Memory Management

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Outline

- Memory Management and Problems
- Some Solutions
  - Coding standards
  - Region-based memory management
  - Smart pointer
  - Garbage collection
- Rust Language

References

- Region-based memory management
- Smart Pointers in C++
- Garbage Collection, Richard Jones
- Inside the JVM
- Rust: book, github

Memory Management

- Static data section
- Stack: stack frame
- Heap: store dynamically-allocated objects
  - C: malloc, free
    - Glibc's ptmalloc, Doug Lea's dlmalloc
    - Efficient concurrent memory allocator
      - jemalloc, TBBmalloc, TCMalloc (gperftools)
  - Java: new、**Garbage Collection**
    - Richard Jones's the Garbage Collection Page

Memory Management Problems

- Safety
  - Do not access freed memory
    - Dangling references to Stack frames or Heap
    - double free, use after free

    ➔ Abnormal program termination, denial-of-service attacks
  - Free dynamically allocated memory when no longer needed
    - Memory leak

    ➔ denial-of-service attack
  - Allocate and copy structures containing a flexible array member dynamically

    ➔Undefined behavior
  - Only free memory allocated dynamically

    ➔ Depend on the implementation

Memory Management Problems

- Safety (cont'd)
  - Allocate sufficient memory for an object

    ➔Buffer overflows, the execution of arbitrary code vulnerabilities
  - ……
- Performance
  - Memory hierarchy
    - Program locality (temporal, spatial)
  - Memory allocator
    - VM-intensive: frequent memory allocation and deallocation

Some Solutions

- Coding standards
- Region-based memory management
- Smart pointers (modern C++)
- Garbage collection

Coding Standards

- SEI CERT Coding Standards (CMU)
  - Android、C、C++、Java、Perl
  - Safety and Security Coding Standards for C，2016

    Author: Robert C. Seacord (SEI, CMU)
    - Securing Coding in C and C++, 2nd Edition, 2013
    - C安全编码标准，译著，机械工业出版社，2015
- C++ Core Guidelines，
  - Editors：Bjarne Stroustrup，Herb Sutter
- MISRA publications (UK): C/C++
  - MISRA: Motor Industry Software Reliability Association
- MITRE (USA)
  - Common Weakness Enumeration 公共弱点（缺陷）枚举
  - Common Vulnerabilities and Exposures 公共漏洞和暴露

Region-based Memory Management

- Region (or zone, arena, memory context)
  - A collection of allocated objects that can be efficiently ***deallocated all at once***

    ➔ **time performance**
- History and concepts
  - 1967 Douglas T. Ross's AED Free Storage Package
  - [SPP 1990] David R. Hanson, explicit regions in C
  - Used in Apache HTTP Server, PostgreSQL, etc.

  but, do not provide memory safety

  ➔ Memory leak, dangling pointer

Region-based Memory Management

- Region inference ➔ safe memory allocation
  - [POPL1988]Ruggieri and Murtagh
    - a *region* is created at the beginning of each function and deallocated at the end
    - use *data flow analysis* to determine a lifetime for each static allocation expression, and assign it to the youngest region
  - [POPL 1994] Tofte and Talpin

    Polymorphic region type and region calculus, used in SML
    - Extended lambda calculus including regions

      e_1 at ρ: Compute the result of the expression e_1 and store it in region ρ;

      letregion ρ in e_2 end: Create a region and bind it to ρ; evaluate e_2; then deallocate the region.

Region-based Memory Management

- Generalization to Other PLs
  - C
    - Cyclone [PLDI 2002], RC, Control-C[CASES 2002]
  - Java
    - Real time Java, combined with ownership types
  - Logic PLs such as Prolog, Mercury
- Disadvantages
  - A large proportion of dead data in large regions
  - Shorter-lifetime regions: difficulty in region inference

Smart Pointers

Enable automatic, exception-safe, object lifetime management

- Dynamic memory management in C++
  - Pointer categories, implemented as class templates

    An object can only be referenced by a single smart pointer

    - **auto_ptr**(removed in C++17): **strict** object **ownership** semantics
    - **unique_ptr**(C++11): **unique** object **ownership** semantics
    - **shared_ptr**(C++11): **shared** object **ownership** semantics
      - Reference counted pointer
    - **weak_ptr**(C++11): weak reference to an object managed by std::shared_ptr, call **wp.*lock*()** to check whether the object is deleted
      - Used to break circular references of std::shared_ptr.
- Smart Pointers - What, Why, Which?

Garbage Collection

- Garbage
  - allocated space that is no longer usable by the program

```
let x = [[1; 2; 3]; [4]] in
let y = [2] :: List.tl x in y
```

    - x is never used again and becomes garbage
- Reachability
  - Roots: pointers that appear in the env.
  - GC: reclaims blocks that are no longer reachable from a set of roots
  - Heap: a directed graph in which the nodes are blocks of memory and the edges are the pointers between these blocks.
  - Reachability: computed as a graph traversal

Why Garbage Collection (GC)

- Eliminate a common source of defects
  - Storage leak
  - Dangling pointer
- Improve abstraction and modularity
  - A class need not include code to deal with storage deallocation
- Programs written in OCaml and Java require GC

GC Techniques

- Requirements
  - Should identify most garbage
  - Anything it identifies as garbage must be garbage
  - Should impose a low added time overhead
  - Program pauses made by GC should be short
- Basic GC Techniques
  - Reference counting
  - Mark-and-sweep: mark all reachable objects from a set of roots, and *sweep* through memory, deallocating all unmarked objects
  - Copying collection: *copy* - move reachable objects from the heap to a new area called the *to-space*
  - Mark-and-compact: *compact* – compute forwarding addresses, update pointers and relocate blocks

Identifying Pointers

- **Way1**: Reserve a tag bit in each word
  - Use up 3% memory
  - Limit the range of integers (and pointers): a 32-bit machine can address about 2GB (2^{31})
  - Small run-time cost: arithmetic or dereference
- **Way2**: Have the compiler record info that the GC can query to find out the types of locations
  - More complicated: tightly coupling the GC and compiler
- **Way3**: GC considers memory unreachable only if there is nothing that looks like it might be a pointer to it
  - Work well in practice: integers are small, pointers look like large integers

Other Issues on GC

- Traversing the heap
  - Recursive traversal: hard to do when low on free space
  - Traversal *without* recursion or external stack:
- Program pauses
  - Generational GC: new and old (long-lived) generations, minor GC (only scan new generation), major GC
    - *Intergenerational* pointers: pointers located in an old-generation block that points to a new-generation block
  - Incremental GC: let the garbage collector run *concurrently* with the rest of the program, instead of pausing the program
    - enabling ***predictable real-time*** performance
    - **Complicated synchronization** needed between the garbage collector and the rest of the program

Garbage Collection and Java

- Parallel GC vs. Concurrent GC
  - Parallel GC: a stop-the-world, multithreaded collector
  - Concurrent GC: a mostly concurrent, low-pause collector
- Hashcodes
  - Object.hashCode(): *typically* the address of the heap block
- Finalization
  - finalize(): the GC calls it before reclaiming an object's block
- Package java.lang.ref
  - GC cannot reclaim ***strong*** references
  - GC can reclaim ***soft*** reference (内存不足时被回收), ***weak*** reference (一旦发现即回收), ***phantom*** reference (加到ReferenceQueue中，使程序可以对队列中引用的对象在回收前采取行动)

内存安全及性能的重要性

Google 2020 全年

电	15.5 太瓦时 总计 30.85 亿美元
软件开发	27,169 名软件工程师 总计 48.56 亿美元

提升软件**开发产能**和**软件性能**

[20220413] How Google plans to use 100% carbon-free energy in its data centers by 2030

Google Environmental Report 2021

编程语言及其内存安全是选择的关键

- **手工内存管理**(MMM)：C、C++等
  需要静态分析工具、Sanitizers来检查安全性
- **所有权转移及检查**(OTC)：Rust等
  Borrow检查很慢，但能更好地支持并发
- **垃圾收集**(GC)：Java, Go, Python, JavaScript等
  开发产能最佳但运行慢；Go高效编译和运行
- **引用计数**（RC）：Swift等
  开发产能佳，但需破RC环

PL的内存管理演变

1972

系统编程
(指针/cast)

1980

抽象ADT
性能
复用

1995

WWW
JVM
并行

2009

云计算
DevOps等

编译型语言+MMM → **字节码虚拟机 (即时编译 + GC)** → **编译型语言+GC**

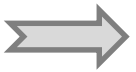
性能高
安全性差

开发快，支持跨平台
性能差

开发快、编译快、性能较高
更灵活广泛的表达能力?

多厂商
异构计算

LLVM
后端



JLang

Ahead-Of-Time 编译

Asul Falcon 云原生
JIT编译

GoLLVM

THANKS