Subtyping (Dynamic Polymorphism)

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References

• PFPL

- Chapter 24 Structural Subtyping
- Chapter 27 Inheritance
- <u>TAPL</u> (pdf)
 - Chapter 15 Subtyping
- [Concepts in PLs]

Subtyping and Inheritance

- Interface
 - The external view of an object
- Subtyping
 - Relation between interfaces
- Implementation
 - The internal representation of an object
- Inheritance
 - Relation between implementations

Various Object-Oriented Languages

- Pure dynamically-typed OO languages
 - Object implementation and run-time lookup
 - Class-based languages (Smalltalk)
 - Prototype-based languages (Self, JavaScript)
- Statically-typed OO languages
 - C++
 - using static typing to eliminate search
 - problems with C++ multiple inheritance
 - Java
 - using Interfaces to avoid multiple inheritance

Smalltalk: Subtyping

• If interface A contains all of interface B, then A objects can also be used B objects.

Point	ColorPoint
x:y:	x:y:
moveDx:Dy:	moveDx:Dy:
X	Х
У	У
draw	color
	draw

ColorPoint interface contains Point ColorPoint is a subtype of Point

Subtyping and Inheritance

- Smalltalk/JavaScript subtyping is implicit
 - Not a part of the programming language
 - Important aspect of how systems are built

- Inheritance is explicit
 - Used to implement systems
 - No forced relationship to subtyping

C++

- C++ is an object-oriented extension of C, Bell Labs
- Object-oriented features
 - Classes
 - Objects, with dynamic lookup of virtual functions
 - Inheritance
 - Single and multiple inheritance
 - Public and private base classes
 - Subtyping
 - Tied to inheritance mechanism
 - Encapsulation
 - Public, private, protected visibility

C++: Virtual functions

- Member functions are either
 - Virtual, if explicitly declared or inherited as virtual
 - Non-virtual otherwise
- Virtual functions
 - Accessed by indirection through ptr in object
 - May be redefined in derived (sub) classes
- Non-virtual functions
 - Are called in the usual way. Just ordinary functions.
 - Cannot redefine in derived classes (except overloading)
- Pay overhead only if you use virtual functions

C++ Subtyping

- Subtyping in principle
 - A <: B if every A object can be used without type error whenever a B object is required
- C++: A <: B if class A has public base class B
 - Independent classes not subtypes

Java

- 1990-95 James Gosling and others at Sun
- Syntax similar to C++
- Object
 - has fields and methods
 - is allocated on heap, not run-time stack
 - accessible through reference (only ptr assignment)
 - garbage collected
- Dynamic lookup
 - Similar in behavior to other languages
 - Static typing => more efficient than Smalltalk
 - Dynamic linking, interfaces => slower than C++

Inheritance

- Similar to Smalltalk, C++
- Subclass inherits from superclass
 - Single inheritance only (but Java has interfaces)
- Some additional features
 - Conventions regarding *super* in constructor and *finalize* methods
 - Final classes and methods cannot be redefined

Interfaces vs Multiple Inheritance

- C++ multiple inheritance
 - A single class may inherit from two base classes
 - Constraints of C++ require derived class representation to resemble *all* base classes
- Java interfaces
 - A single class may implement two interfaces
 - No inheritance (of implementation) involved
 - Java implementation (discussed later) does not require similarity between class representations

Subtyping Principles

• Subtyping judgement $\tau' <: \tau$

 $\overline{\tau <: \tau}$

- Subsumption rule
- Numeric types

$$\frac{\Gamma \vdash e: \tau' \quad \tau' <: \tau}{\Gamma \vdash e: \tau}$$

 $\tau'' <: \tau' \quad \tau' <: \tau$

 $\tau'' <: \tau$

- int <: rat <: real</pre>
- Product types, Sum types

 $\frac{J \subseteq I}{\langle \tau_i \rangle_{i \in I} <: \langle \tau_j \rangle_{j \in J}}$

$$\frac{J \subseteq I}{[\tau_i]_{i \in I} <: [\tau_j]_{j \in J}}$$

Width subtyping (较宽积类型是较窄积类型的子类型)

Subtyping Principles

- Variance:
 - Product and sum types: Depth subtyping (Covariance)

 $\frac{\tau'_{i} <: \tau_{i} (\forall i \in I)}{\langle \tau'_{i} \rangle_{i \in I}} \qquad \qquad \frac{\tau'_{i} <: \tau_{i} (\forall i \in I)}{\left[\tau'_{i}\right]_{i \in I} <: \left[\tau_{i}\right]_{i \in I}}$

- Partial function types
 - covariant in its range. $\frac{\tau'_2 <: \tau_2}{\tau_1 \rightharpoonup \tau'_2 <: \tau_1 \rightharpoonup \tau_2}$
 - contravariant in its domain position

$$\frac{\tau_1' <: \tau_1}{\tau_1 \rightharpoonup \tau_2 <: \tau_1' \rightharpoonup \tau_2}$$

Subtyping Principles

Quantified Types

 $\frac{\Delta, t \text{ type } \vdash \tau' <: \tau}{\Delta \vdash \forall (t, \tau') <: \forall (t, \tau)} \qquad \qquad \frac{\Delta, t \text{ type } \vdash \tau' <: \tau}{\Delta \vdash \exists (t, \tau') <: \exists (t, \tau)}$

- Substitution: If Δ, t type $\vdash \tau_1 <: \tau_2$, and $\Delta \vdash \tau$ type, then $\Delta \vdash [\tau/t]\tau_1 <: [\tau/t]\tau_2$

THANKS