



中国科学技术大学
University of Science and Technology of China

Subtyping

(Dynamic Polymorphism)

《程序语言设计和程序分析》

张昱

yuzhang@ustc.edu.cn

中国科学技术大学
计算机科学与技术学院



[PFPL](#)

- Chapter 24 Structural Subtyping
- Chapter 27 Inheritance

[TAPL \(pdf\)](#)

- Chapter 15 Subtyping

[\[Concepts in PLs\]](#)



Subtyping and Inheritance

Interface

- The external view of an object

Subtyping

- Relation between interfaces

Implementation

- The internal representation of an object

Inheritance

- Relation between implementations



Pure dynamically-typed OO languages

- Object implementation and run-time lookup
- Class-based languages (Smalltalk)
- Prototype-based languages (Self, JavaScript)

Statically-typed OO languages

- C++
 - using static typing to eliminate search
 - problems with C++ multiple inheritance
- Java
 - using Interfaces to avoid multiple inheritance



Smalltalk: Subtyping

- If interface **A** contains all of interface **B**, then **A** objects can also be used **B** objects.

Point	ColorPoint
x:y:	x:y:
moveDx:Dy:	moveDx:Dy:
x	x
y	y
draw	color
	draw

ColorPoint interface contains Point
ColorPoint is a subtype of Point



□ Smalltalk/JavaScript subtyping is implicit

- Not a part of the programming language
- Important aspect of how systems are built

□ Inheritance is explicit

- Used to implement systems
- No forced relationship to subtyping



C++

- **C++ is an object-oriented extension of C, Bell Labs**
- **Object-oriented features**
 - **Classes**
 - **Objects, with dynamic lookup of virtual functions**
 - **Inheritance**
 - Single and multiple inheritance
 - Public and private base classes
 - **Subtyping**
 - Tied to **inheritance** mechanism
 - **Encapsulation**
 - Public, private, protected visibility



- **Member functions are either**
 - Virtual, if explicitly declared or inherited as virtual
 - Non-virtual otherwise
- **Virtual functions**
 - Accessed by indirection through ptr in object
 - May be redefined in derived (sub) classes
- **Non-virtual functions**
 - Are called in the usual way. *Just ordinary functions.*
 - Cannot redefine in derived classes (except overloading)
- **Pay overhead only if you use virtual functions**



□ Subtyping in principle

- $A <: B$ if every A object can be used without type error whenever a B object is required

□ C++: $A <: B$ if class A has public base class B

- Independent classes not subtypes



- **1990-95 James Gosling and others at Sun**
- **Syntax similar to C++**
- **Object**
 - has fields and methods
 - is allocated on heap, not run-time stack
 - accessible through reference (only ptr assignment)
 - garbage collected
- **Dynamic lookup**
 - Similar in behavior to other languages
 - Static typing => more efficient than Smalltalk
 - Dynamic linking, interfaces => slower than C++



- **Similar to Smalltalk, C++**
- **Subclass inherits from superclass**
 - Single inheritance only (but Java has interfaces)
- **Some additional features**
 - Conventions regarding *super* in constructor and *finalize* methods
 - Final classes and methods cannot be redefined



Interfaces vs Multiple Inheritance

□ C++ multiple inheritance

- A single class may inherit from two base classes
- Constraints of C++ require derived class representation to resemble *all* base classes

□ Java interfaces

- A single class may implement two interfaces
- No inheritance (of implementation) involved
- Java implementation (discussed later) does not require similarity between class representations



□ Subtyping judgement $\tau' <:\tau$



$$\overline{\tau <:\tau}$$

$$\frac{\tau'' <:\tau' \quad \tau' <:\tau}{\tau'' <:\tau}$$

□ Subsumption rule

$$\frac{\Gamma \vdash e:\tau' \quad \tau' <:\tau}{\Gamma \vdash e:\tau}$$

□ Numeric types

■ $\text{int} <:\text{rat} <:\text{real}$

□ Product types, Sum types

$$\frac{J \subseteq I}{\langle \tau_i \rangle_{i \in I} <:\langle \tau_j \rangle_{j \in J}}$$

$$\frac{J \subseteq I}{[\tau_i]_{i \in I} <:[\tau_j]_{j \in J}}$$

Width subtyping

(较宽积类型是较窄积类型的子类型)



□ Variance:

- Product and sum types: Depth subtyping (Covariance)

$$\frac{\tau'_i <: \tau_i (\forall i \in I)}{\langle \tau'_i \rangle_{i \in I} <: \langle \tau_i \rangle_{i \in I}}$$

$$\frac{\tau'_i <: \tau_i (\forall i \in I)}{[\tau'_i]_{i \in I} <: [\tau_i]_{i \in I}}$$

- Partial function types

- covariant in its range.

$$\frac{\tau'_2 <: \tau_2}{\tau_1 \rightarrow \tau'_2 <: \tau_1 \rightarrow \tau_2}$$

- contravariant in its domain position

$$\frac{\tau'_1 <: \tau_1}{\tau_1 \rightarrow \tau_2 <: \tau'_1 \rightarrow \tau_2}$$



□ Quantified Types

$$\frac{\Delta, t \text{ type} \vdash \tau' <: \tau}{\Delta \vdash \forall(t. \tau') <: \forall(t. \tau)}$$

$$\frac{\Delta, t \text{ type} \vdash \tau' <: \tau}{\Delta \vdash \exists(t. \tau') <: \exists(t. \tau)}$$

- Substitution: If $\Delta, t \text{ type} \vdash \tau_1 <: \tau_2$, and $\Delta \vdash \tau \text{ type}$, then $\Delta \vdash [\tau/t]\tau_1 <: [\tau/t]\tau_2$



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