \triangleleft u \cap v$.

Check: $\frac{\mu(U \cap v)}{\mu(u \cap v)} \stackrel{\sim}{=} \frac{u \cap v}{(u \cap v)(U \cap v)} \stackrel{\sim}{=} \frac{e(U \cap v)}{(u \cap v)}$
 ↗ symmetric.

Use isomorphism theorems.

$$\frac{\mu(U \cap v)}{\mu(u \cap v)} \stackrel{\sim}{=} \frac{(\mu \circ U \cap v)(U \cap v)}{\mu(U \cap v)} \stackrel{\text{2nd}}{=} \frac{u \cap v}{\mu(U \cap v) \cdot \mu(u \cap v)}$$

$$\stackrel{\sim}{=} \frac{u \cap v}{(u \cap v)(U \cap v)}$$

Observation: Every Normal Tower can be refined so that $\frac{G_i}{G_{i+1}}$ is simple

Def of thm 1: $G_1 \triangleright G_2 \triangleright G_3 \triangleright \dots$ (Trick)

insert H_1, H_2, \dots

Namely: $G_1 \triangleright G_{11} = G_2(H_1 \cap G_1) \triangleright G_{12} = G_2(H_2 \cap G_1) \triangleright \dots \triangleright G_2 \triangleright \dots$
 $G_{ij} = G_{i+1}(H_j \cap G_i)$ See butterfly.

Similarly: $G_{ij} = H_{j+1}(G_i \cap H_j)$.

$$\frac{G_{ij}}{G_{i,j+1}} \cong \frac{G_{i+1} H_j}{H_{j+1}} \quad \triangleright \text{By Butterfly.}$$

- $(-1), (s-1), +1$ elements
- end with $\{e\}$

Proof of thm 2: Do refinement as above.

$\forall i \exists \text{ one } j: \frac{G_i}{G_{i+1}} = \frac{G_{ij}}{G_{i,j+1}}$ (You got many)

"trivial refinement"

You may refer to P144 for more about "solvable groups"



习题 1.2 习题解答 (part)

1.2.1. $\varepsilon: f: A \rightarrow G, d \mapsto 1$

逆: $(f^{-1})(a) = (f(a))^{-1}$

1.2.2.

1.2.3 解析几何.

1.2.5. 只有 (3)

1.2.6. $\forall x, y \in \bigcup_{n \geq 1} M_n \quad \exists N: x, y \in M_N$

1.2.7 例子: $\{(M, N) : n \in \mathbb{Z}\}$

1.2.8. ✓

1.2.9. $a, b \in G^X$ 则 $ab^{-1} \in G^X$

1.2.10. $a_1, \dots, a_n \in G \quad |G| = n \Rightarrow \exists 1 \leq p \leq q \leq n: a_p a_{p+1} \dots a_q = 1$

Pf: 考虑集合 $\{a_1, a_2, \dots, a_1 \dots a_n\}$ 若 $1 \in S$, done. 若不然, $\exists i, j: a_i \dots a_i = a_1 \dots a_j$ 从而 $a_i \dots a_j = 1$

1.2.11 若不然, 取 $a \in A \setminus B \quad b \in B \setminus A$ 则 $a^{-1} \in A \quad b^{-1} \in B \Rightarrow ab \notin A \cup B$, 矛盾.

1.2.12 $x^2 = 1$ 偶数阶群. 若有偶数个解.

~~G~~ $G = \{1\} \cup \{2 \text{阶元}\}$ 是偶数 $[g \text{ 和 } g^{-1} \text{ 配对}]$
 $\cup \{x^2 = 1\}$

1.2.13. $O_{p,q}(\mathbb{R}) := \{A \in GL_n(\mathbb{R}) : A^T \begin{pmatrix} I_p & \\ & -I_q \end{pmatrix} A = \begin{pmatrix} I_p & \\ & -I_q \end{pmatrix}\}$

Pf: $A^{-1} = A^{-T} \quad \begin{pmatrix} I_p & \\ & -I_q \end{pmatrix}^{-1} = \begin{pmatrix} I_p & \\ & -I_q \end{pmatrix}$

$S_{p,2n}(\mathbb{R}) = \{A : A^T \begin{pmatrix} 0 & I_n \\ -I_n & 0 \end{pmatrix} A = \begin{pmatrix} -I_n & \\ & I_n \end{pmatrix}\}$

1.2.14, 1.2.15 : easy (由 1.2.11)

1.2.16. (\Rightarrow) AB 是 G 的子群 \Rightarrow ~~$(ab)^{-1} = b^{-1}a^{-1} \in AB$~~
 $ba = [a^{-1}b^{-1}]^{-1} \in AB \Rightarrow BA \subseteq AB$

$AB = (AB)^{-1} = B^{-1}A^{-1} = BA$

(\Leftarrow) $AB = BA \Rightarrow \forall a_1, a_2 \in A, b_1, b_2 \in B: (b_1 b_2)(b_2^{-1} a_2^{-1}) = a_1 (b_1 b_2^{-1}) a_2^{-1} = a_1 a_2 b_2^{-1} b_1 \in AB$



1.2.17 $A \rightarrow A^{-1} \rightarrow A^{-1}g$ 是双射,

从而 $|A^{-1}g| + |B| > |g|$

$\Rightarrow A^{-1}g \cap B \neq \emptyset \Rightarrow \exists a \in A \quad b \in B \quad a^{-1}g = b \quad i.e. g = a \cdot b$

1.2.18. (1) ~~洛朗级数~~ 利用欧几里德除法 + 最小 argument.

(2) $\frac{U}{h\mathbb{Z}} = \langle I \rangle_n = \langle M_n \rangle$ 故所有子群为 $\langle M_n^k \rangle \quad k=1, \dots, n.$

1.2.19. 1.2.20 easy.

1.2.21 (1) $H \times K \rightarrow \mathbb{R}^x$
 $(a, b) \mapsto ab.$

(2) $\text{Diag}_n(\mathbb{F}) \times T_n(\mathbb{F}) \rightarrow B_n(\mathbb{F})$
 $(\lambda I, A) \mapsto \lambda A.$

(3) $H \times K \rightarrow G.$
 $(e^{i\theta}, r) \mapsto re^{i\theta} \quad r > 0.$

1.2.22. ~~(R, +) 中每个元素都是有限阶元~~ (R, +) 中的有限阶元 \checkmark
~~(R, \times) 中, 1 不是有限阶元~~ (R, \times) 中有限阶元: 有限且

1.2.23. $(a, b): a \neq 0$
 $\varphi: (a, b) \mapsto \begin{pmatrix} a & b \\ & 1 \end{pmatrix}.$

$$\begin{pmatrix} a & b \\ & 1 \end{pmatrix} \begin{pmatrix} c & d \\ & 1 \end{pmatrix} = \begin{pmatrix} ac & ad+bc \\ & 1 \end{pmatrix} = \varphi(ac, ad+bc) \\ = \varphi((a, b) \cdot (c, d))$$

1.2.24. G 群有一无不动点的自同构 α , 且 $\alpha^2 = I$. 证明:

G 是奇数阶阿贝尔群.

pf: 只要证明 $\alpha(g) = g^{-1} \quad \forall g \in G.$ 单满
 为此注意到 α 无不动点, 从而 $h \mapsto \alpha(h)h^{-1}$ 是双射, 故 $g = \alpha(h)h^{-1} \quad (\exists h)$
 两边用 α 作用, 得 $\alpha(g) = \alpha(\alpha(h)h^{-1})^{-1} = g^{-1}$ 得证.



1.3.1. $A = 4 \times n$ $B = 3 \times n$ AB : 无穷 BA : 无穷阶.

1.3.4. $k=0 \dots 1$ 不全: $d \in \bar{n} \bar{k}^{-1}$, $d \leq n$.

$$1.3.6. f^2 = 1_G \quad g^3 = \frac{\frac{x-1}{x} - 1}{\frac{x-1}{x}} = 1$$

Check: $f g f^{-1} = g^{-1}$.

1.3.7. (1). $g = e^{i\theta}$ 取 θ 最小的那个元子, 它是生成元.

(2) \mathbb{Q} 不是循环群: 设 $\mathbb{Q} = \langle \frac{1}{p} \rangle$ \mathbb{R} 对 $\frac{1}{\pi}$ 无法生成. 注意: 无穷阶
有限生成子群: $\langle \mathbb{Q} \frac{1}{a_1}, \dots, \frac{1}{a_n} \rangle$ 取 $g = \frac{1}{\text{lcm}(a_1, \dots, a_n)}$ 即为生成元.

(3) 注意到 G 的子群有链结构:

$$\{1\} \leq G_1 \leq G_2 \leq G_3 \leq \dots$$

$$G_i = \{x \in G : x^{p^i} = 1\}$$

$$1.3.3 \quad aba^{-1} = a^m b a^{-m} = a^3 \dots (a^m b a^{-m}) \dots a^{-3} = b$$

$$\Rightarrow ab = ba.$$

$$1.3.8. (ab)^k = 1 \Rightarrow a^k b^k = 1. \quad \begin{matrix} k \mid m \\ \Rightarrow m \mid k, n \mid k. \end{matrix} \quad \begin{matrix} m, n \mid k \\ \Rightarrow m \mid k \end{matrix} \quad \& \quad (ab)^{mn} = 1$$

$$\Rightarrow \text{order of } ab \text{ is } 1.$$

1.3.9. TBD

$$1.3.11. \quad n = \sum_{d|n} \varphi(d). \quad = \sum_{d|n} \varphi(d). \quad \varphi(d) \text{ 是 } d \text{ 阶元的个数}$$

则由 $\varphi(d) \leq \varphi(d)$ [d阶元有一个 \Rightarrow 有 $\varphi(d)$ 个]. 由条件, $\varphi(d) \leq d$.
故 $\varphi(d) = \varphi(d)^0$, 故 $\varphi(n) = 1$



1.3.12: ~~no~~ \mathbb{Z} .

1.3.13. (1). $(ab^{-1})^{mn} = a^m b^{-n} = 0$.

(2) 由 1.3.1 利得.

1.3.14 (1) Abel $\Rightarrow \varphi \in \text{End}(G)$.

$x^2 = y^2 \Rightarrow (x-y)^2 = 1 \xrightarrow{\text{odd}} x=y$.

injective & finite \Rightarrow surjective.

(2). $\varphi_d(x) = x^d$. $d(n) = 1$ when $n = \phi(G) < \infty$.

1.3.15. G abel $\alpha \in \text{Aut}(G)$ $\alpha\alpha = \text{id}$.

$G_1 = \{g \in G : \alpha(g) = g\}$ $G_{-1} = \{g \in G : \alpha(g) = g^{-1}\}$

(1). $\alpha(g) = g = g^{-1} \Rightarrow g = 1$ since $|G|$ odd.

$g = [g\alpha(g^{-1})]_h = a(g)$
 $\alpha(h) = a(g)g^{-1} = h^{-1}$ since $\alpha \circ \alpha = 1$.

(2) 特别 $\exists! h \in G : h^2 = 1 \Rightarrow h = 1 \Rightarrow$ no non-trivial order-2-element.

$\Rightarrow G \cap G_{-1} = 1$

(i) (ii): 取 α 为转置, $x \mapsto -x$.

1.3.16. (1) $\text{Aut}(\mathbb{Q}, +) : \varphi \in \text{Aut}(\mathbb{Q}, +)$

Classical Problem for Beginners determine $GL(\mathbb{F}_p)$?

$\varphi \Rightarrow \varphi(1) = g$ Then $g \leftrightarrow \varphi$ correspondence

$\varphi \circ \psi(1) = g\psi(1) = gq$ Thus $\text{Aut}(\mathbb{Q}, +) \cong \mathbb{Q}^*$.

(2). $\varphi(1) = \pm 1 \Rightarrow \text{Aut}(\mathbb{Z}) \cong \mu_2$.

(3). ~~$\mathbb{Z} \times \mathbb{Z}$~~ $\frac{\mathbb{Z}}{2\mathbb{Z}} \times \frac{\mathbb{Z}}{2\mathbb{Z}} \cong \mathbb{F}_2$ $\text{Aut}(\mathbb{Q}) \cong GL_{\mathbb{Z}}(\mathbb{F}_2)$ 加法群同构, 即子模同构

$\cong GL_{\mathbb{F}_2}(\mathbb{F}_2)$: And hence we consider Matrix Group of.

We have: $\{(0 \ 1), (1 \ 1), (1 \ 1), (1 \ 1), (1 \ 1), (1 \ 1)\} \cong S_3$



$$A f_i \rightarrow x f_i$$

New

1.4.1. $(a, b)(c, d) = (ac, ad+bc)$ $a \in \mathbb{R}^x, b \in \mathbb{R}$
 注意到 $G \cong \left\{ \begin{pmatrix} a & b \\ 0 & 1 \end{pmatrix} \dots \right\}$

1.4.2 $G = GL_n(\mathbb{R})$ $M = \det > 0$
 $\frac{G}{M} \cong G \rightarrow \mathbb{R} \rightarrow \mathbb{R}_+$

1.4.3. (2) 反例: $K_4 \triangleleft A_4 \triangleleft S_4$ 但 $K_4 \not\triangleleft S_4$

1.4.4. (1) $aba^{-1}b^{-1} \stackrel{=1}{\neq} \Rightarrow \tau^{-1}a\tau b \tau^{-1}a^{-1}\tau b^{-1}\tau^{-1} = \tau^{-1}aa^{-1}\tau = 1$
 (2) 注意到 $aN = Na$ $\forall a$

1.4.5. 定义

1.4.6. 反证法 $Z(G) \neq \{e\} \Rightarrow \frac{G}{Z(G)} = \langle \text{非平凡} \rangle \langle a + Z(G) \rangle$

claim: $ab=ba \forall a, b$ 是矛盾.
 $a \in Z(G)$.
 • $\forall a, b \in Z(G) \Rightarrow ab=ba$
 • $\forall b \notin Z(G) \Rightarrow b = a^n c$ $c \in Z(G)$ 由循环群
 $\Rightarrow ab = a a^n c = a^{n+1} c = (a^{n+1}) a = b a$

1.4.7 (外直积) 显然.

1.4.8. (2) $\text{Inn}(G) = \langle \text{conjugation} \rangle = \langle \text{inner automorphisms} \rangle = \langle \text{conjugation} \rangle$
 (3) $\text{I}(G) \cong \frac{G}{Z(G)}$; $\bullet G \rightarrow Z(G)$
 $x \mapsto z_x$

1.4.9. $GL_n(\mathbb{R})$: 先用 $I+E_{ij}$, 再用 $I+E_{ii}$ 作为 "test function".





1.4.11. (1) $M \cap N = \{1\}$ $a \in M, b \in N$ $ab = ba$.
 注意到 $aba^{-1}b^{-1} \in M \cap N$

(2) $\phi: M \times N \rightarrow G$ $(m, n) \mapsto mn$. Check: ϕ 是同态 (用到了 (1))
 $\ker \phi = 1.$

1.4.12. $\text{ord}(g) \mid m + n \mid |G| = 1$.
 $\Rightarrow g = g^{\text{ord}(g)} = g^{m+n} = (g^m)^{n+1} \in N$ why?

1.4.13 利用 1.4.6 和 1.4.8 立得.

1.4.14 TBP

1.3, 20 k 1.3-24

(加) 用群作用更容易理解: $G/K \cong G$ 关于 K 的左陪集.

- H act on $G/K \Rightarrow HxK = \bigcup_{g \in H} gK$ (轨道中并起来)
- $\text{Stab}(xK) = \{g \in H : gxK = xK\} = \{g \in H : x^{-1}gx \in K\} = H \cap xKx^{-1}$.

从而: $|HxK| = |H| \cdot \#\text{orbit} = |K| \cdot \frac{|H|}{|\text{Stab}|} = \frac{|K||H|}{|H \cap xKx^{-1}|} = \frac{|K||H|}{|x^{-1}Hx \cap K|}$

$[g(sNT)g^{-1} = gsg^{-1}NgTg^{-1} \quad \forall g \in G, s, T \in G]$

1.3.24 $x \sim y \Leftrightarrow x \in HyK. \Rightarrow G = \bigcup_{g \in G} HgK$. $\{b_i g a_i, g_i\}$ units
 $G = \bigcup_{g \in G} A_g A \quad A_g A = \bigcup_{g \in G} A g a_i = \bigcup_{g \in G} b_i g A$ 而 $A = A b_i = a_i A$
 $\Rightarrow A A = \bigcup_{g \in G} A b_i g a_i = b_i g A$

